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# Parallelism and Concurrency

## Midterm Exam

Wednesday, April 18, 2018

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Your points are *precious*, don't let them go to waste!

**Your Time** All points are not equal. Note that we do not think that all exercises have the same difficulty, even if they have the same number of points.

**Your Attention** The exam problems are precisely and carefully formulated, some details can be subtle. Pay attention, because if you do not understand a problem, you can not obtain full points.

| Exercise     | Points | Points Achieved |
|--------------|--------|-----------------|
| 1            | 20     |                 |
| 2            | 20     |                 |
| 3            | 20     |                 |
| <b>Total</b> | 60     |                 |

## Exercise 1: Parallel Reduce-By-Key (20 points)

In this exercise you are tasked to implement a version of the reduce-by-key operation using `ParSeq.aggregate`. The function takes a sequence `data` of tuples consisting of keys and values, and a function `h` that *reduces* two values into one. As its output, the function returns the mapping from keys to the reduced subsequences of values in `data` which have their key in common.

### Example

Given the following input sequence of tuples

```
(1, "do"), (2, "re"), (1, "mi"), (4, "fa"), (4, "so"), (2, "la"), (1, "ti")
```

and the reduction function `+` (i.e., string concatenation) we want to produce the mapping:

```
1 -> "domiti"  
2 -> "rela"  
4 -> "faso"
```

### Question 1

Your task is to implement the reduce-by-key operation in a way that is *generic*, *efficient* and *parallel*. The operation takes as its inputs a `ParSeq` of  $(K, V)$  tuples, where the  $K$  component is the key, along with a reduction function `h` that takes two values of type  $V$  and returns a reduced  $V$ . Your implementation should reduce all the values in `data` that have the same key in the order that they occur in `data`, and produce a `Map` going from the keys to the reduced values.

**Your method should be implemented in terms of the `ParSeq.aggregate` method.** Remember that the `aggregate` method has certain requirements on its parameters, i.e., `z` and the two functions `f` and `g`.

**Hint:** You can refer to the appendix for an API of `Map` and other useful classes and methods.

Please answer on the next page.

```
def reduceByKey[K, V](data: ParSeq[(K, V)])(h: (V, V) => V): Map[K, V] =
```

## Question 2

In general, which of the following properties should hold for the parameters  $z$ ,  $f$  and  $g$  of `ParSeq.aggregate[S]`?

1.  $f(g(u, v), x) == g(f(u, x), v)$
2.  $g(u, z) == z$
3.  $g(u, z) == u$
4.  $f(f(z, x), x) == f(g(f(z, x), z), x)$
5.  $(xs ++ ys).foldLeft(z)(f) == g(xs.foldLeft(z)(f), ys.foldLeft(z)(f))$
6.  $xs.foldLeft(z)(f) == xs.foldRight(z) \{ (x, u) => f(u, x) \}$

where  $u, v: S, x: T$  and  $xs, ys: ParSeq[T]$  are arbitrary.

**Please circle the correct answer(s).**



## Exercise 2: Concurrency (20 points)

In concurrent systems, there is a common category of resources that require the system to satisfy the following two properties:

- **Property 1:** The resource can be read in parallel by multiple readers, however,
- **Property 2:** Writing to the resource requires exclusive access.

In another word, read-read may be overlapped, but read-write and write-write must be mutually exclusive.

One solution to implement the two properties is to use a read lock `rl`, a write lock `wl`, and a number `rn` to remember the number of concurrent readers. A tentative algorithm is given below:

```
val rl = new Lock // read lock
val wl = new Lock // write lock
var rn = 0       // readers count

def read(): Int = {
  rl.acquire()
  rn = rn + 1
  if (rn == 1) wl.acquire()

  // read resource, compute result, omitted

  rn = rn - 1
  if (rn == 0) wl.release()
  rl.release()

  result
}

def write(param: Int): Unit = {
  wl.acquire()
  rl.acquire()

  // write resource using param, omitted

  rl.release()
  wl.release()
}
```

### Question 1

Does the algorithm above satisfy the two properties? Briefly explain.

## Question 2

Is the code vulnerable to deadlocks ? If it is the case, describe an execution schedule that leads to a deadlock. Otherwise, motivate your answer.

## Question 3

Improve `read` and `write` so that the two properties are satisfied and there is no deadlock. You may introduce extra variables.

