



# Using partial-evaluation to efficiently implement a language in GraalVM

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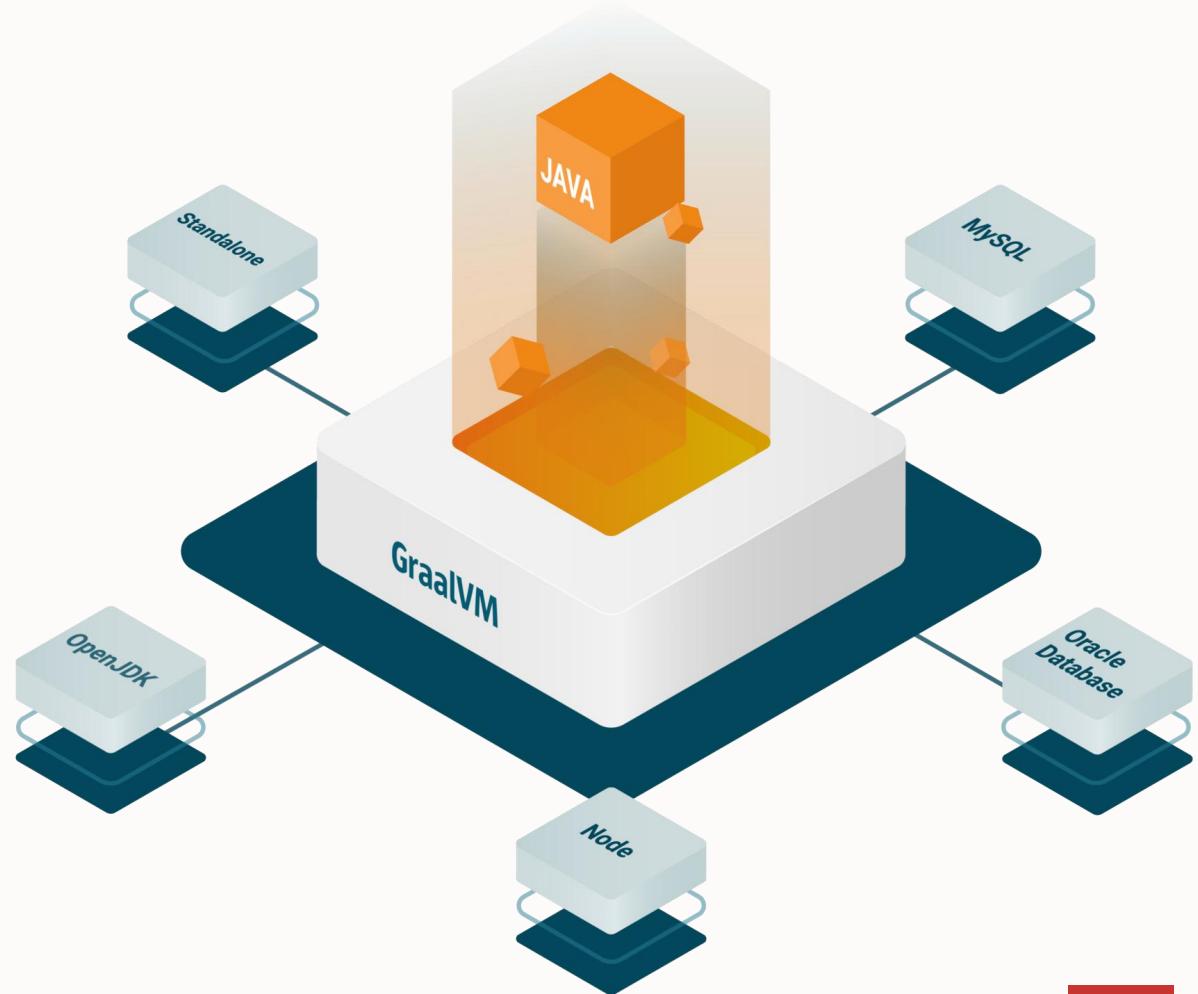
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# Implementing a language in GraalVM

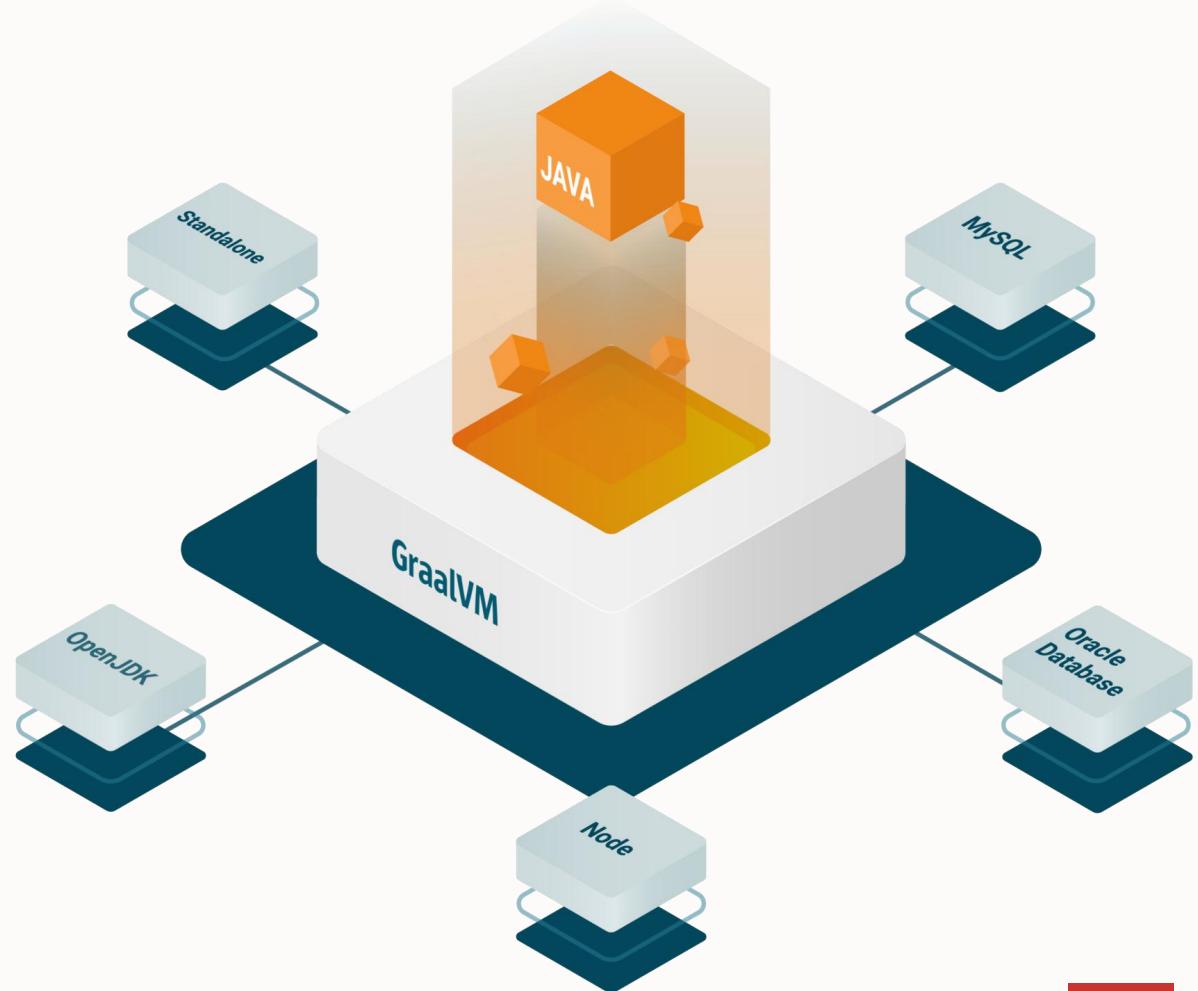
GraalVM is a shiny high-performance next-generation multi-lingual embeddable VM.



# Implementing a language in GraalVM

GraalVM is a shiny high-performance next-generation multi-lingual embeddable VM.

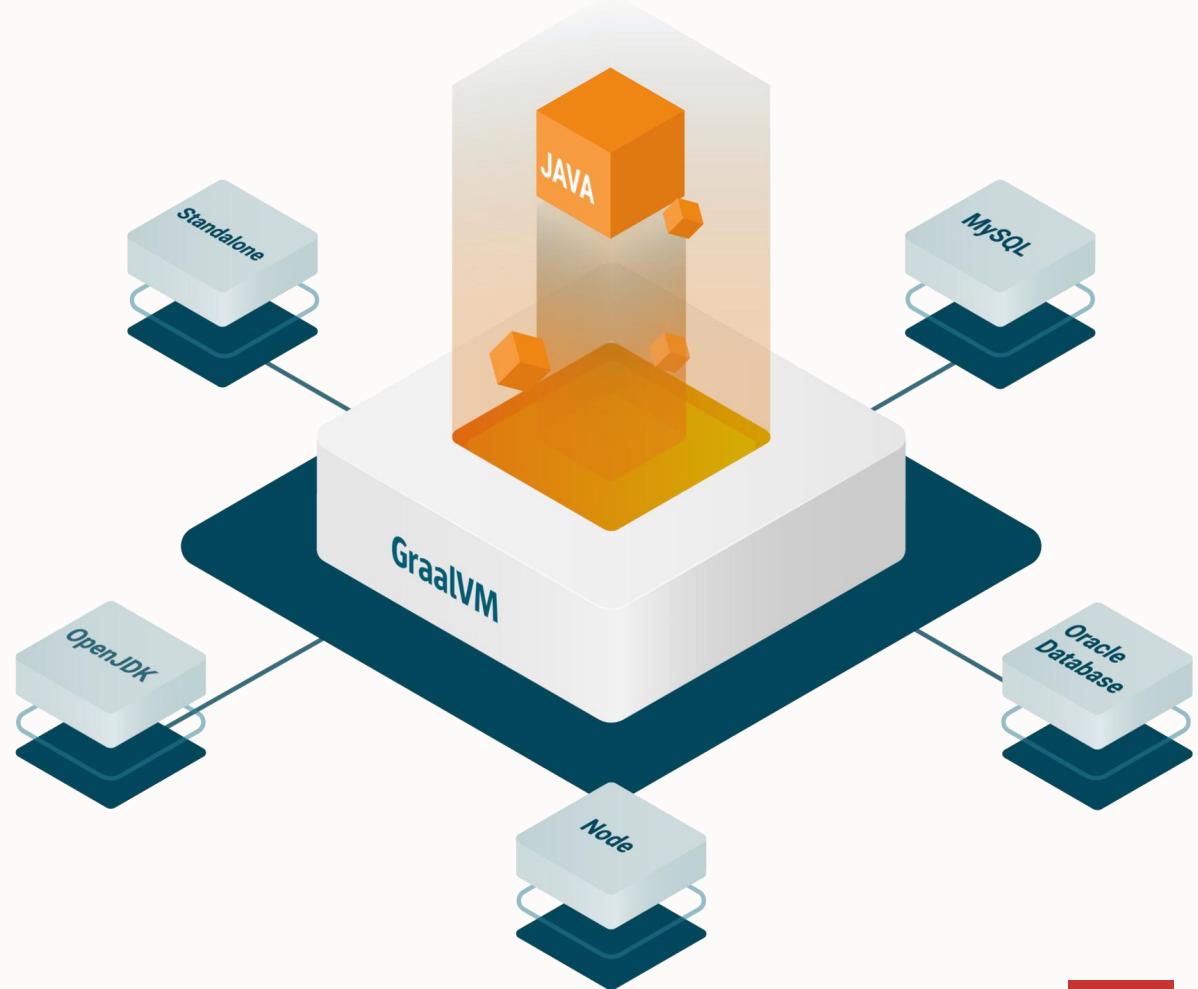
How do you implement a new language for it?



# Implementing a language in GraalVM

## Steps to create a language

1. Implement an executable AST.
2. Implement a parser that creates the AST for an input program.



# Running example

Expression-stack languages with arithmetic and blocks.

# Expression-stack language with arithmetic and blocks

---

```
<instr> =  
    i32.const <n>  
    i32.add  
    i32.eq  
    block  
        <instr>*  
    end
```

# Expression-stack language with arithmetic and blocks

```
<instr> =  
    i32.const <n>  
    i32.add  
    i32.eq  
    block  
        <instr>*  
    end
```

Grammar

```
block  
    i32.const 7  
    i32.const 9  
    i32.add  
    block  
        i32.const 11  
        i32.const 5  
        i32.add  
    end  
    i32.eq  
end
```

Example program

# Expression-stack language with arithmetic and blocks

Let's show the semantics  
on an example.

```
block
i32.const 7
i32.const 9
i32.add
block
i32.const 11
i32.const 5
i32.add
end
i32.eq
end
```

Program state

Expression stack and program counter

Example program

# Expression-stack language with arithmetic and blocks

7

Program state

block  
→ i32.const 7  
i32.const 9  
i32.add  
block  
i32.const 11  
i32.const 5  
i32.add  
end  
i32.eq  
end

Example program

# Expression-stack language with arithmetic and blocks



Program state

block  
i32.const 7  
i32.const 9  
i32.add  
block  
i32.const 11  
i32.const 5  
i32.add  
end  
i32.eq  
end

A red arrow points from the word "block" in the "Program state" section to the first "block" keyword in the "Example program" code block.

Example program

# Expression-stack language with arithmetic and blocks

16

Program state

```
block
i32.const 7
i32.const 9
i32.add
block
i32.const 11
i32.const 5
i32.add
end
i32.eq
end
```



Example program

# Expression-stack language with arithmetic and blocks

16

Program state

```
block
i32.const 7
i32.const 9
i32.add
block
i32.const 11
i32.const 5
i32.add
end
i32.eq
end
```



Example program

# Expression-stack language with arithmetic and blocks

5  
11  
16

Program state

block  
i32.const 7  
i32.const 9  
i32.add  
block  
i32.const 11  
i32.const 5  
i32.add  
end  
i32.eq  
end



Example program

# Expression-stack language with arithmetic and blocks

16  
16

Program state

block  
i32.const 7  
i32.const 9  
i32.add  
block  
i32.const 11  
i32.const 5  
i32.add  
end  
i32.eq  
end



Example program

# Expression-stack language with arithmetic and blocks

1

Program state

```
block
i32.const 7
i32.const 9
i32.add
block
i32.const 11
i32.const 5
i32.add
end
i32.eq
end
```



Example program

# Expression-stack language with arithmetic and blocks

Let's capture the semantics  
in an interpreter.

1

Program state

```
block
i32.const 7
i32.const 9
i32.add
block
i32.const 11
i32.const 5
i32.add
end
i32.eq
end
```



Example program

# Implementing a language with an interpreter loop

## Step 1: extend the Node class

```
class BlockNode extends Node {  
    final int offset;  
    final byte[] program;  
  
    BlockNode(int offset, byte[] program) {  
        this.offset = offset;  
        this.program = program;  
    }  
  
    Object execute(Frame frame) { ... }  
}
```

```
0: block  
1: i32.const 7  
2: i32.const 9  
3: i32.add  
4: block  
5: i32.const 11  
6: i32.const 5  
7: i32.add  
8: end  
9: i32.eq  
10: end
```

# Implementing a language with an interpreter loop

## Step 2: implement the execute method

```
Object execute(Frame frame) {
```

Think of it as an array of  
local variables

```
}
```

```
0: block
1: i32.const 7
2: i32.const 9
3: i32.add
4: block
5: i32.const 11
6: i32.const 5
7: i32.add
8: end
9: i32.eq
10: end
```

# Implementing a language with an interpreter loop

## Step 2: implement the execute method

```
Object execute(Frame frame) {  
    int pc = offset;  
    while (program[pc] != END) {  
        }  
        return null;  
    }
```

```
0: block  
1: i32.const 7  
2: i32.const 9  
3: i32.add  
4: block  
5: i32.const 11  
6: i32.const 5  
7: i32.add  
8: end  
9: i32.eq  
10: end
```

# Implementing a language with an interpreter loop

## Step 2: implement the execute method

```
Object execute(Frame frame) {  
    int pc = offset;  
    while (program[pc] != END) {  
        switch (program[pc++]) {  
            }  
        }  
        return null;  
    }
```

0 : block	
1 : i32.const 7	
2 : i32.const 9	
3 : i32.add	
4 : block	
5 : i32.const 11	
6 : i32.const 5	
7 : i32.add	
8 : end	
9 : i32.eq	
10 : end	

# Implementing a language with an interpreter loop

## Step 2: implement the execute method

```
Object execute(Frame frame) {
    int pc = offset;
    while (program[pc] != END) {
        switch (program[pc++]) {
            case I32_CONST: ...
            case I32_ADD: ...
            case I32_EQ: ...
            case BLOCK: ...
        }
    }
    return null;
}
```

0: block	
1: i32.const 7	
2: i32.const 9	
3: i32.add	
4: block	
5: i32.const 11	
6: i32.const 5	
7: i32.add	
8: end	
9: i32.eq	
10: end	

# Implementing a language with an interpreter loop

## Step 2: implement the execute method

```
Object execute(Frame frame) {
    int pc = offset;
    while (program[pc] != END) {
        switch (program[pc++]) {
            case I32_ADD:
                int x = pop(frame);
                int y = pop(frame);
                push(frame, x + y);
                break;
        }
    }
    return null;
}
```

0: block	
1: i32.const 7	
2: i32.const 9	
3: i32.add	
4: block	
5: i32.const 11	
6: i32.const 5	
7: i32.add	
8: end	
9: i32.eq	
10: end	

# Implementing a language with an interpreter loop

## Step 2: implement the execute method

```
Object execute(Frame frame) {
    int pc = offset;
    while (program[pc] != END) {
        switch (program[pc++]) {
            case BLOCK:
                BlockNode b = children[pc];
                b.execute(frame);
                break;

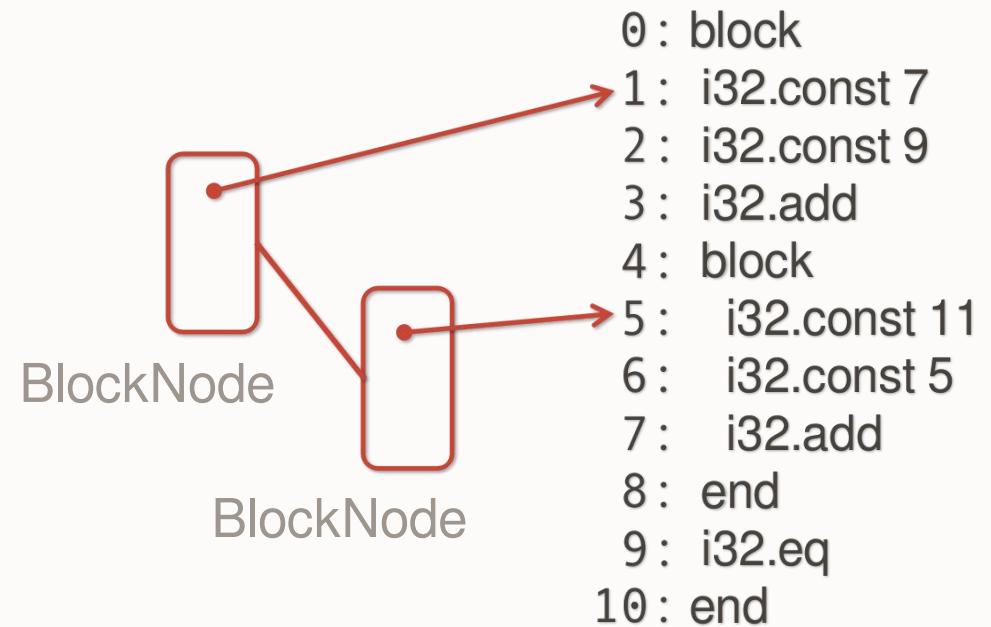
            }

        return null;
    }
}
```

0: block  
1: i32.const 7  
2: i32.const 9  
3: i32.add  
4: block  
5: i32.const 11  
6: i32.const 5  
7: i32.add  
8: end  
9: i32.eq  
10: end

# Implementing a language with an interpreter loop

AST holds offsets into the instructions array

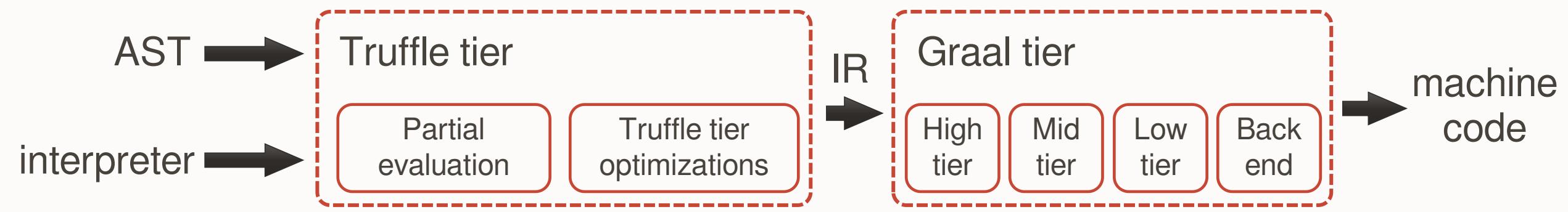


# Congratulations

You're now the local WebAssembly expert.

