

Omega Test

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Decision Procedures

Decision procedures for Presburger Arithmetic

- Quantifier Elimination
- Cooper's Algorithm
- Automata based
- *Omega Test*
- ...

Description

- Satisfiability of the conjunction of two type of constraints

$$\sum_{i=1}^n a_i b_i = b$$

- Variable elimination using a method based on Euclid's algorithm

$$\sum_{i=1}^n a_i b_i \leq b$$

- Variable elimination using the Fourier-Motzkin method

Euclid's Algorithm

$$\text{GCD}(a, b) = \begin{cases} a & \text{if } b = 0 \\ \text{GCD}(b, a \bmod b) & \text{otherwise} \end{cases}$$

Example

$$\text{GCD}(22,6) = \text{GCD}(6,4) = \text{GCD}(4,2) = \text{GCD}(2,0) = 2$$

Least Remainder Algorithm

$$\text{GCD}(a, b) = \begin{cases} a & \text{if } b = 0 \\ \text{GCD}(b, \widehat{a \bmod b}) & \text{otherwise} \end{cases}$$

- Idea: Instead of replacing b by $(a \bmod b)$ during the division step, replace it by $(a \bmod b) - b$ if $|a \bmod b| > \frac{1}{2}b$

$$\widehat{a \bmod b} \stackrel{\text{def}}{=} a - b \lfloor \frac{a}{b} + \frac{1}{2} \rfloor$$

Example

$$\text{GCD}(22, 6) = \text{GCD}(6, -2) = \text{GCD}(-2, 0) = -2$$

Equality Constraints

We want to eliminate the variable x_1 in the system using the following equation

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n = b$$

- 1 if $a_1 = 1$: move all the terms other than x_1 to the right hand side
- 2 if $a_1 > 1$:
 - ▶ if a_1 divides $a_2 \cdots a_n$ then there is a solution only if it divides also b
 - ★ Divide both sides by a_1 and compute as the first case
 - ▶ if a_1 fails to divide $a_2 \cdots a_n$
 - ★ Introduce a new variable x'_1 by

$$\lfloor a_1/a_1 \rfloor x_1 + \lfloor a_2/a_1 \rfloor x_2 + \cdots + \lfloor a_n/a_1 \rfloor x_n = x'_1$$

- ★ Solve the equation for x_1 and replace the original equation with the following

$$a_1x'_1 + (a_2 \bmod a_1)x_2 + \cdots + (a_n \bmod a_1)x_n = b$$

Equality Constraints

Exercise

Remove the equality constraints from the following system

$$\begin{cases} x_1 + 3x_2 + 5x_3 = 9 \\ 2x_1 + 3x_2 = 6 \\ x_1 + x_3 \leq 7 \end{cases}$$

The Omega Test method uses the least remainder algorithm instead of the Euclid's algorithm

Fourier-Motzkin Elimination

$$(\exists x. ax \leq \alpha \wedge \beta \leq bx) \equiv a\beta \leq b\alpha$$

$$\left\{ \begin{array}{l} ax \leq \alpha \\ \beta \leq bx \end{array} \right\} \begin{array}{c} \updownarrow \\ (a, b > 0) \end{array} a\beta \leq b\alpha$$

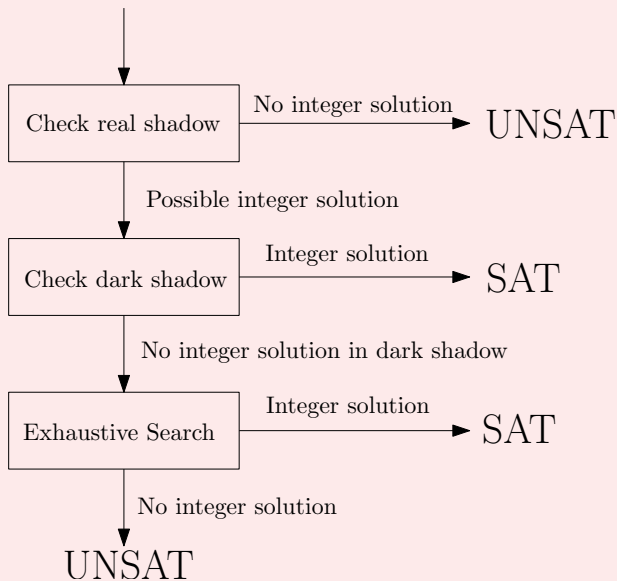
Does the equivalence hold for the integers?

$$(\exists x \in \mathbb{Z}. 3 \leq 2x \wedge 2x \leq 3) \not\equiv 6 \leq 6$$

Exercise

A sufficient (but not necessary) condition for existence of integer x in the equivalence is the following $b\alpha - a\beta \geq (a-1)(b-1)$

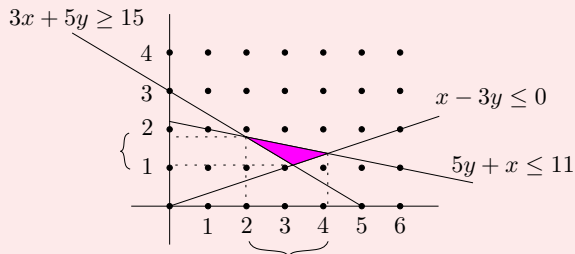
Inequality Constraints



Real Shadow

Check if there are integer solutions in the real solution space

$$\begin{cases} 3x + 5y \geq 15 \\ x - 3y \leq 0 \\ 5y + x \leq 11 \end{cases}$$



Dark Shadow

The subset of the real shadow that further satisfies the integer constraint

$$(b\alpha - a\beta \geq (a - 1)(b - 1)) \Rightarrow (\exists x. ax \leq \alpha \wedge \beta \leq bx)$$

This can eliminate the false positive from the previous example

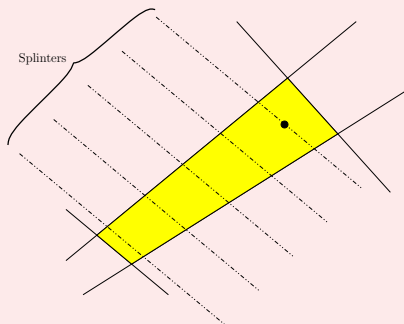
$$\begin{cases} 5y & \leq 11 - x \\ 15 - 3x & \leq 5y \end{cases}$$

3 does not satisfy the following constraint:

$$5(11 - x) - 5(15 - 3x) \geq (4)(4)$$

Omega Nightmare

- No integer solution in the dark shadow is not enough
- If the real shadow is non-empty but the dark shadow is empty an exhaustive search is done ($\exists x. ax \leq \alpha \wedge \beta \leq bx$)
- Try all the integers i such that $0 \leq i \leq \frac{ab-a-b}{a}$
- See if any of the equations of the form $bx = \beta + i$ has an integer solution



References

- “Decision procedures: an algorithmic point of view”; Daniel Kroening, Ofer Strichman, EATCS Springer - 2008
- “Dependence Analysis”; Utpal K. Banerjee, Kluwer Academic Publishers - 1996