### Exercise 3: Parallel String Match (20 points)

Your task in this exercise is to count the number of appearance of a given word in a sentence. The function to implement has the following signature:

```
def count(word: Array[Char], sentence: Array[Char]): Int
```

Here are two examples where count returns 2:

```
count("you".toArray, "you know what you're doing".toArray) == 2
//                ^^^ 1st one  ^^^ 2nd one

count("lala".toArray, "lalala".toArray) == 2
//                ^^^^^ 1st one
//                ^^^^^ 2nd one
```

Instead of implementing `count` directly, you will have to implement two helper methods, `countWithin` and `countWithinPar`, which in addition to the two input arrays take a start index and an end index to define the slice of the `sentence` where the count should be performed. Afterwards, `count` can simply be implemented as:

```
def count(word: Array[Char], sentence: Array[Char]): Int =
  countWithin(word, sentence, 0, sentence.length)
```

Or, for a parallel implementation:

```
def countPar(word: Array[Char], sentence: Array[Char]): Int =
  countWithinPar(word, sentence, 0, sentence.length)
```

Please see the next pages for the questions.

## Question 1

Implement a sequential version of `count` called `countWithin`, which takes as additional parameters the array indices at which to `start` and `end` the search within `sentence`. The index `start` is inclusive and `end` is exclusive. Your implementation should perform at most  $O(n^2)$  work, where  $n$  is the size of the of the array slice, i.e.,  $n = \text{end} - \text{start}$ .

```
def countWithin(word: Array[Char], sentence: Array[Char], start: Int, end: Int): Int =
```

## Question 2

Implement an efficient parallel version of `countWithin` called `countWithinPar` using the `parallel` method presented in class (and in the appendix). Your implementation should be sequential for chunks of `THRESHOLD` characters or less. You can assume that the length of the `word` will always be strictly shorter than the `THRESHOLD`. Also, remember that the index `start` is inclusive and `end` is exclusive.

For this question, you can also assume that you have a correct sequential implementation of `countWithin` from question 1. Make sure to use it!

```
def countWithinPar(word: Array[Char], sentence: Array[Char], start: Int, end: Int): Int =
```

### Question 3

What is the *depth* of `countPar` in terms of Big-O notation, assuming unbounded parallelism and constant size words and `THRESHOLD`? Briefly explain.

Remember that depth is defined as the length of the longest series of operations that have to be performed sequentially.



## Appendix

The provided API is not meant to be exhaustive. You can, of course, use methods and classes that are not presented in this API.

### Parallel API

- `def parallel[A, B](op1: => A, op2: => B): (A, B)`: Executes the two given computations in parallel and returns the pair of results.

### Array API

Relevant API for `Array[T]`:

- `def apply(index: Int): T`: Gets the element at a given index. Indexes start at 0.
- `def length: Int`: Returns the length of the array.

### ParSeq API

Relevant API for `ParSeq[A]`. You can assume that the implementations of the following methods are similar to what was seen during the lectures.

- `def map[B](f: (A) => B): ParSeq[B]`: Applies, in parallel, a function on all elements of the sequence, and returns the sequence of results.
- `def reduce(op: (A, A) => A): A`: Reduces the elements of this sequence using the specified operator.
- `def aggregate[B](z: => B)(f: (B, A) => B, g: (B, B) => B): B`: Aggregates the results of applying an operator to the elements.
- `def foldLeft[S](z: S)(f: (S, T) => S): S`: Folds the elements of this sequence sequentially from left to right using the specified operator `f`.
- `def foldRight[S](z: S)(f: (T, S) => S): S`: Folds the elements of this sequence sequentially from right to left using the specified operator `f`.

### Map API

Relevant API for `Map[K, V]`. A `Map[K, V]` is an immutable collection mapping keys `K` to values `V` and supports the following operations efficiently:

- `def ++(other: Map[K, V]): Map[K, V]`: Concatenates two mappings. The second mappings has priority over the first in the case of duplicate keys.
- `def +(elem: (K, V)): Map[K, V]`: Appends an entry in the mappings. If the key previously existed, it is overridden.

- `def contains(key: K): Boolean`: Checks whether a given key exists in this mapping.
- `def apply(key: K): V`: Retrieve the value associated to the key if it exists, otherwise throws an exception.
- `def get(key: K): Option[V]`: Retrieve the value associated to the key, wrapped in `Some`, if it exists, otherwise returns `None`.
- `def keys: Seq[K]`: Returns all the map's keys in some unspecified order.

To create an empty `Map`, you can use the following method:

- `Map.empty[K, V]: Map[K, V]`: Creates an empty map.

## Seq API

Relevant minimal API for `Seq[A]`.

- `def map[B](f: (A) => B): ParSeq[B]`: Applies, in parallel, a function on all elements of the sequence, and returns the sequence of results.
- `def foldLeft[S](z: S)(f: (S, T) => S): S`: Folds the elements of this sequence sequentially from left to right using the specified operator `f`.
- `def foldRight[S](z: S)(f: (T, S) => S): S`: Folds the elements of this sequence sequentially from right to left using the specified operator `f`.