# How compilation works in Graal and Truffle

# Truffle compiler pipeline

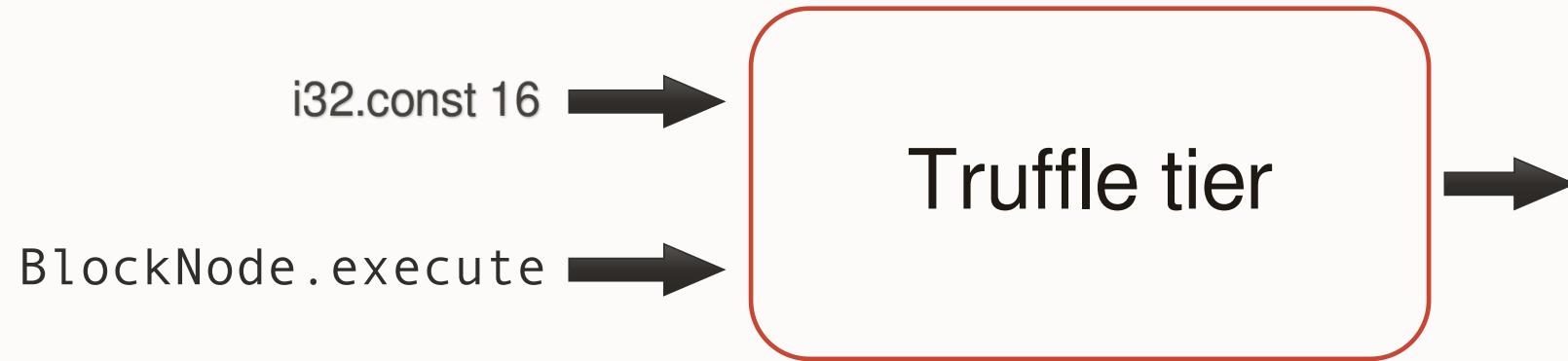


# Partial evaluation



Specializes the interpreter code to a given AST.

# Partial evaluation



# Partial evaluation



# Partial evaluation



# Partial evaluation



# Partial evaluation



# Partial evaluation



# Partial evaluation



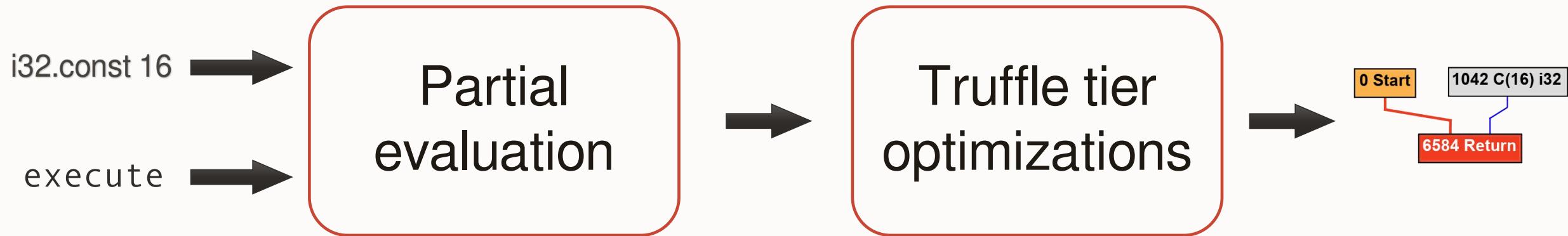
# Partial evaluation



# Partial evaluation



# Partial evaluation

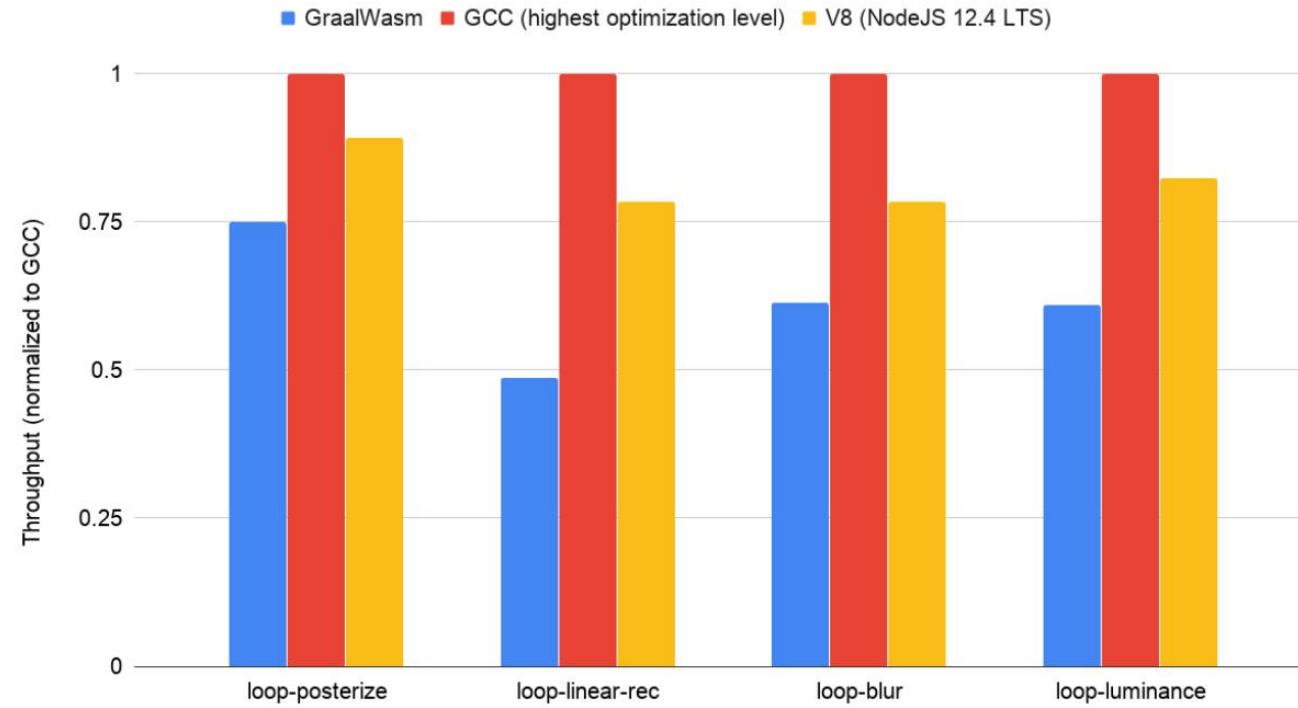


Most of the work is done in partial evaluation,  
and then optimizations simplify the IR.



# Stories in performance engineering

# Start by comparing with another VM



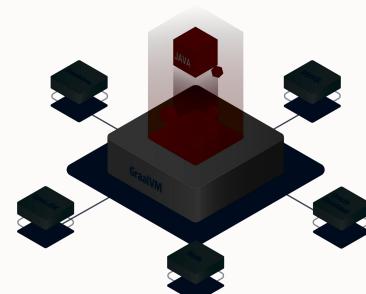
# How do I improve the performance?

# Look at the data

Understanding what's going on will help you  
solve the problem.

# Look at the data

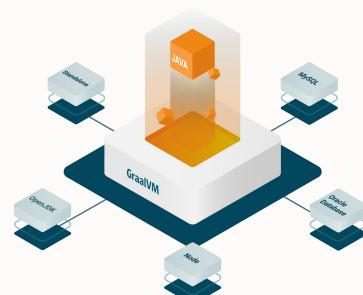
Understanding what's going on will help you solve the problem.



Your problem is not a black box.

# Look at the data

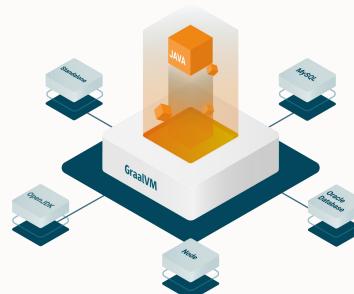
Understanding what's going on will help you solve the problem.



Your problem is not a  
black box.

# Look at the compiler IR

Understanding what the compiler works with  
will help you solve the problem.



# Memory out-of-bounds access profiling

---

Example program: invert filter



# Memory out-of-bounds access profiling

## Example program: invert filter

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = 0xFF - ((color & 0xFF000000) >> 24);  
    uint8_t G = 0xFF - ((color & 0x00FF0000) >> 16);  
    uint8_t B = 0xFF - ((color & 0x0000FF00) >> 8);  
    result[pixel % IMAGE_SIZE] = (R << 24) + (G << 16) + (B << 8) + 0xFF;  
}
```

C program

# Memory out-of-bounds access profiling

Example program: invert filter



# Memory out-of-bounds access profiling

## Example program: invert filter

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = 0xFF - ((color & 0xFF000000) >> 24);  
    uint8_t G = 0xFF - ((color & 0x00FF0000) >> 16);  
    uint8_t B = 0xFF - ((color & 0x0000FF00) >> 8);  
    result[pixel % IMAGE_SIZE] = (R << 24) + (G << 16) + (B << 8) + 0xFF;  
}
```

C program

# Memory out-of-bounds access profiling

## Example program: invert filter

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE]; ← Memory read  
    uint8_t R = 0xFF - ((color & 0xFF000000) >> 24);  
    uint8_t G = 0xFF - ((color & 0x00FF0000) >> 16);  
    uint8_t B = 0xFF - ((color & 0x0000FF00) >> 8);  
    result[pixel % IMAGE_SIZE] = (R << 24) + (G << 16) + (B << 8) + 0xFF;  
}
```

Memory write

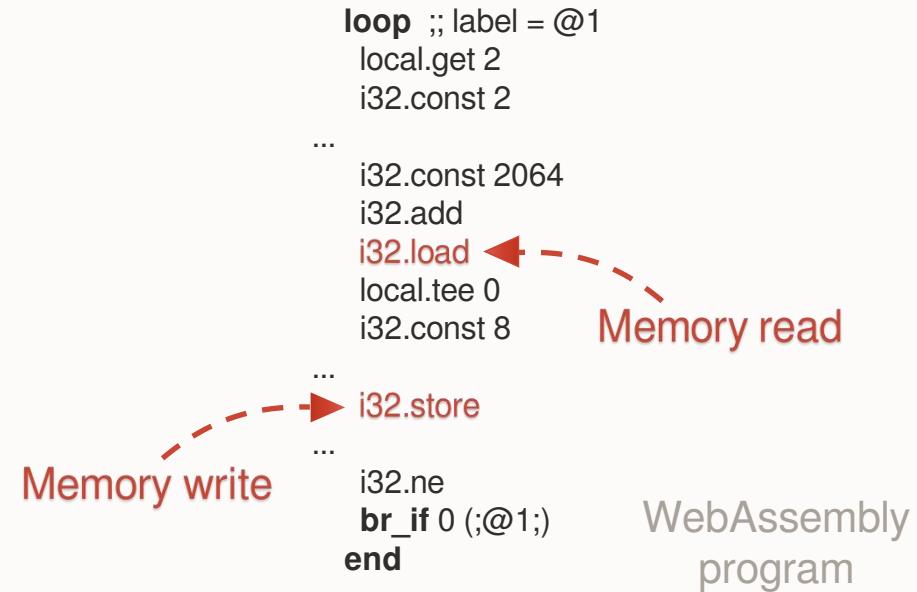
C program

# Memory out-of-bounds access profiling

## Example program: invert filter

```
loop ;, label = @1
  local.get 2
  i32.const 2
  ...
  i32.const 2064
  i32.add
  i32.load ← ----- Memory read
  local.tee 0
  i32.const 8
  ...
  i32.store → ----- Memory write
  ...
  i32.ne
  br_if 0 (;@1;)
end
```

WebAssembly  
program



# Memory out-of-bounds access profiling

---

## How to implement memory reads and writes in a Truffle interpreter?

1. Allocate a region of memory for the Truffle program.

```
public UnsafeWasmMemory(int pageCount) {  
    this.startAddress = unsafe.allocateMemory(pageCount * PAGE_SIZE);  
}
```

# Memory out-of-bounds access profiling

## How to implement memory reads and writes in a Truffle interpreter?

1. Allocate a region of memory for the Truffle program.
2. Implement methods that read and write to memory.

i32.load  
i32.store  
and a few more...  
WebAssembly

```
int load_i32(long address) {
    if (address < 0 || address + 4 > this.pageCount * PAGE_SIZE) {
        trapOutOfBounds(address, 4);
    }
    int value = unsafe.getInt(this.startAddress + address);
    return value;
}
```

GraalWasm

# Memory out-of-bounds access profiling

---

## How to implement memory reads and writes in a Truffle interpreter?

1. Allocate a region of memory for the Truffle program.
2. Implement methods that read and write to memory.
3. Call those methods when interpreting loads and stores.

```
case I32_LOAD: {
    int address = frame[--stackPointer];
    int value = memory.load_i32(address);
    frame[stackPointer] = value;
    stackPointer++;
    break;
}
```

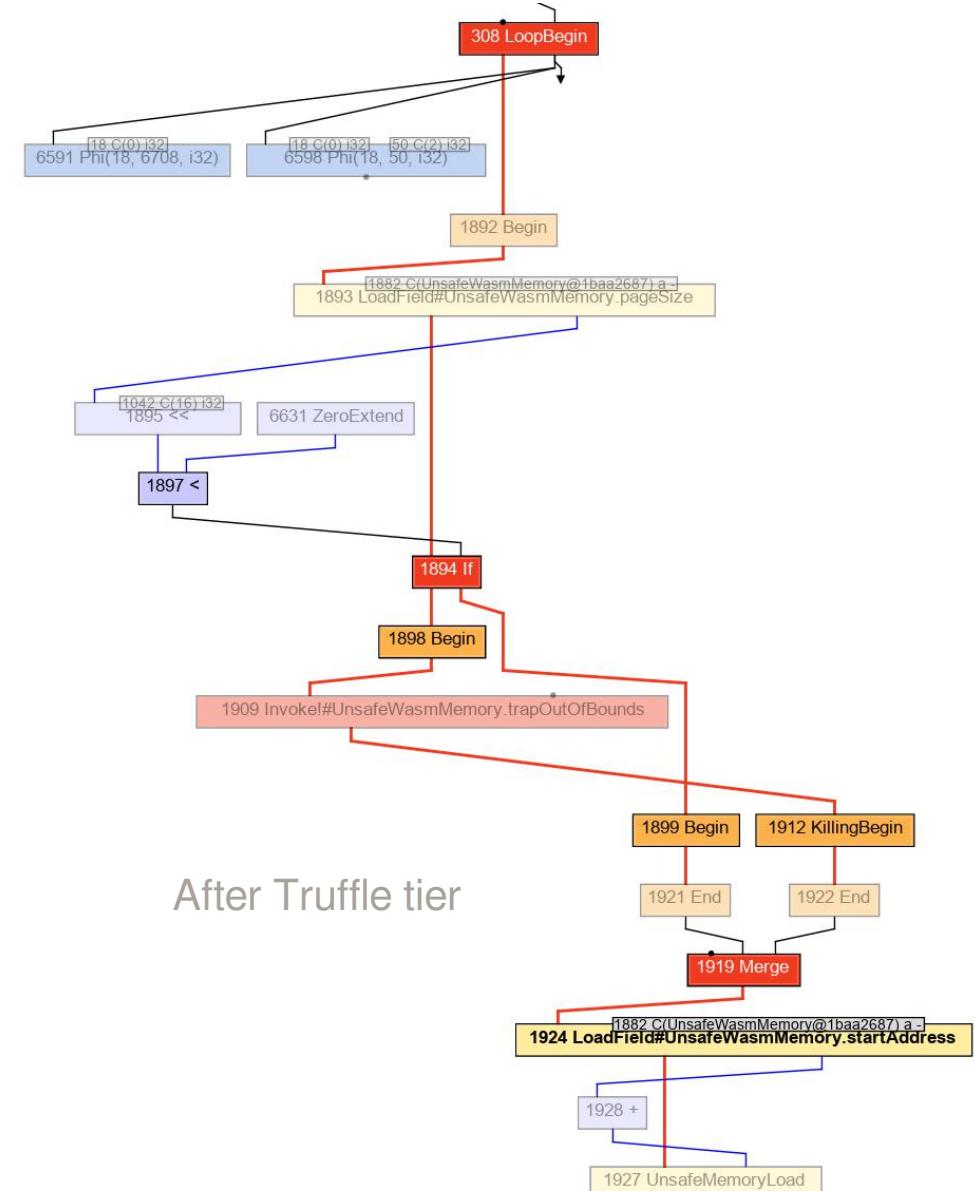
GraalWasm

# Memory out-of-bounds access profiling

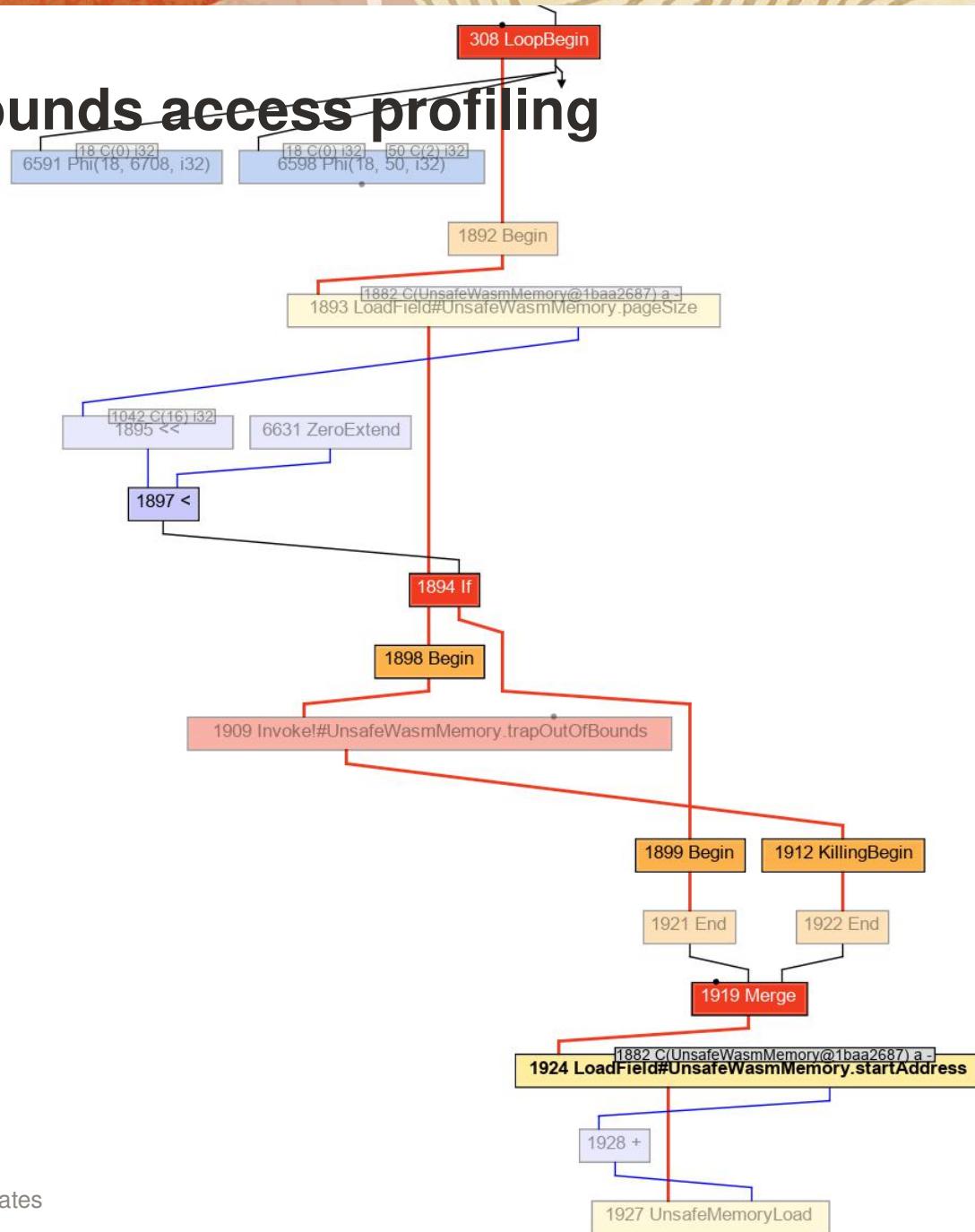
What does our program look like after partial evaluation completes?

