

# Lecture 9

## Solutions to exercises

2013

## Exercise

We call a relation  $r \subseteq S \times S$  functional if  $\forall x, y, z \in S. (x, y) \in r \wedge (x, z) \in r \rightarrow y = z$ .

For each of the following statements either give a counterexample or prove it.

In the following, assume  $Q \subset S$ .

- (i) for any  $r$ ,  $wp(r, S \setminus Q) = S \setminus wp(r, Q)$
- (ii) if  $r$  is functional,  $wp(r, S \setminus Q) = S \setminus wp(r, Q)$
- (iii) for any  $r$ ,  $wp(r, Q) = sp(Q, r^{-1})$
- (iv) if  $r$  is functional,  $wp(r, Q) = sp(Q, r^{-1})$
- (v) for any  $r$ ,  $wp(r, Q_1 \cup Q_2) = wp(r, Q_1) \cup wp(r, Q_2)$
- (vi) if  $r$  is functional,  $wp(r, Q_1 \cup Q_2) = wp(r, Q_1) \cup wp(r, Q_2)$
- (vii) for any  $r$ ,  $wp(r_1 \cup r_2, Q) = wp(r_1, Q) \cup wp(r_2, Q)$
- (viii) Alice has the following conjecture: For all sets  $S$  and relations  $r \subseteq S \times S$  it holds:

$$(S \neq \emptyset \wedge dom(r) = S \wedge \Delta_S \cap r = \emptyset) \rightarrow (r \circ r \cap ((S \times S) \setminus r) \neq \emptyset)$$

She tried many sets and relations and did not find any counterexample. Is her conjecture true?

If so, prove it, otherwise provide a counterexample for which  $S$  is smallest.

## Exercise Solution

- (i) counterexample:  $S = \{a, b, c, d\}$  and  $r = \{(a, b), (b, c), (c, a)\}$
- (ii) as above,  $r$  is functional
- (iii) counterexample:  $S = \{a, b\}$  and  $r = \{(a, b)\}$   $r^{-1} = \{(b, a)\}$   $Q = \{b\}$
- (iv) counterexample as in (iii)
- (v) counterexample:  $S = \{a, b\}$  and  $r = \{(a, b), (a, a)\}$  and  $Q_1 = \{a\}$  and  $Q_2 = \{b\}$
- (vi) true:

$$\begin{aligned}wp(r, Q_1 \cup Q_2) &\Leftrightarrow \{s \mid \forall s'. (s, s') \in r \rightarrow (s' \in Q_1 \vee s' \in Q_2)\} \\ &\Leftrightarrow \{s \mid \forall s'. (s, s') \notin r \vee s' \in Q_1 \vee s' \in Q_2\} \\ &\Leftrightarrow \{s \mid \exists s'. (s, s') \notin r \vee s' \in Q_1 \vee s' \in Q_2\} \text{ (since } r \text{ functional)} \\ &\Leftrightarrow \{s \mid (\exists s'. (s, s') \notin r \vee s' \in Q_1) \vee (\exists s'. (s, s') \notin r \vee s' \in Q_2)\} \\ &\Leftrightarrow wp(r, Q_1) \cup wp(r, Q_2)\end{aligned}$$

- (vii) counterexample:  $S = \{a, b\}$  and  $r_1 = \{(a, b)\}$  and  $r_2 = \{(b, a)\}$  and  $Q = \{b\}$

# Exercise Solution

(viii)  $S = \mathbb{N}\{0\}$

$r = \{(s, s') \mid \exists c. s + c = s'\}$ , i.e. (1, 2), (1, 3), (1, 4), ..., (2, 3), (2, 4), ...

$S \neq \emptyset \wedge \text{dom}(r) = S \wedge \Delta_S \cap r = \emptyset$  is satisfied

$r \circ r \cap ((S \times S) \setminus r) \neq \emptyset$  is equivalent to  $r \circ r \subseteq r$ , i.e. the relation is transitive, which is the case for our  $r$ .

$S$  cannot be finite, since  $r$  is closed under composition, i.e. it is a transitive closure. If  $S$  was finite, we'd necessarily have a loop somewhere violating the  $\Delta_S \cap r = \emptyset$  requirement.