```
case I32_LOAD: {
    int address = frame[--stackPointer];
    int value = memory.load_i32(address);
    frame[stackPointer] = value;
    stackPointer++;
    break;
}
```

GraalWasm



# Memory out-of-bounds access profiling

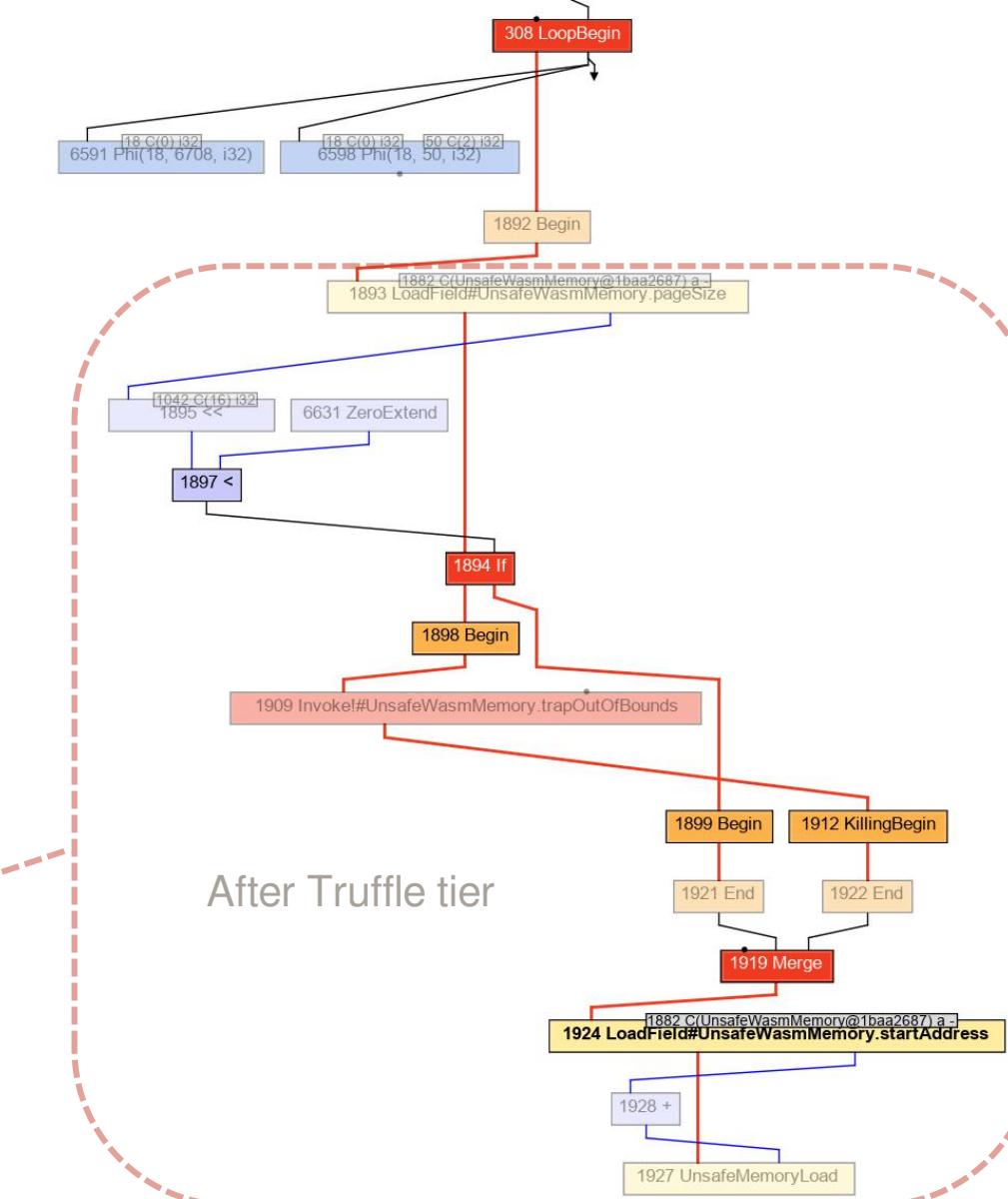


# Memory out-of-bounds access profiling

The part that corresponds to the bounds check and the memory read

```
case I32_LOAD: {
    int address = frame[--stackPointer];
    int value = memory.load_i32(address);
    frame[stackPointer] = value;
    stackPointer++;
    break;
}
```

GraalWasm

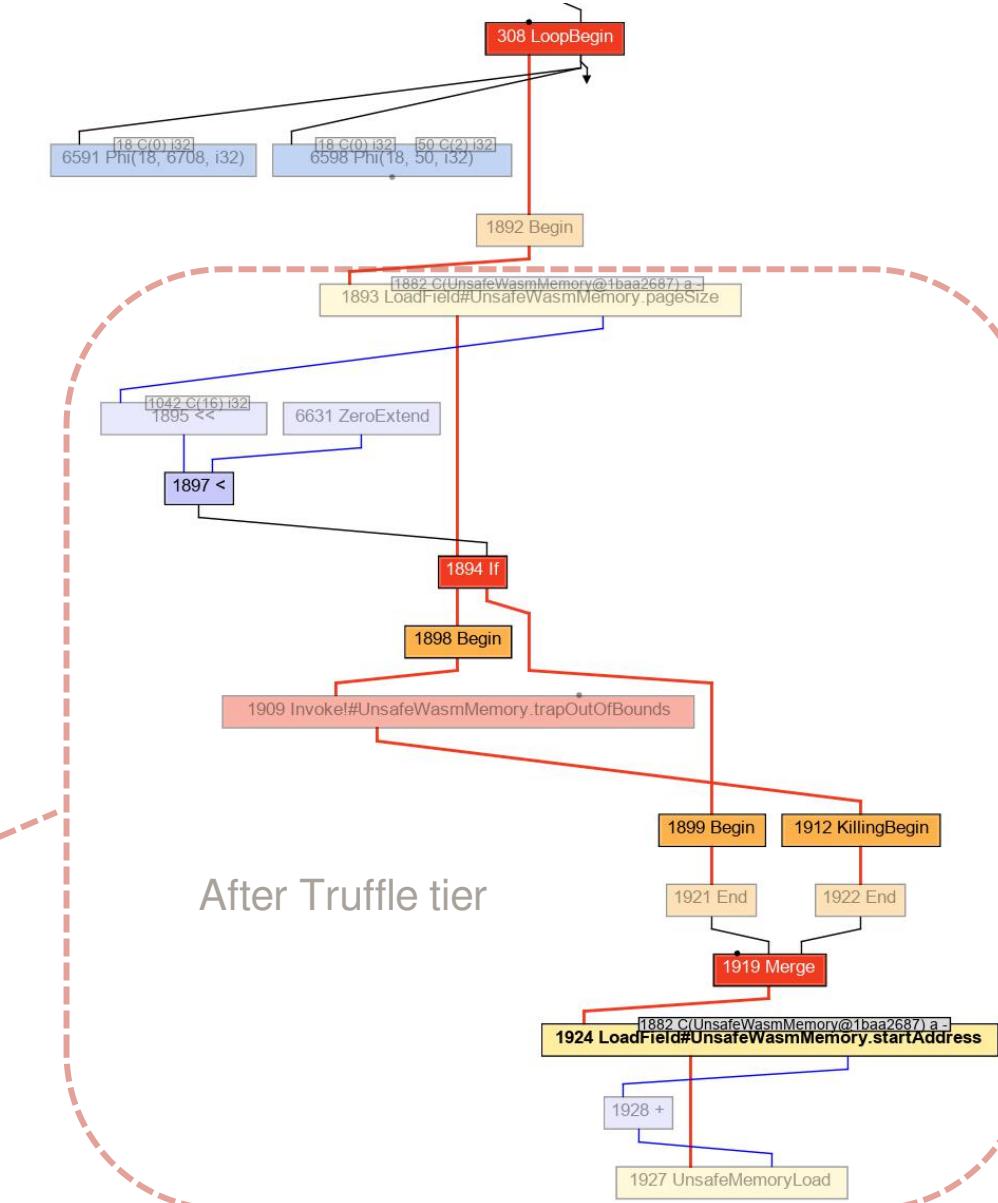


# Memory out-of-bounds access profiling

The part that corresponds to the bounds check and the memory read

```
if (address + 4 > pageCount * PAGE_SIZE) {  
    trapOutOfBounds(address, 4);  
}  
return unsafe.getInt(startAddress + address);
```

GraalWasm

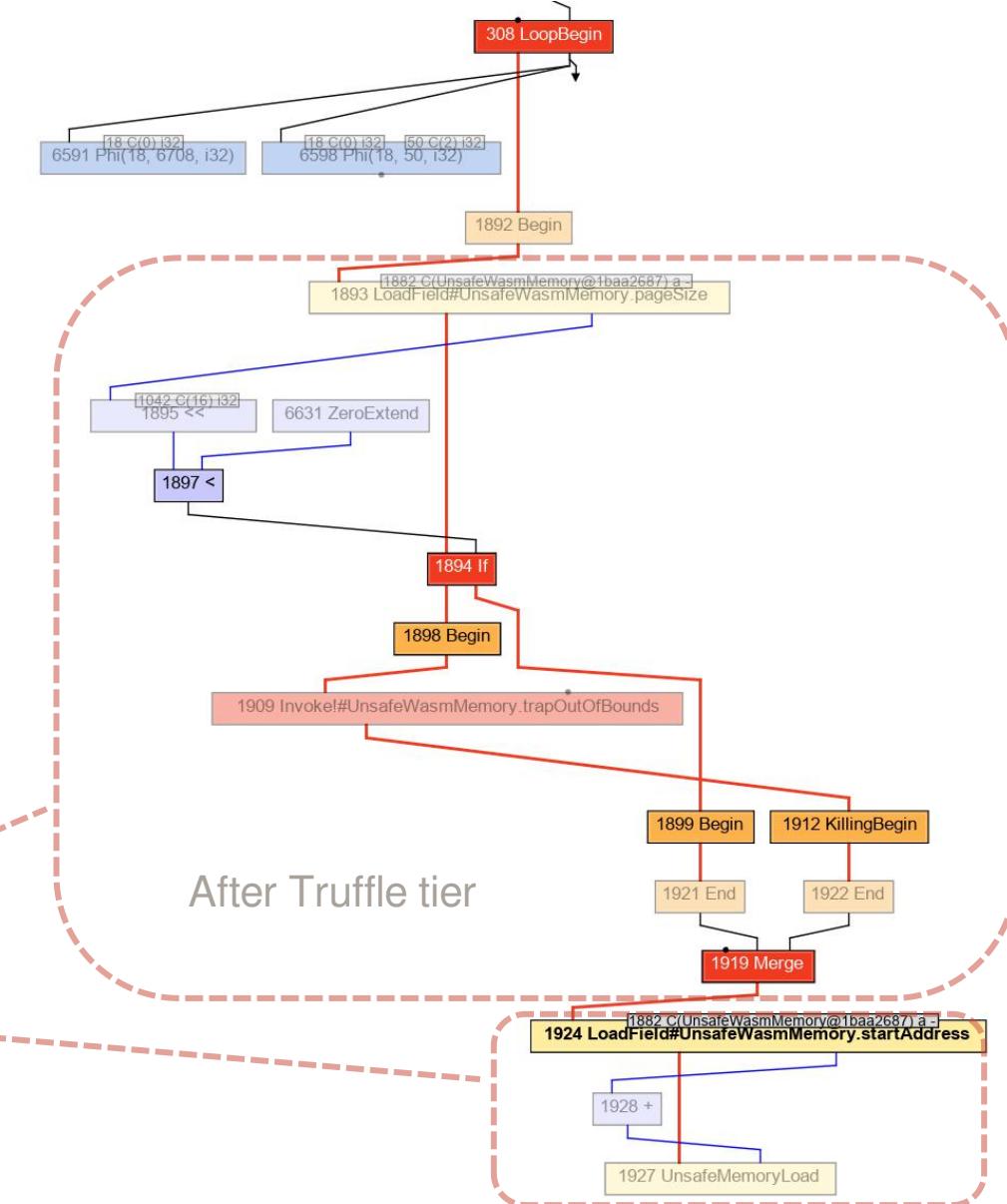


# Memory out-of-bounds access profiling

The part that corresponds to the bounds check and the memory read

```
if (address + 4 > pageCount * PAGE_SIZE) {  
    trapOutOfBounds(address, 4);  
}  
  
return unsafe.getInt(startAddress + address);
```

GraalWasm



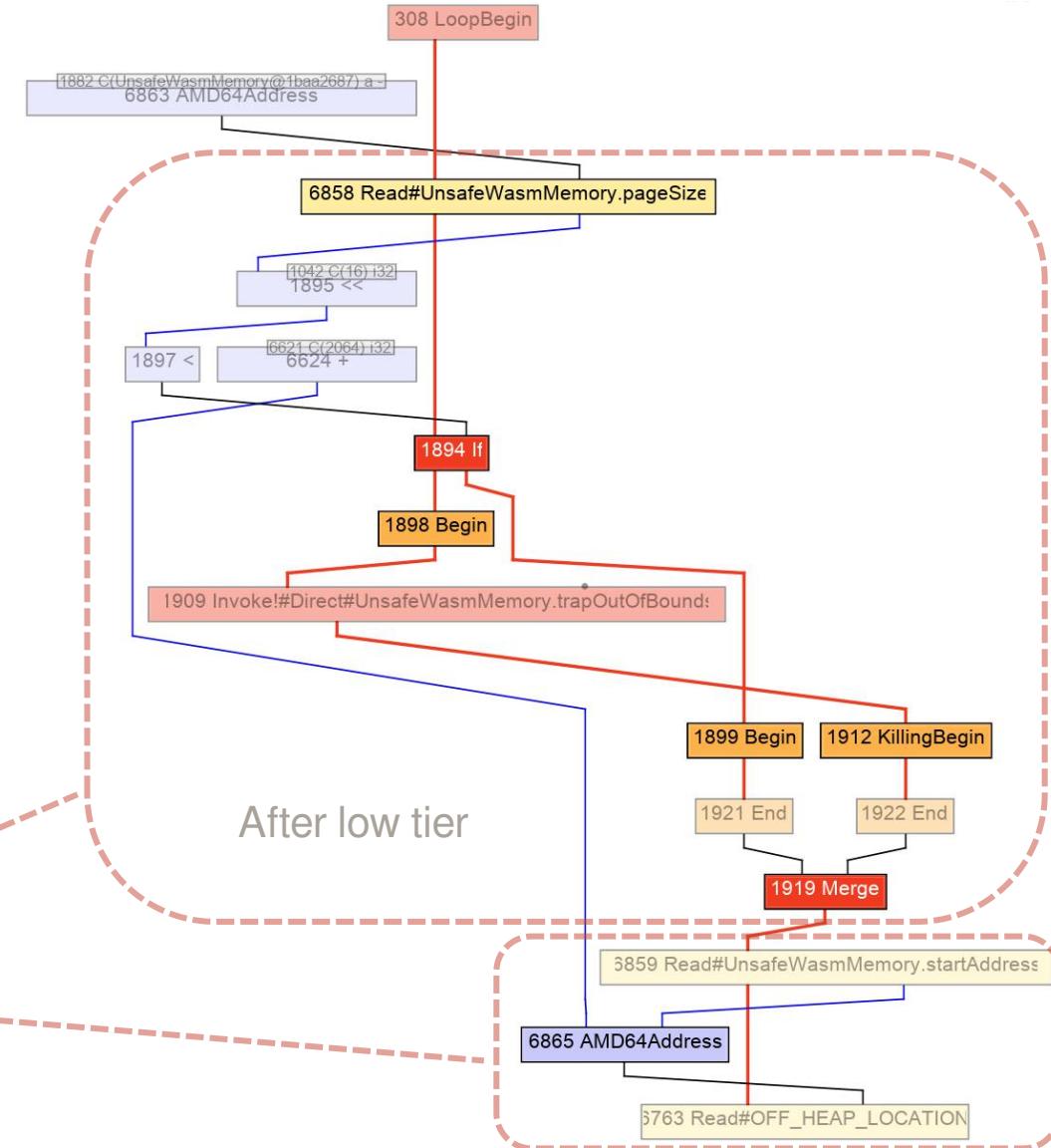
# Memory out-of-bounds access profiling

It is important to take a look at the IR after the low tier as well

LoadField nodes on Java objects are converted into Read nodes, which may have been moved around the IR.

```
if (address + 4 > pageCount * PAGE_SIZE) {  
    trapOutOfBounds(address, 4);  
}  
  
return unsafe.getInt(startAddress + address);
```

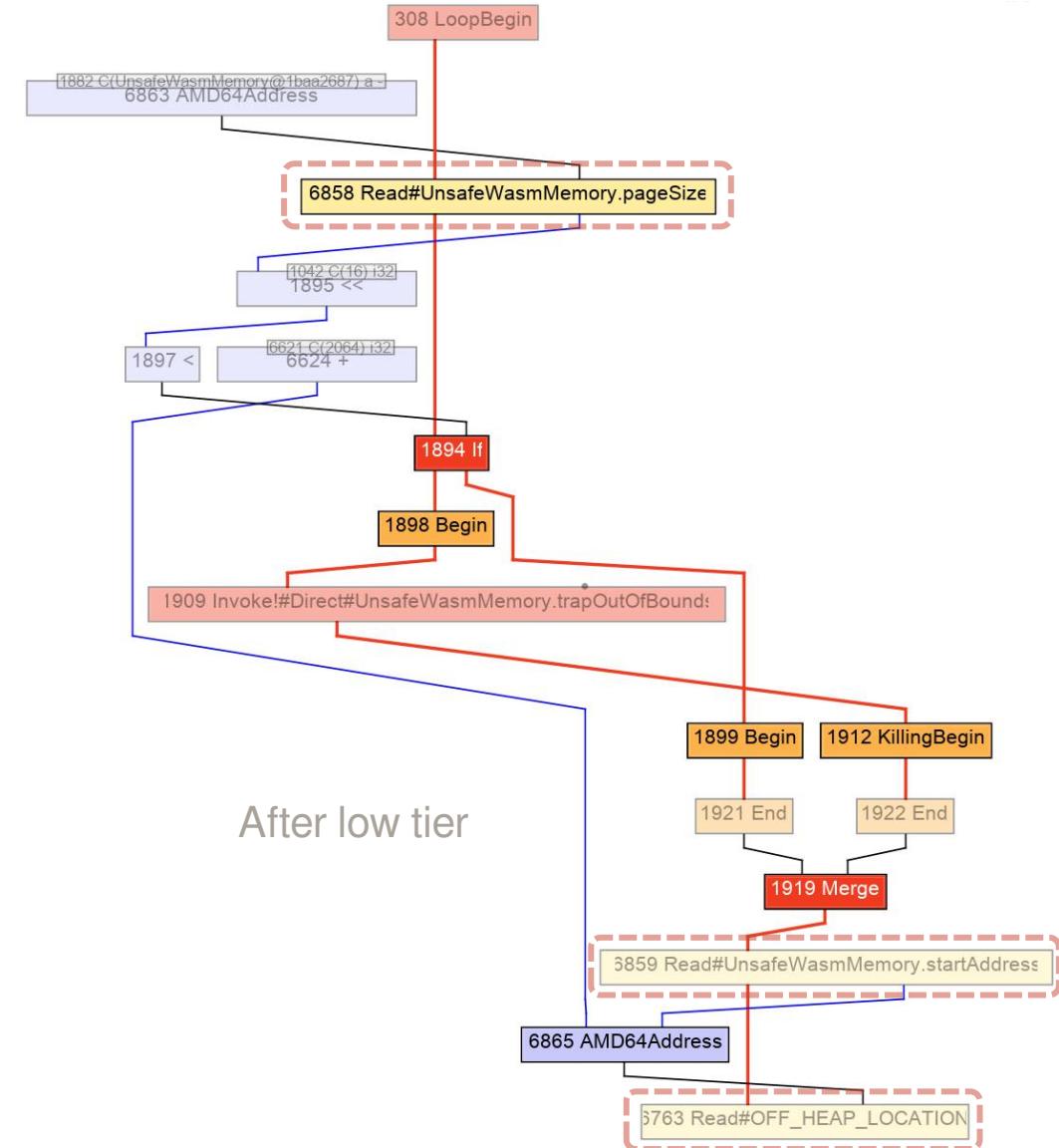
GraalWasm



# Memory out-of-bounds access profiling

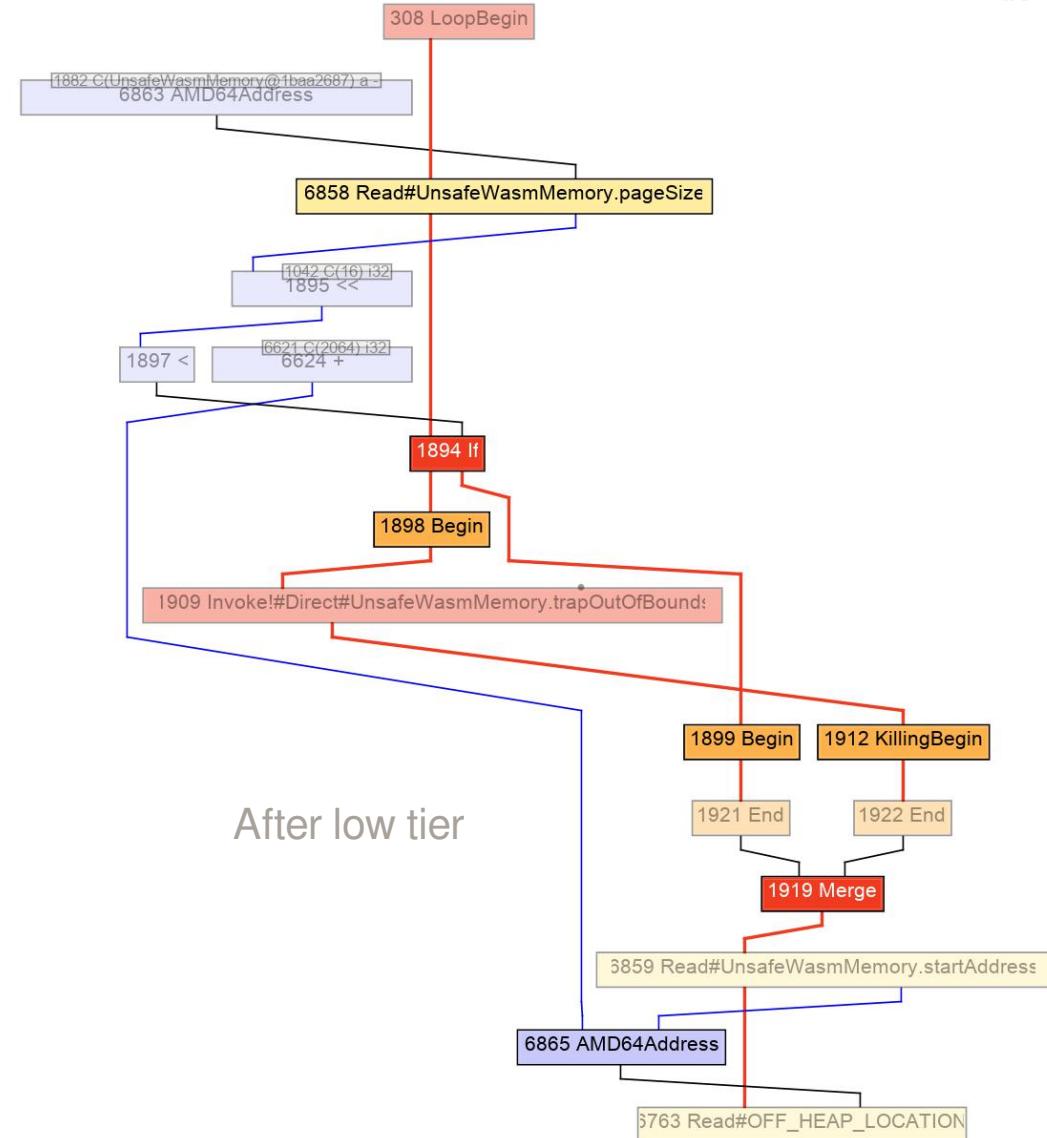
We have 3 memory-read nodes for every memory read in the program

We would expect that the size of the heap and the start-address of the heap are loop-invariant.

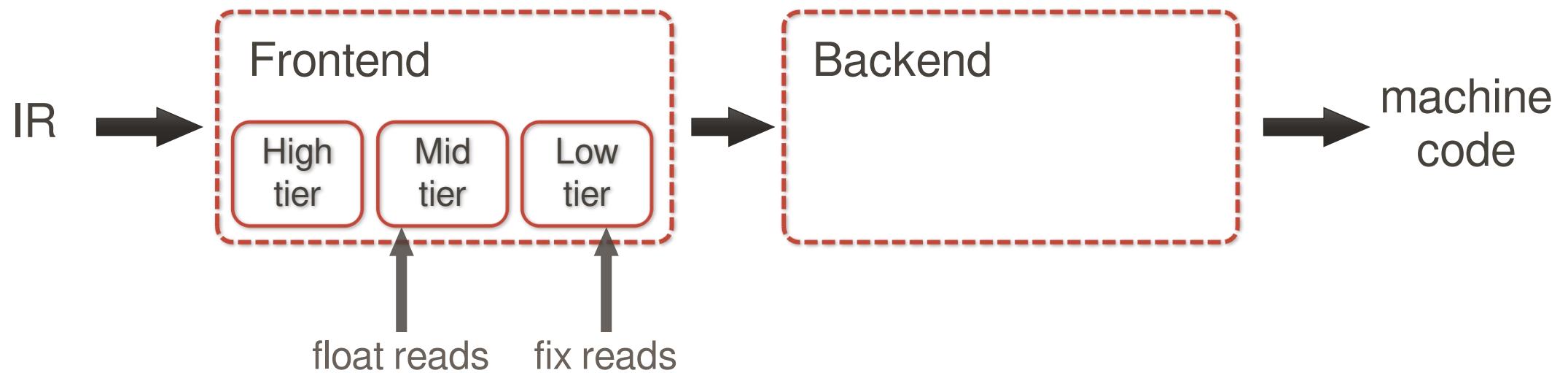


# Memory out-of-bounds access profiling

The compiler decides not to float  
the reads out of the loop - why?



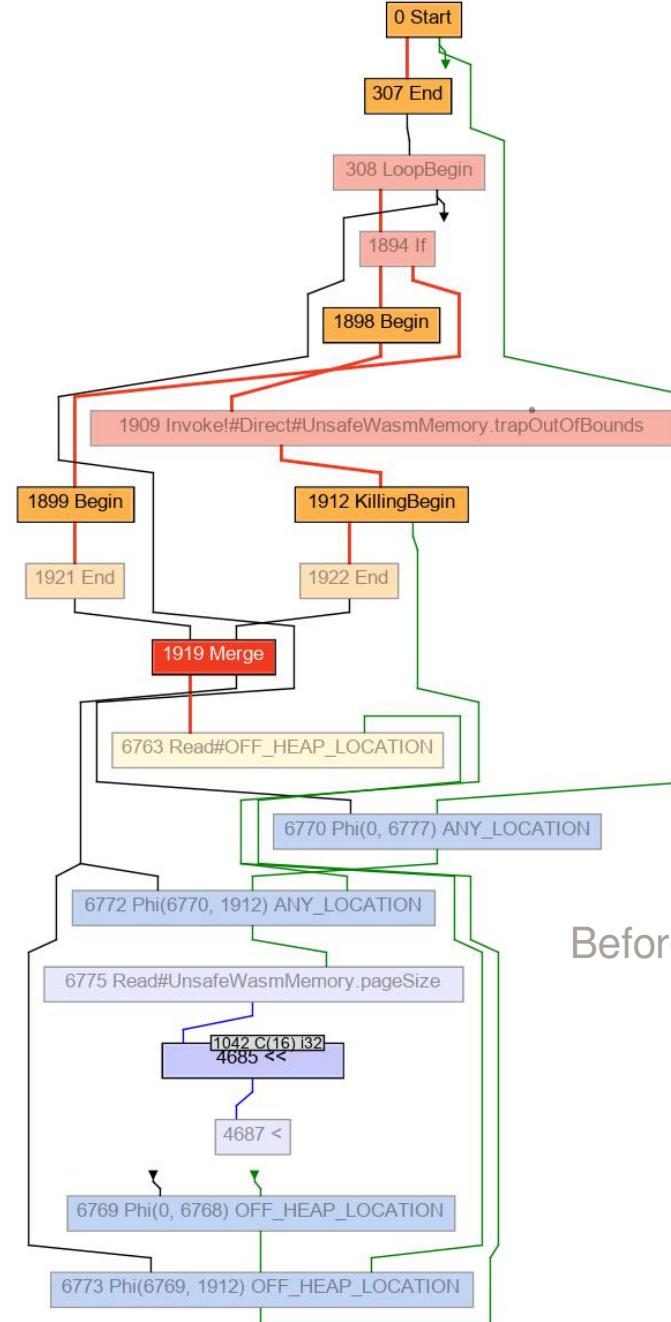
# Memory out-of-bounds access profiling



# Memory out-of-bounds access profiling

The compiler decides not to float the reads out of the loop - why?

To understand why, we need to inspect the memory graph that gets produced once the reads get floated.



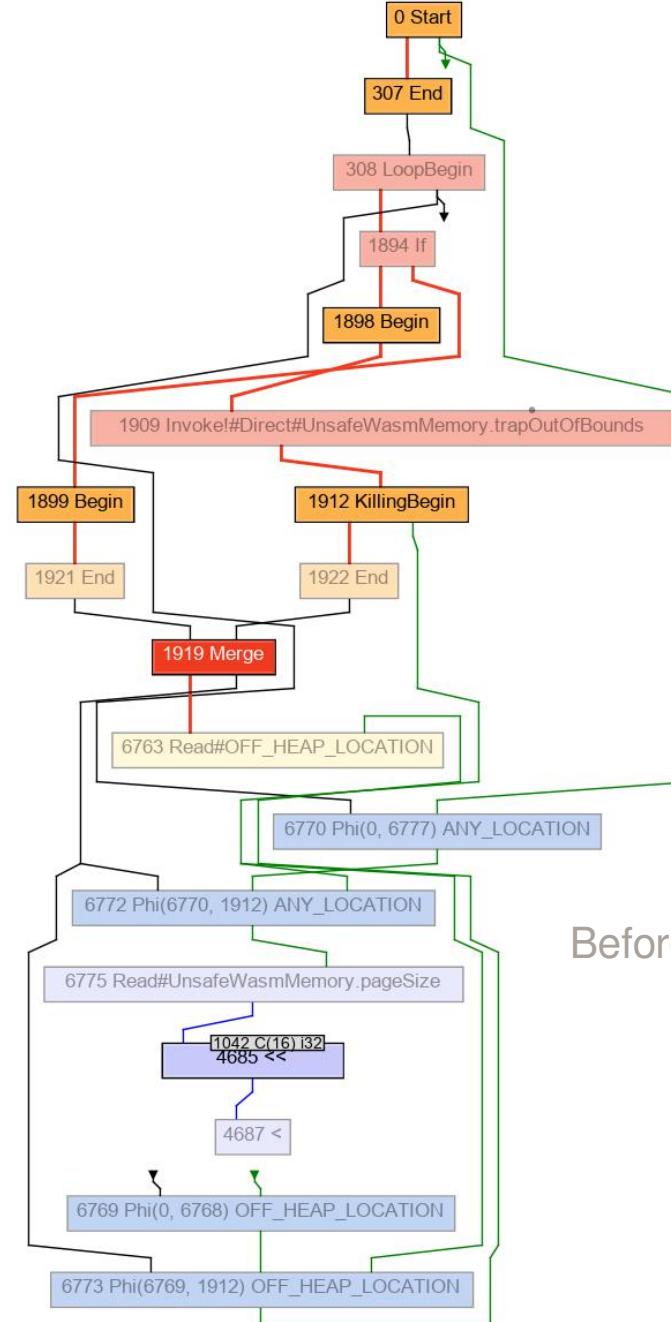
Before fix-reads

# Memory out-of-bounds access profiling

The compiler decides not to float the reads out of the loop - why?

To understand why, we need to inspect the memory graph that gets produced once the reads get floated.

Heavily simplified model - a read cannot be scheduled before its (either value or memory) inputs.

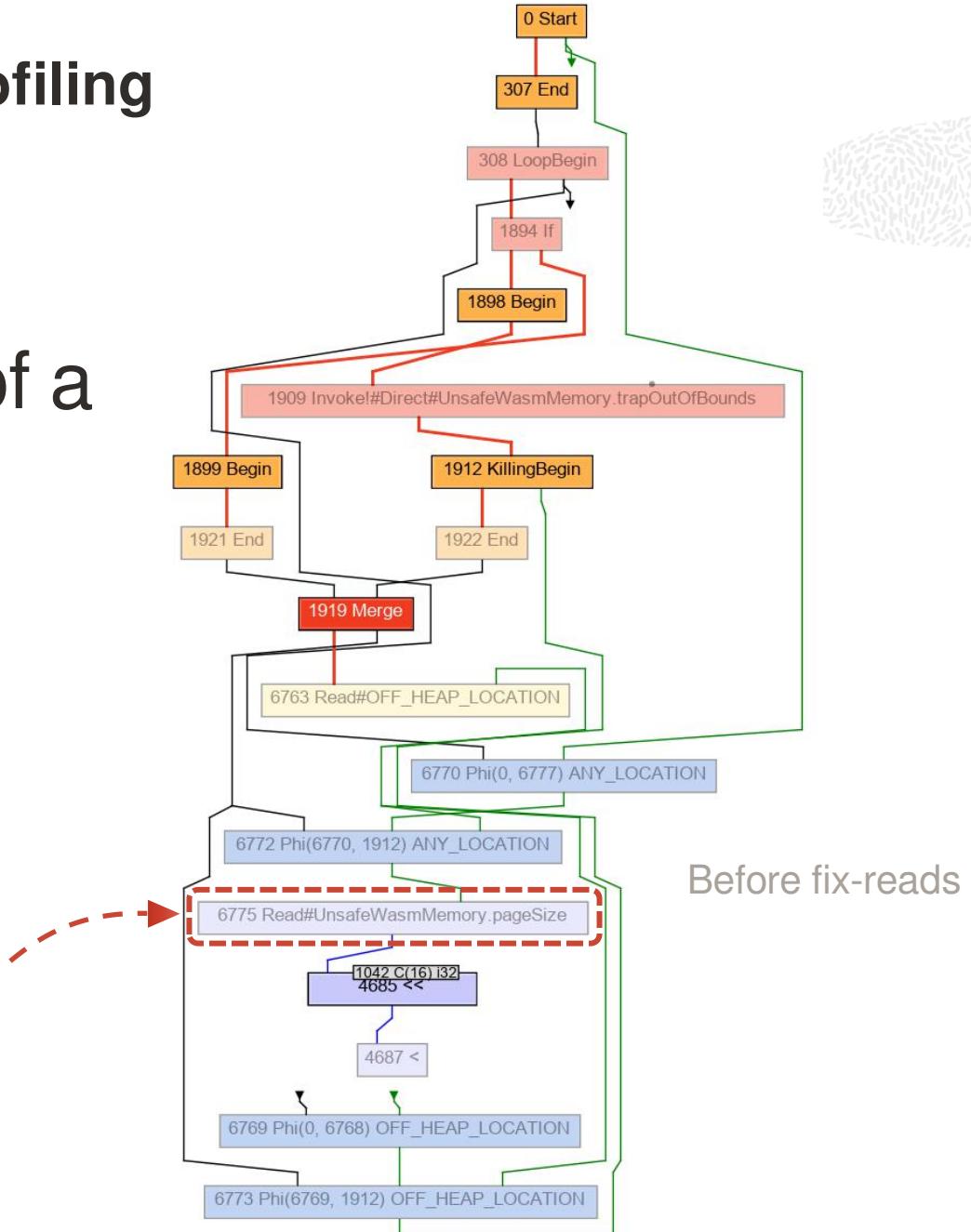


Before fix-reads

# Memory out-of-bounds access profiling

The memory graph consists of a set of ordered effects

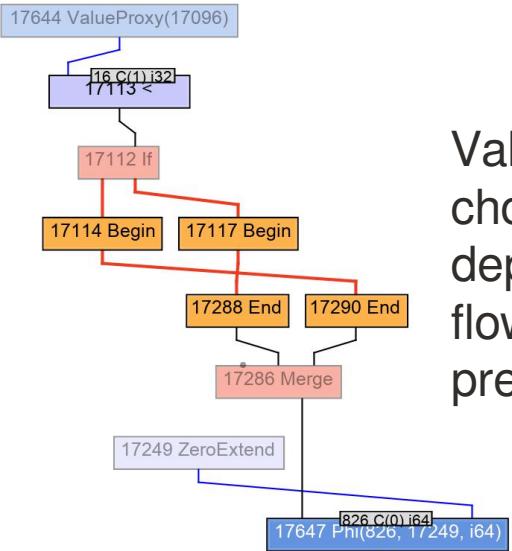
We find a read, and follow its inputs.



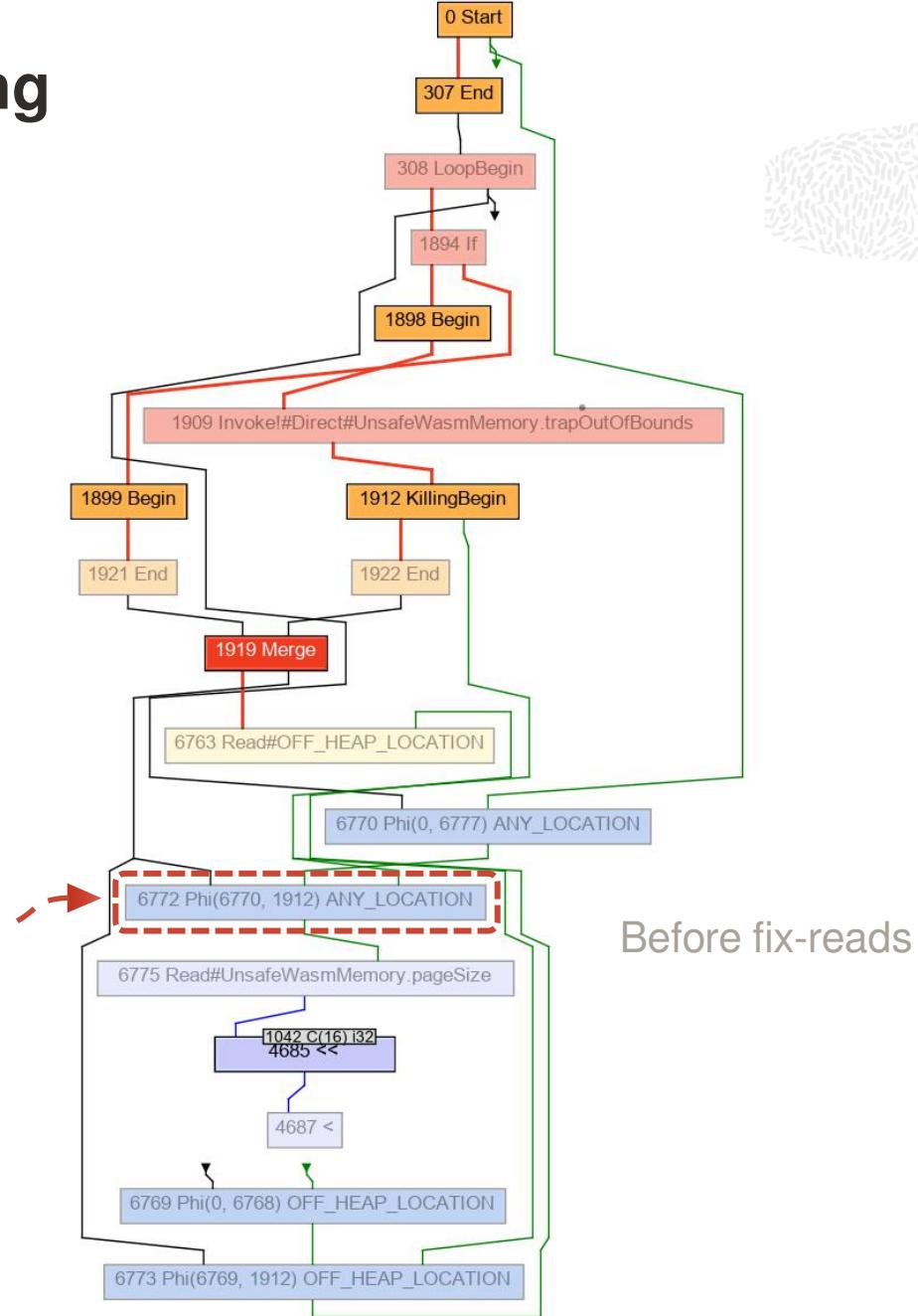
# Memory out-of-bounds access profiling

The memory graph consists of a set of ordered effects

We find a read, and follow its inputs.



Value-phi nodes represent a choice between two values, depending on which control flow the program has previously traversed.



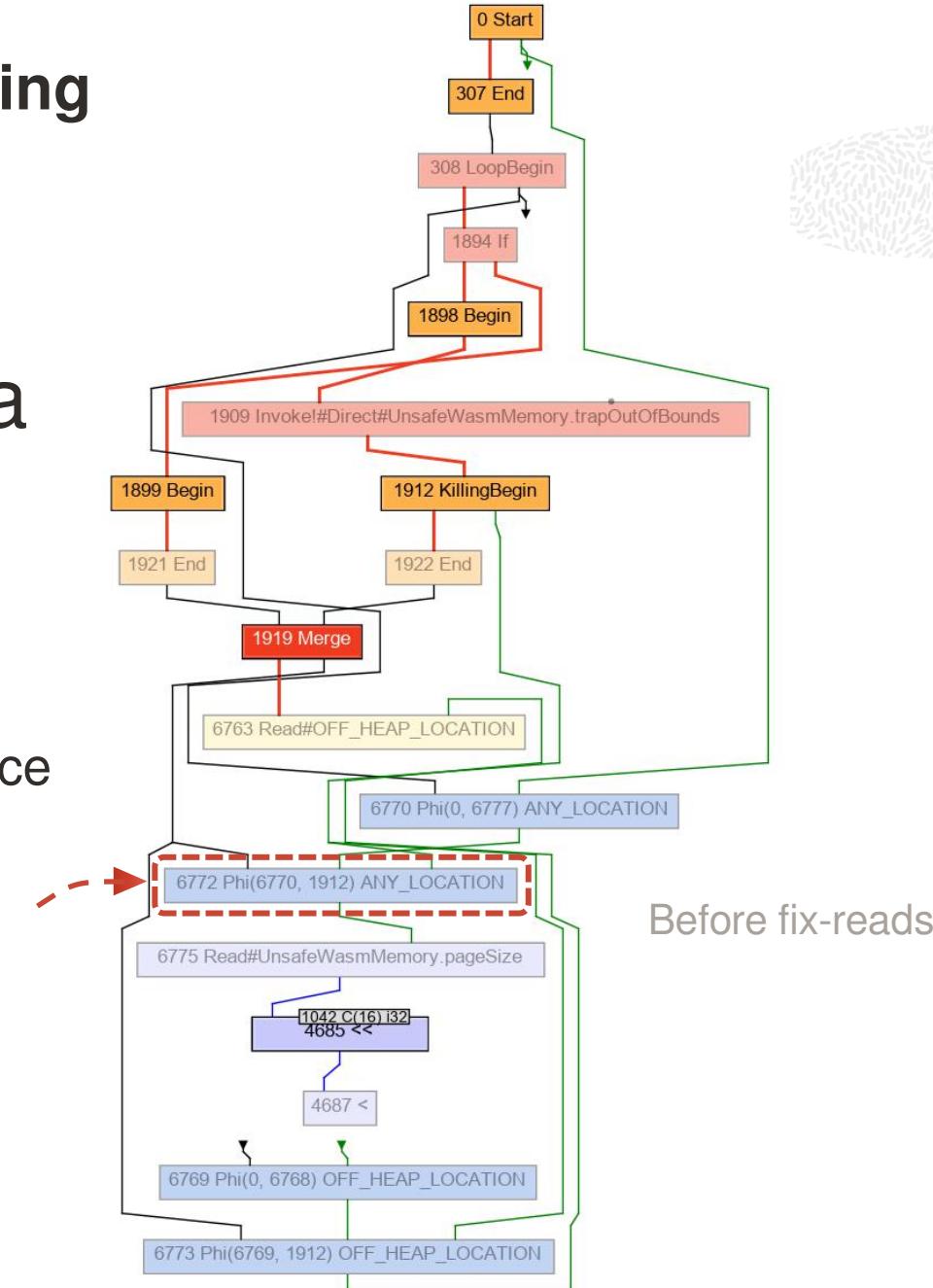
Before fix-reads

# Memory out-of-bounds access profiling

The memory graph consists of a set of ordered effects

We find a read, and follow its inputs.

Memory-phi nodes (green edges) represent a choice between two *memory states*, depending on which control flow the program has previously traversed.



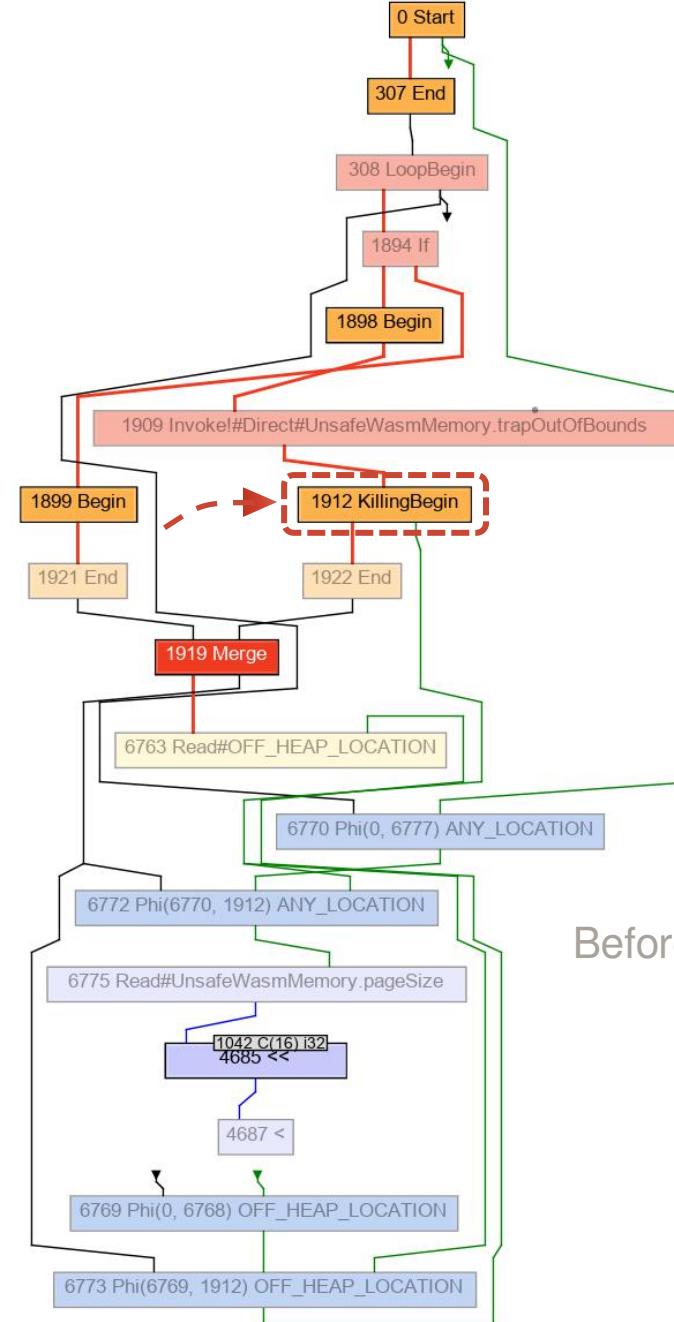
# Memory out-of-bounds access profiling

The memory graph consists of a set of ordered effects

We find a read, and follow its inputs.

Memory-phi nodes (green edges) represent a choice between two *memory states*, depending on which control flow the program has previously traversed.

We reach a KillingBegin node - beginning of a basic block that could follow a write to a memory location.



Before fix-reads

# Memory out-of-bounds access profiling

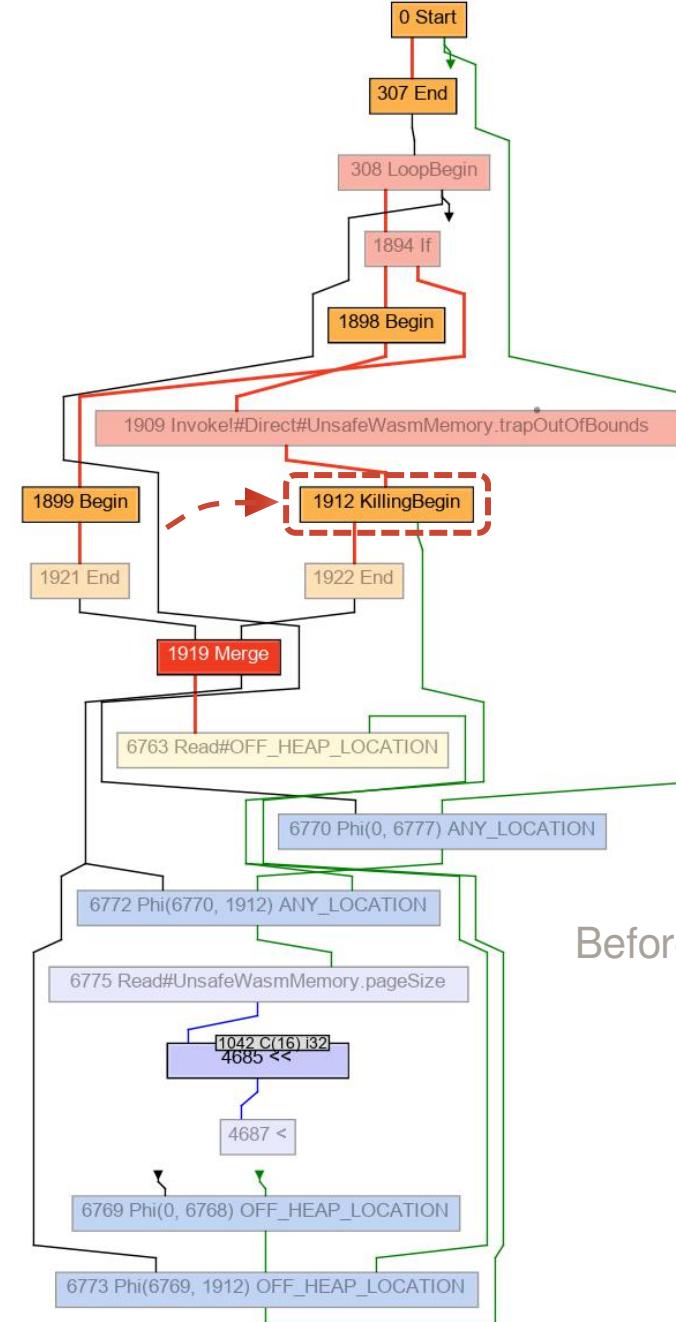
The memory graph consists of a set of ordered effects

We find a read, and follow its inputs.

Memory-phi nodes (green edges) represent a choice between two *memory states*, depending on which control flow the program has previously traversed.

We reach a KillingBegin node - beginning of a basic block that could follow a write to a memory location.

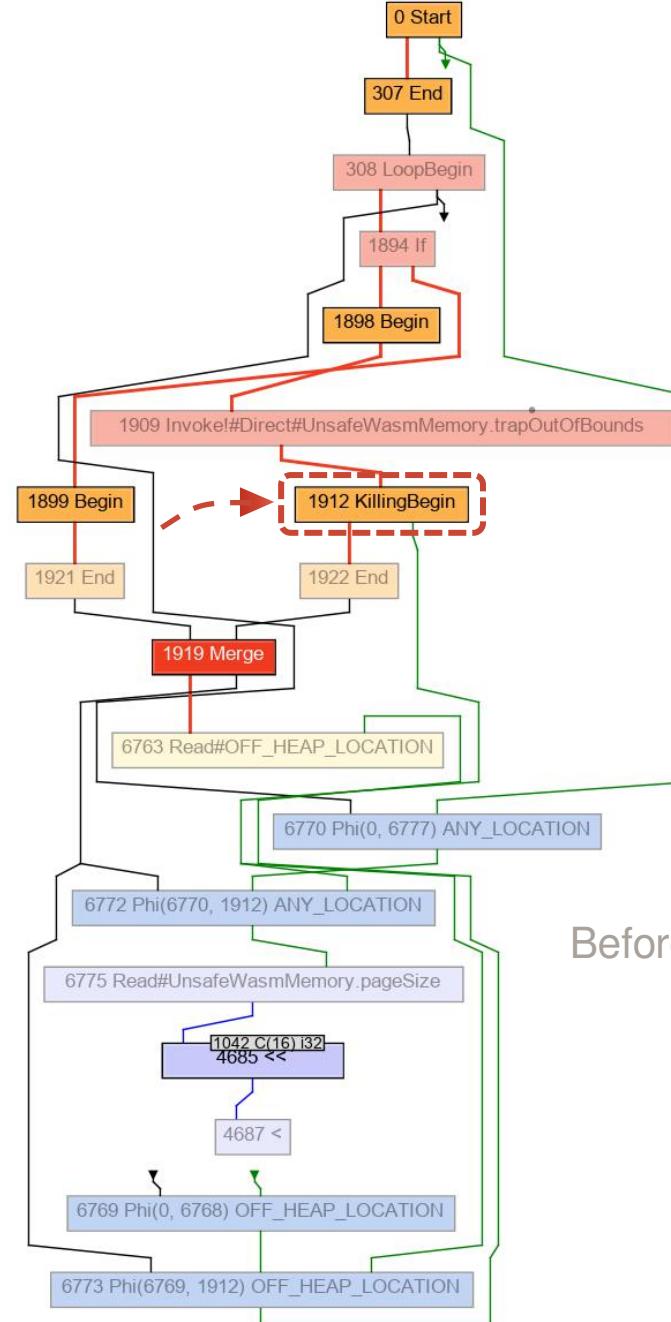
Unsurprising: preceding Invoke may write to memory.



# Memory out-of-bounds access profiling

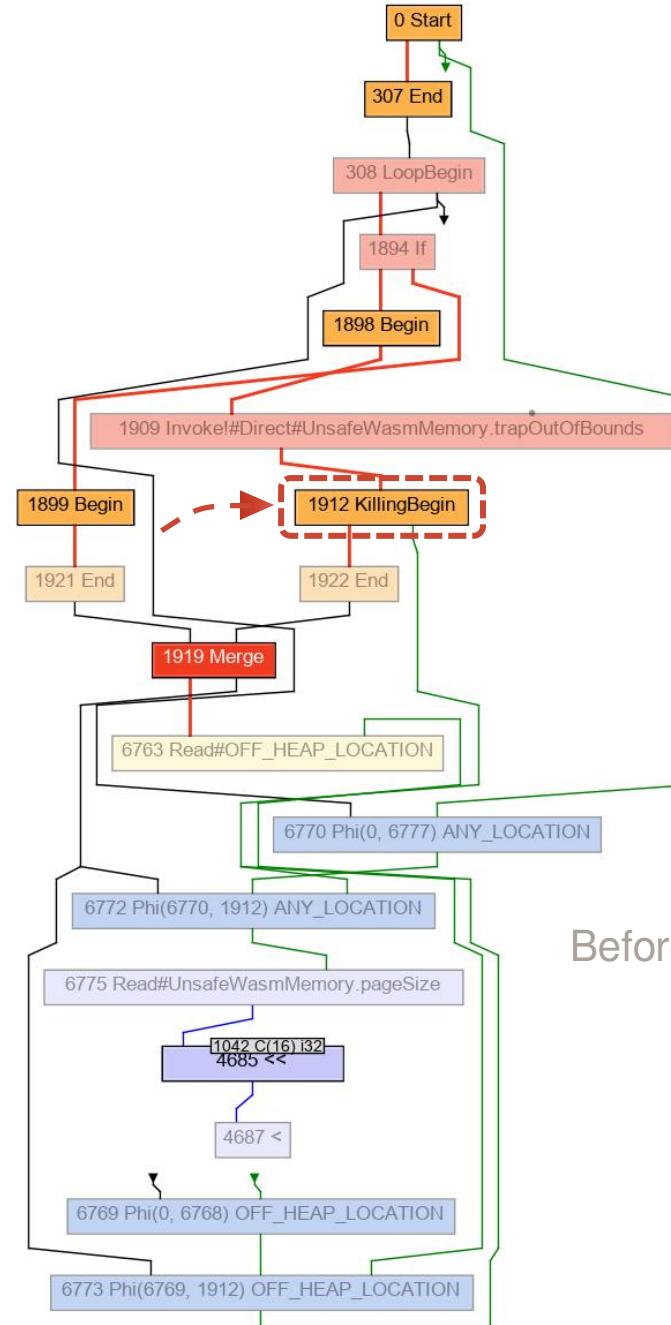
Conclusion: we cannot schedule the read before a call to a method that may have written to memory

If we did schedule the read above the Invoke, then we would read a stale value.



# Memory out-of-bounds access profiling

Could the compiler inline the call  
to trapOutOfBounds?



Before fix-reads

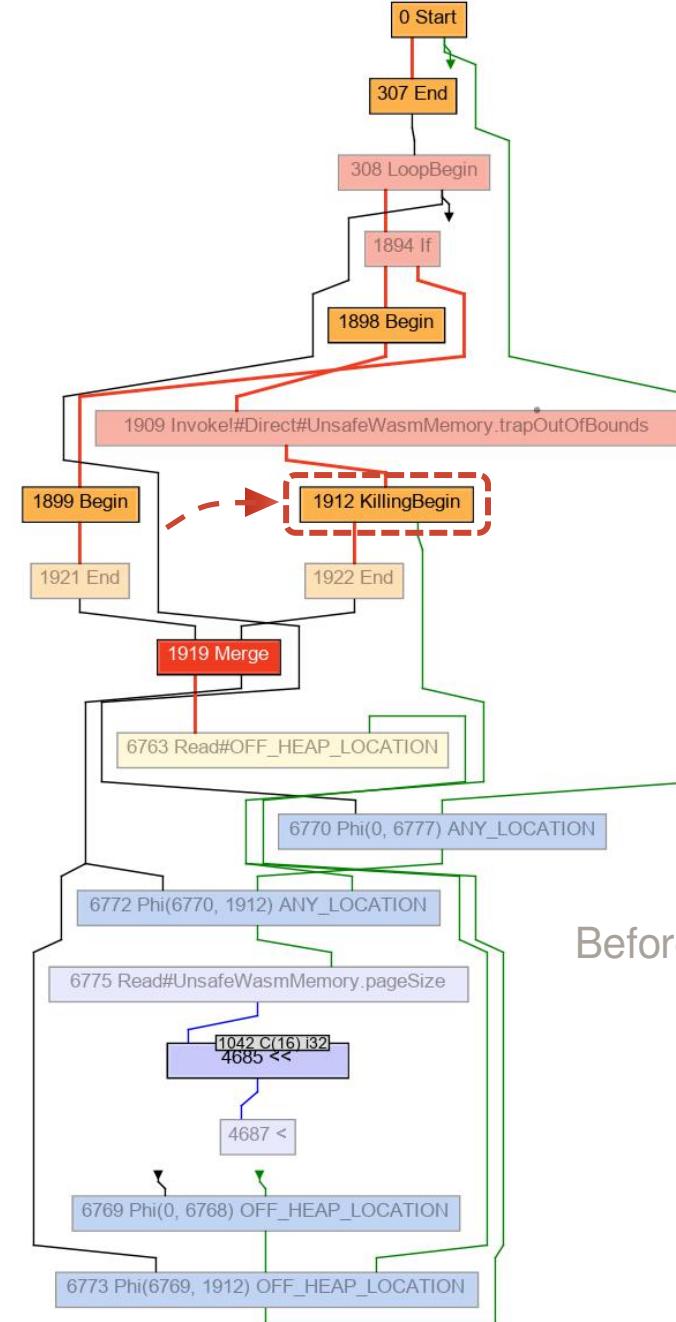
# Memory out-of-bounds access profiling

Could the compiler inline the call to trapOutOfBounds?

No, because this call is a Truffle boundary, which was added to prevent the exception creation code from getting partially evaluated.

```
@TruffleBoundary
void trapOutOfBounds(long address, long offset) {
    throw new WasmTrap(
        "Out-of-bounds: " + (address + offset));
}
```

GraalWasm



Before fix-reads

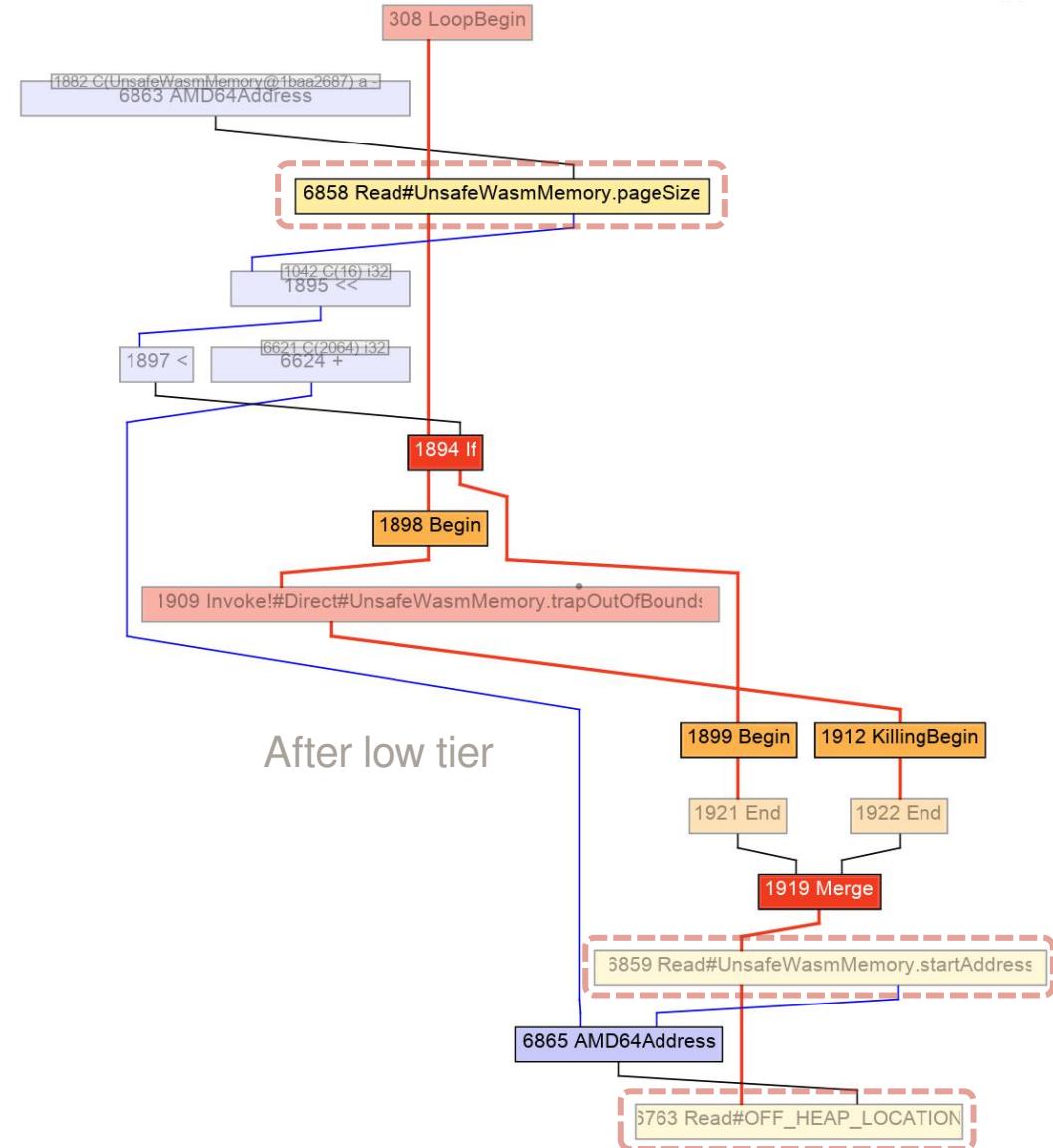
# Memory out-of-bounds access profiling

We have 3 memory-read nodes for every memory read in the program

Normally, the language should annotate these values as `@CompilationFinal` values, so that partial evaluation can inline them into the code.

```
class WasmMemory {  
    @CompilationFinal long startAddress;  
    @CompilationFinal long pageSize;  
    ...
```

GraalWasm



# Memory out-of-bounds access profiling

In WebAssembly, memory can grow

There is an instruction that can attempt to resize the memory of the program.

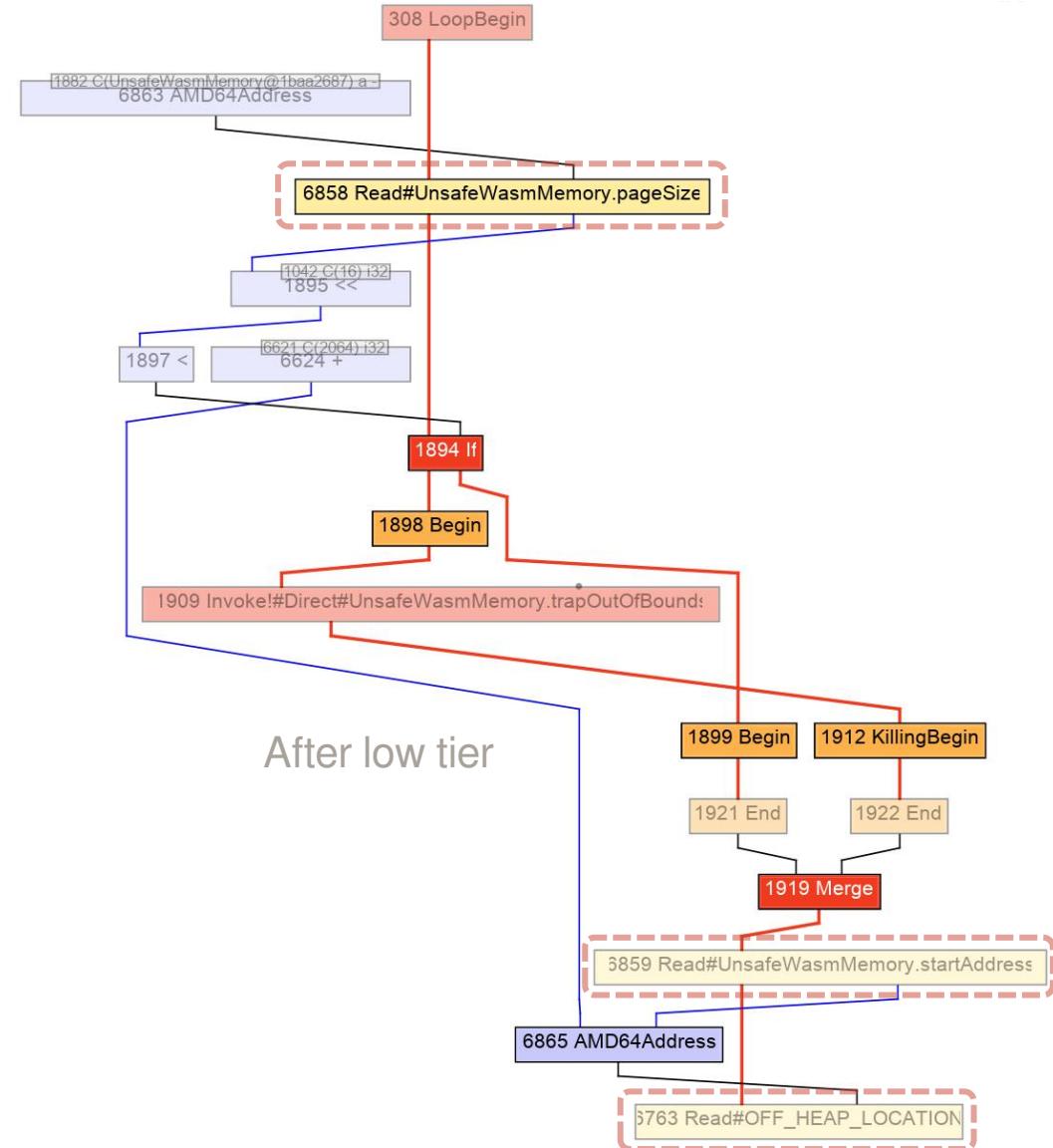
The startAddress and pageSize cannot be constant!

```
class WasmMemory {  
    long startAddress;  
    long pageSize;  
    ...
```

GraalWasm

...  
i32.load  
i32.store  
memory.grow  
...

WebAssembly



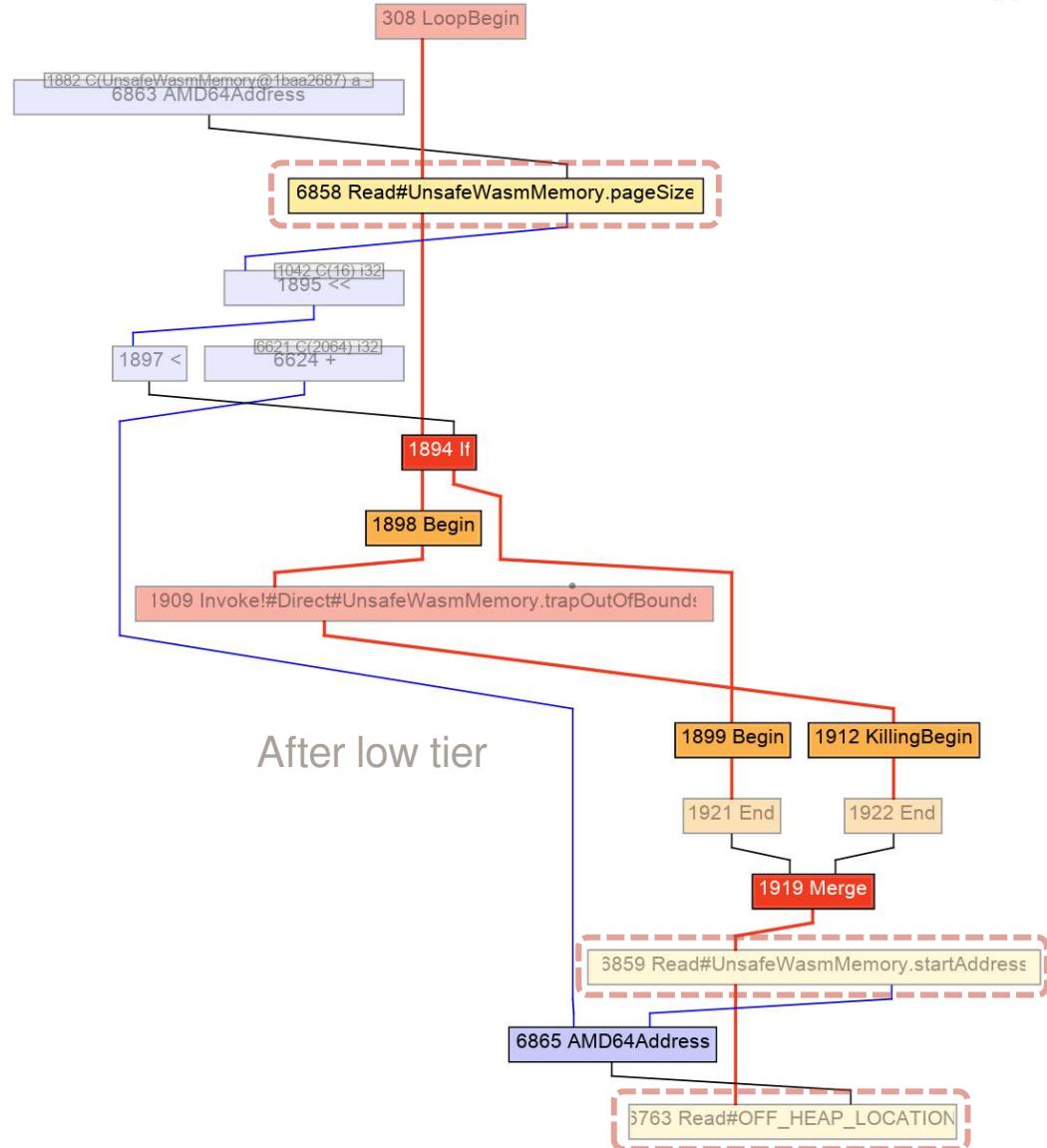
# Memory out-of-bounds access profiling

So, is there anything we can do?

```
class WasmMemory {  
    long startAddress;  
    long pageSize;  
    ...  
}
```

GraalWasm

...  
i32.load  
i32.store  
memory.grow  
...  
WebAssembly



# Memory out-of-bounds access profiling

So, is there anything we can do?

When in doubt, speculate.

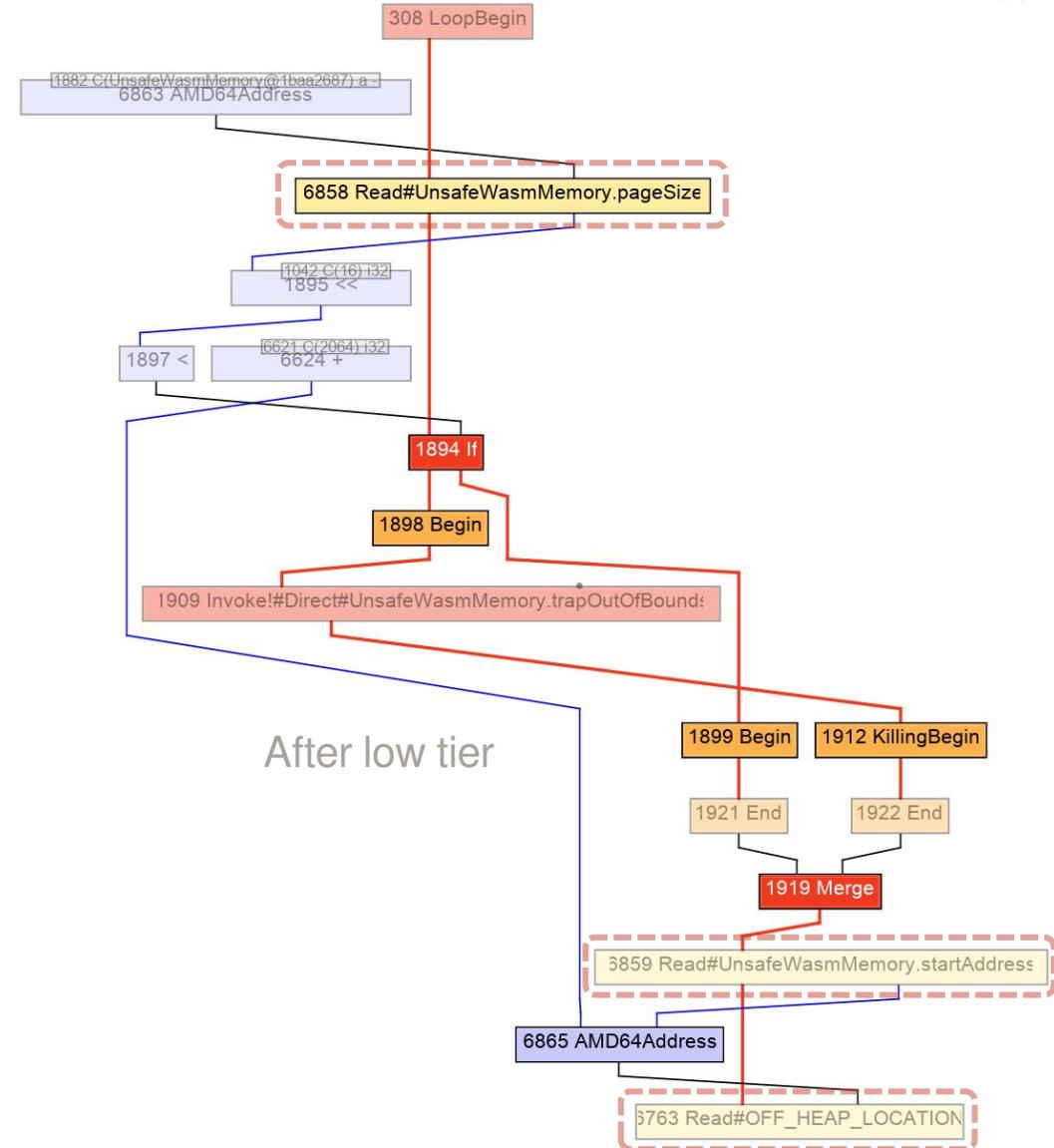
How often does an out-of-bounds access error happen?

```
class WasmMemory {  
    long startAddress;  
    long pageSize;  
    ...  
}
```

GraalWasm

...  
i32.load  
i32.store  
memory.grow  
...

WebAssembly

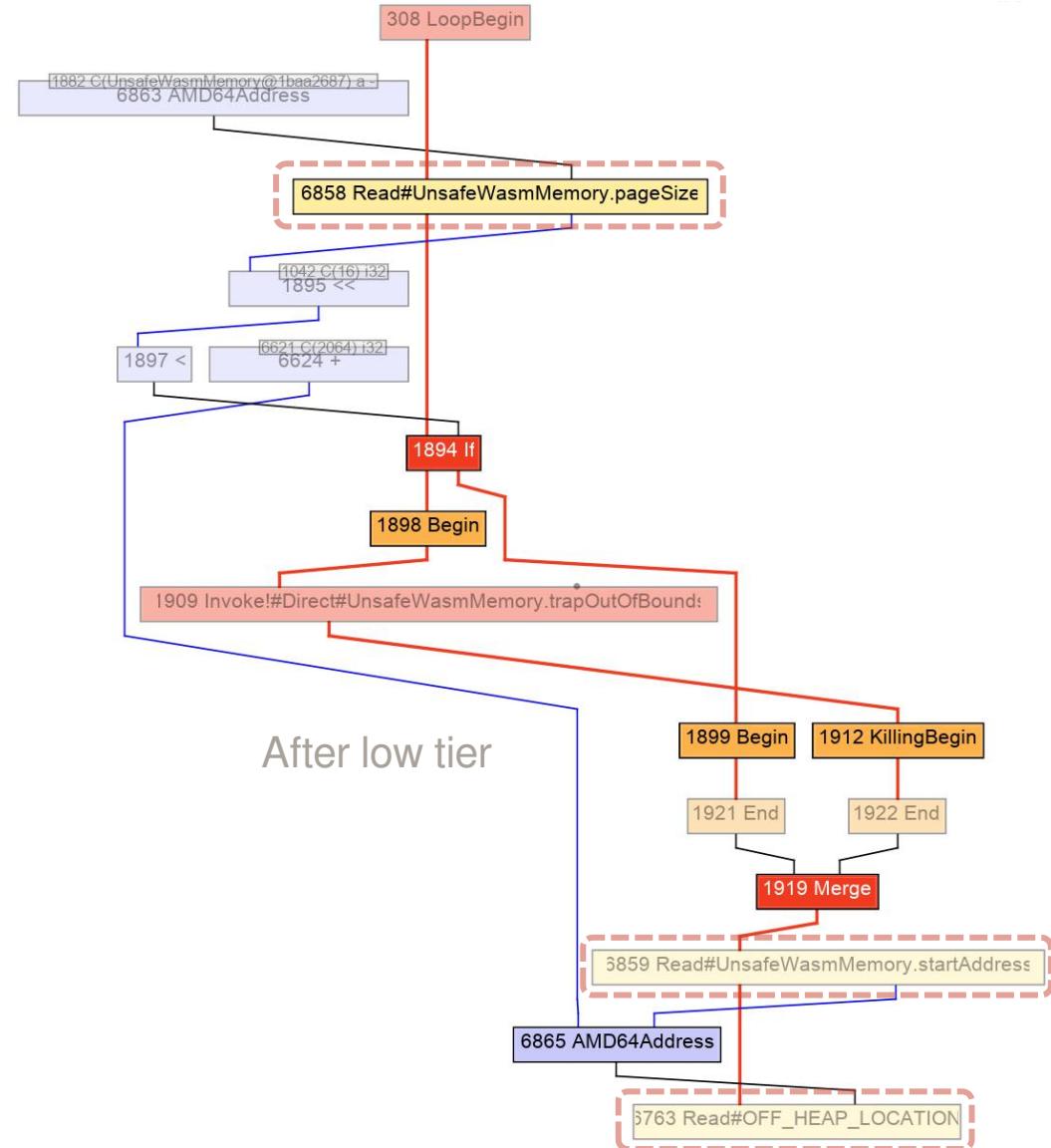


# Memory out-of-bounds access profiling

Out-of-bounds memory access must terminate the program

```
int load_i32(long address) {
    if (address + 4 > pageCount * PAGE_SIZE) {
        trapOutOfBounds(address, 4);
    }
    return unsafe.getInt(startAddress + address);
}
```

GraalWasm



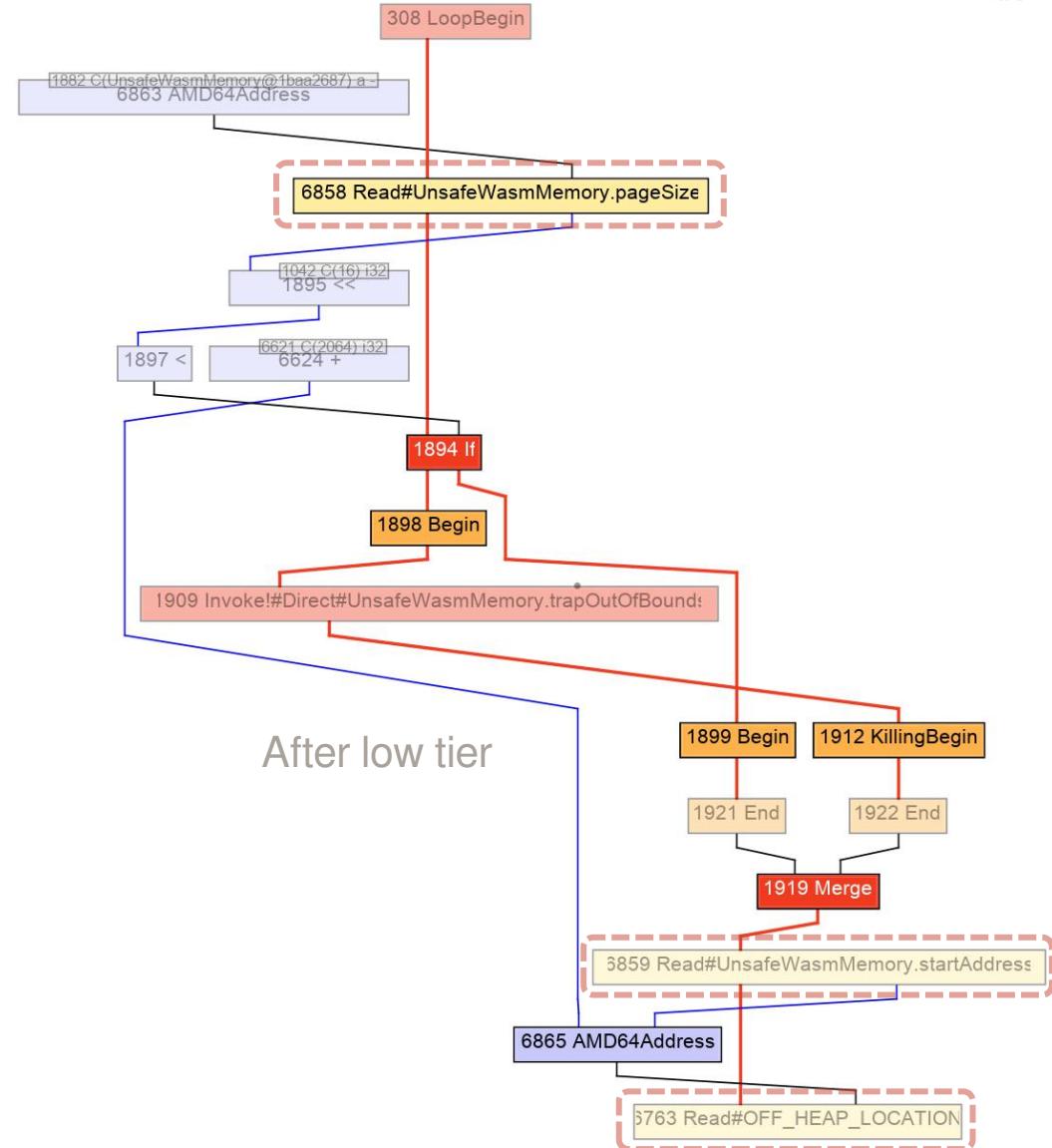
# Memory out-of-bounds access profiling

Out-of-bounds memory access must terminate the program

Therefore, it will never happen while the program is running. We can therefore speculate that the branch that throws the out-of-bounds error will never happen.

```
int load_i32(long address) {
    if (address + 4 > pageCount * PAGE_SIZE) {
        trapOutOfBounds(address, 4);
    }
    return unsafe.getInt(startAddress + address);
}
```

GraalWasm



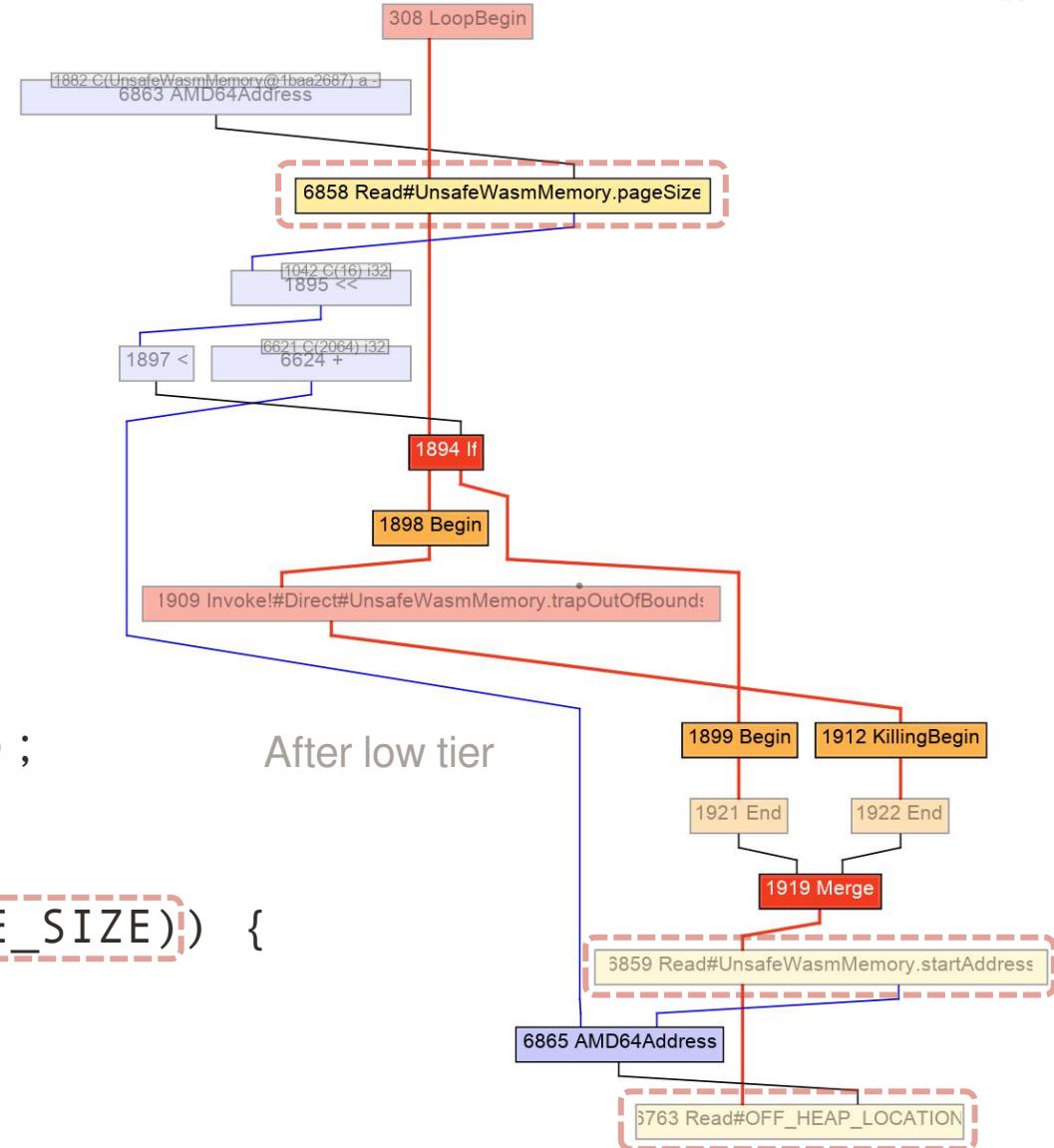
# Memory out-of-bounds access profiling

We speculate that the out-of-bounds branch with a ConditionProfile

```
ConditionProfile oob = ConditionProfile.create();

int load_i32(long address) {
    if (oob.profile(address + 4 > pageCount * PAGE_SIZE)) {
        trapOutOfBounds(address, 4);
    }
    return unsafe.getInt(startAddress + address);
}
```

GraalWasm



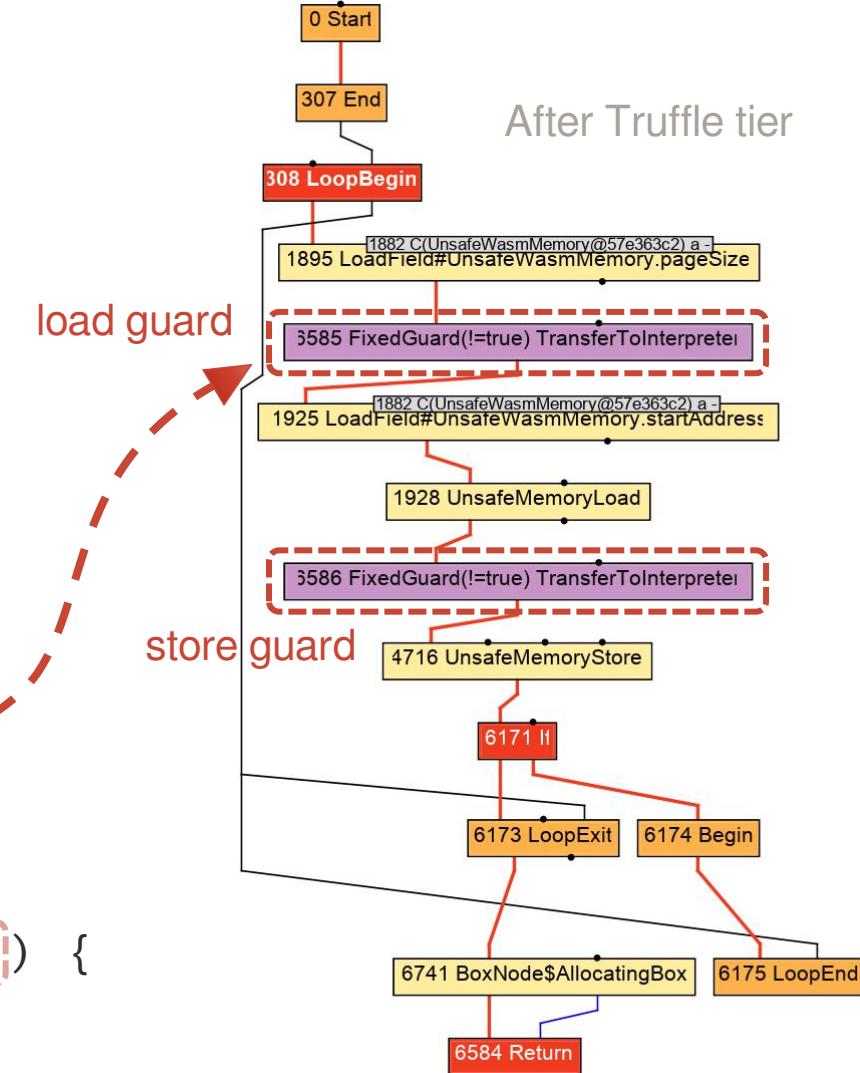
# Memory out-of-bounds access profiling

A profile whose count is 0 will introduce Guard nodes into the IR

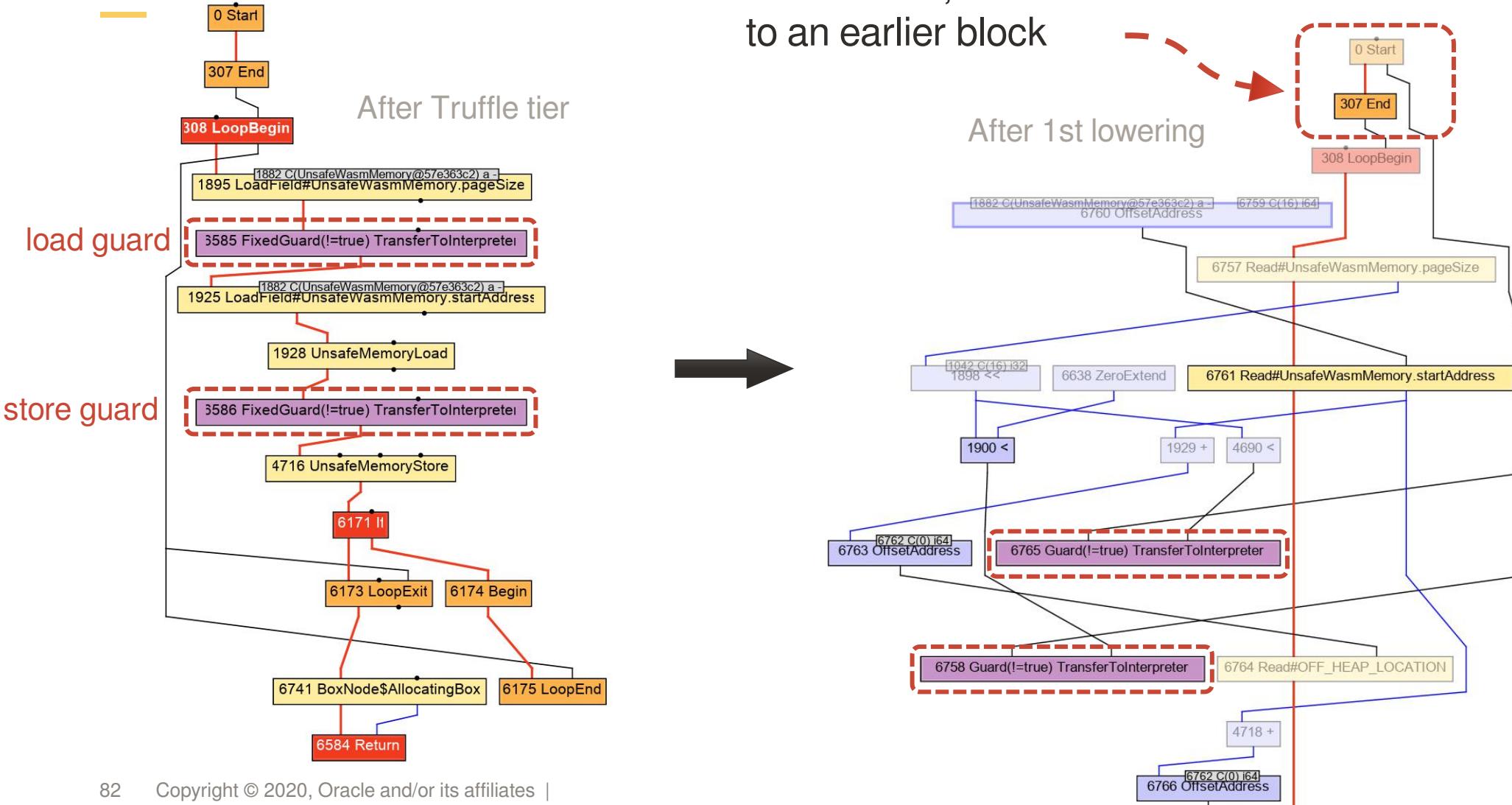
The compiler moves the guard out of the loop when its condition does not depend on the values produced in that loop.

```
int load_i32(long address) {
    if (oop.profile(address + 4 > pageCount * PAGE_SIZE)) {
        trapOutOfBounds(address, 4);
    }
    return unsafe.getInt(startAddress + address);
}
```

GraalWasm

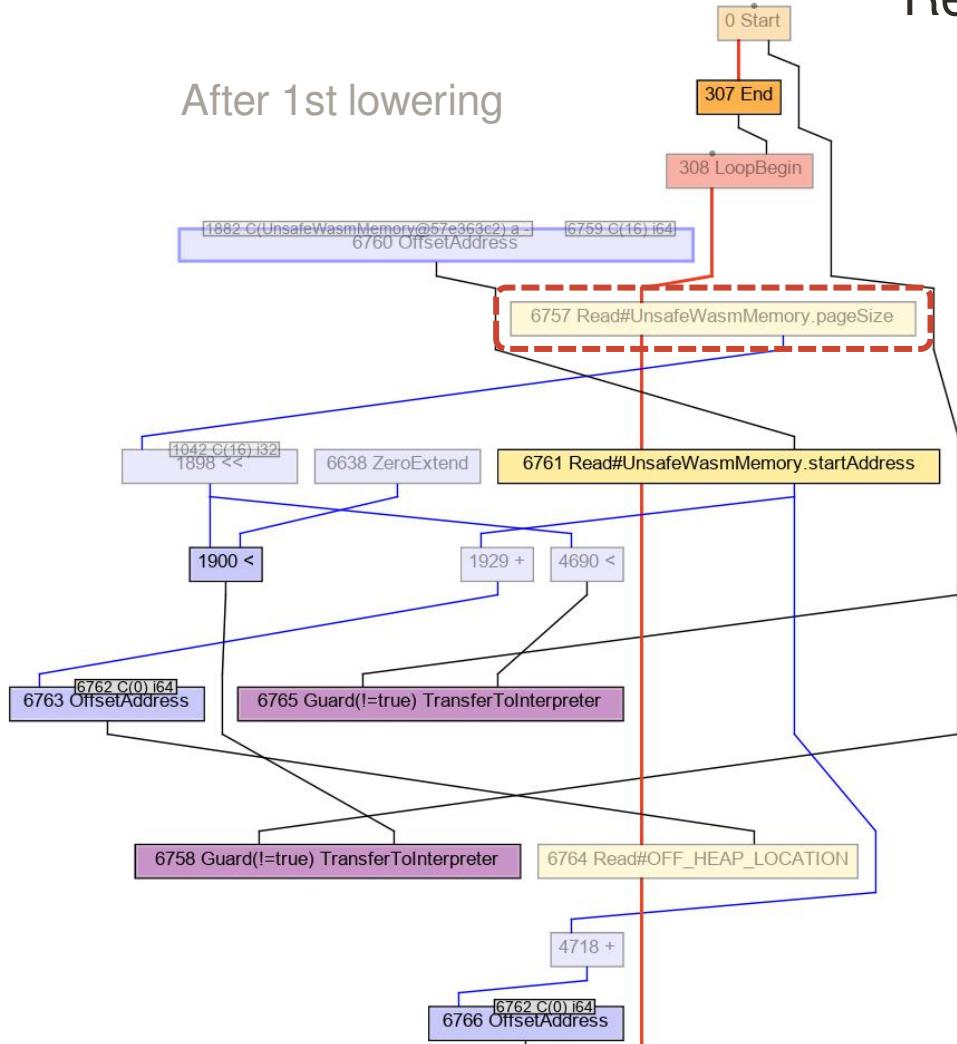


# Memory out-of-bounds access profiling



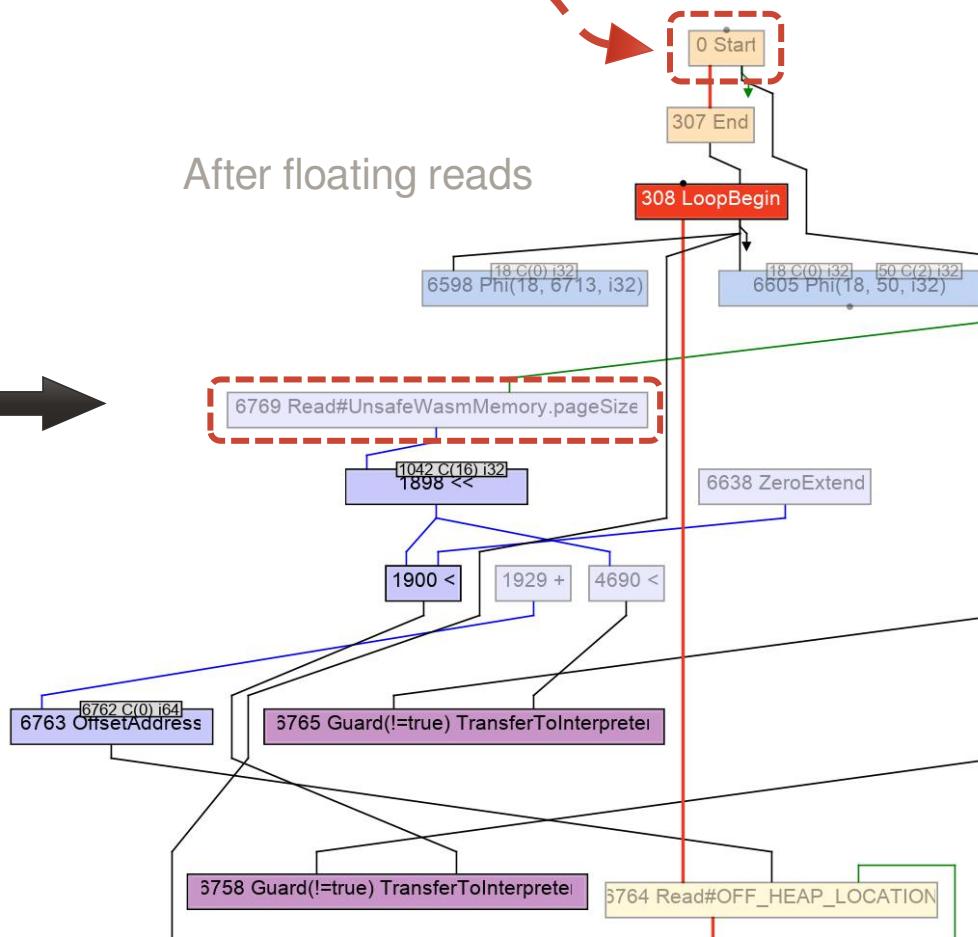
# Memory out-of-bounds access profiling

After 1st lowering



Reads float, and take the previous  
memory effect as input

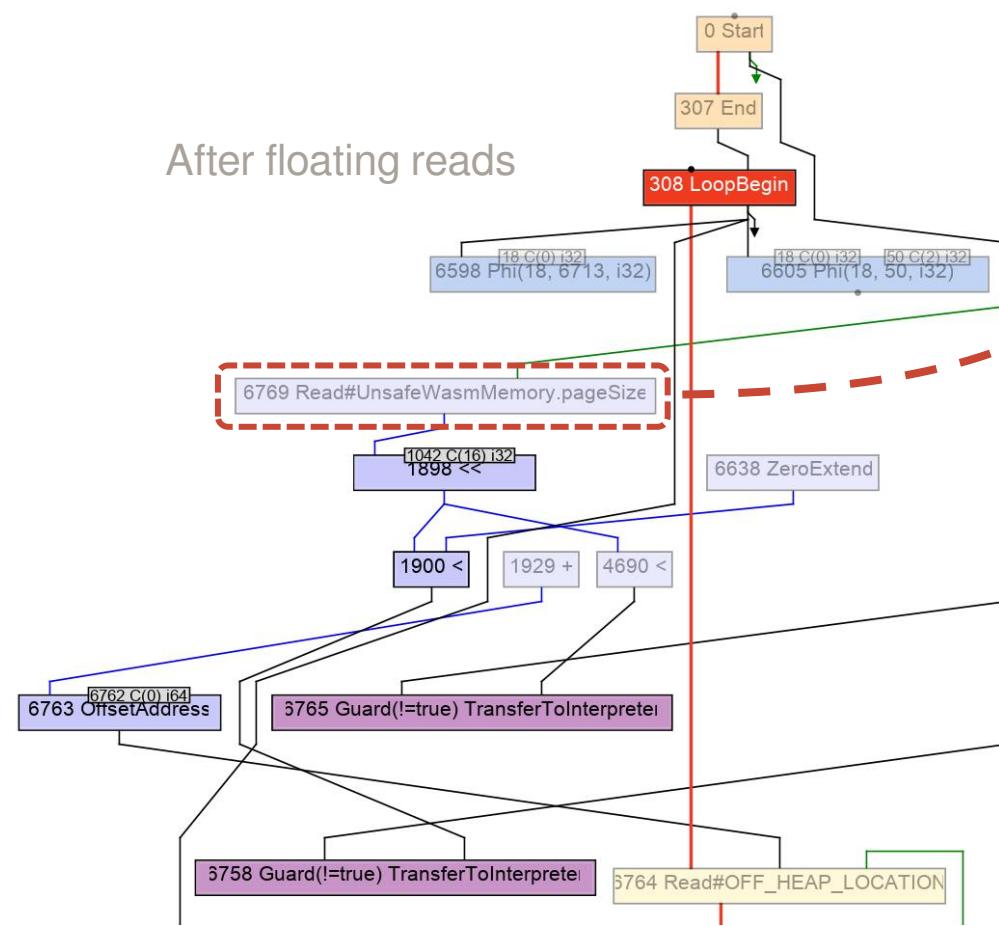
After floating reads



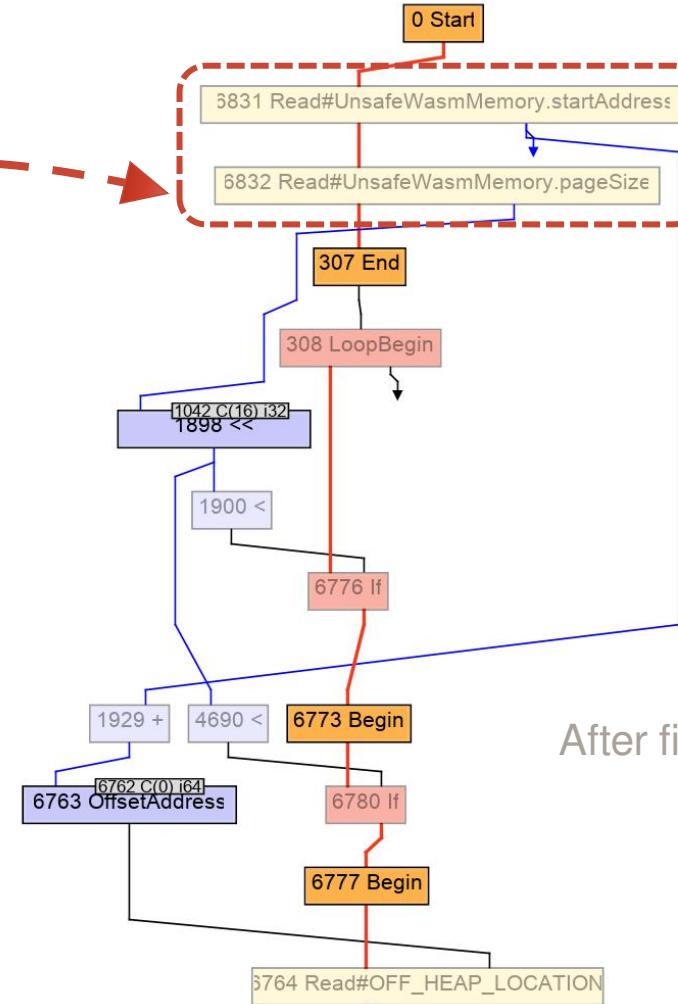
# Memory out-of-bounds access profiling

Reads are then scheduled into basic blocks according to their inputs

After floating reads



After fix-reads

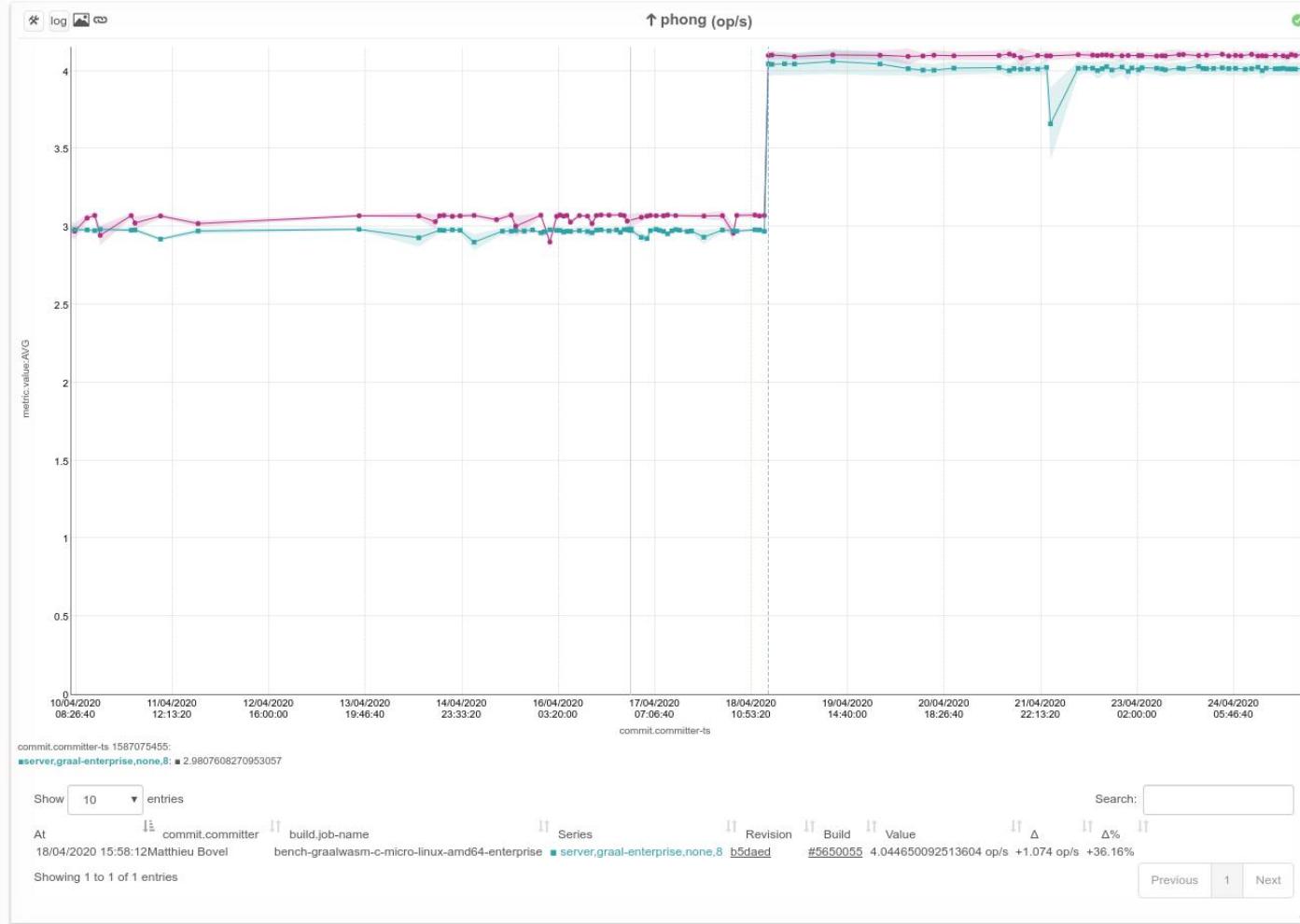


# Memory out-of-bounds access profiling



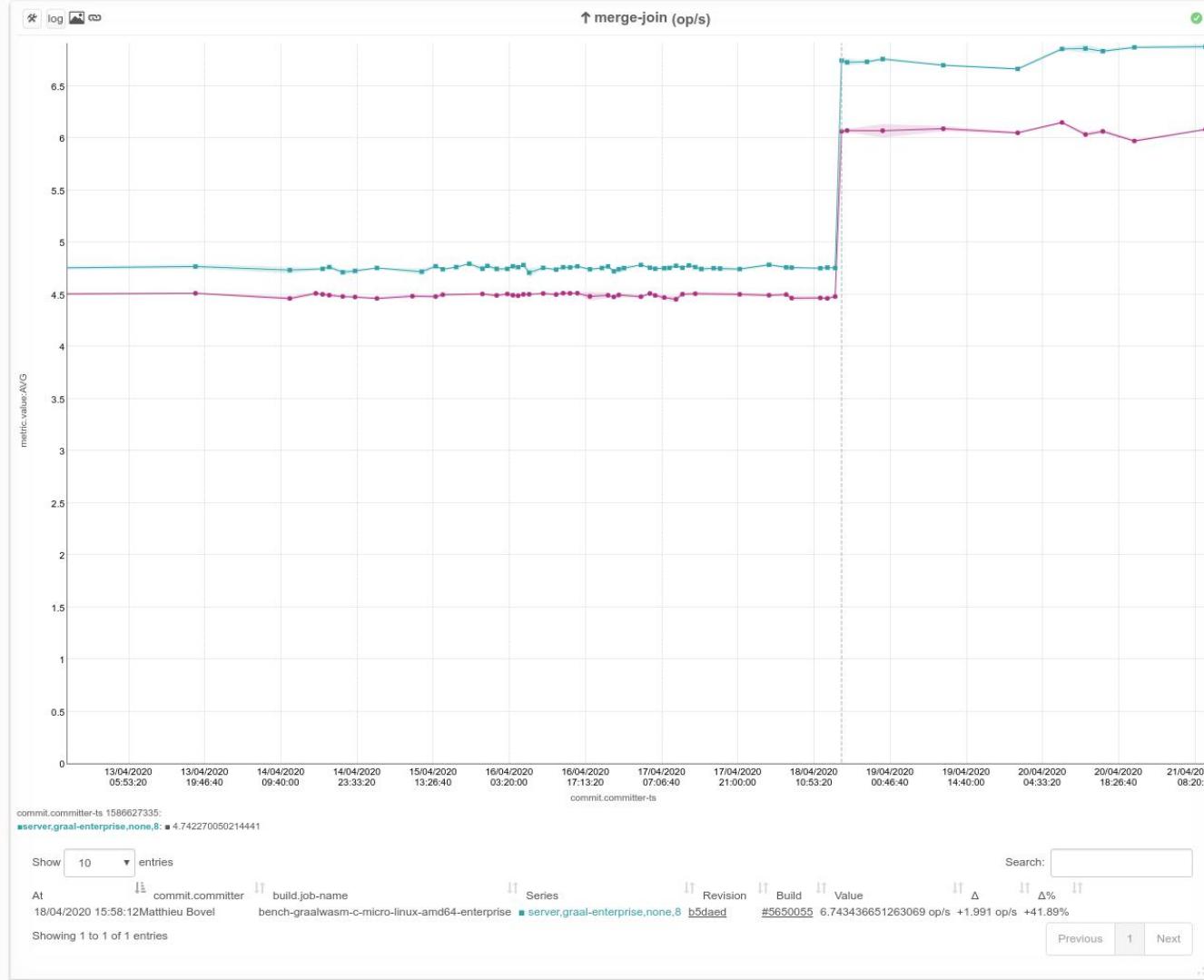
- 14% improvement on a discrete-event simulation benchmark (*event-sim*)

# Memory out-of-bounds access profiling



- 14% improvement on a discrete-event simulation benchmark (*event-sim*)
- 39% improvement on a 3d-renderer benchmark (*phong*)

# Memory out-of-bounds access profiling



- 14% improvement on a discrete-event simulation benchmark (*event-sim*)
- 39% improvement on a 3d-renderer benchmark (*phong*)
- 42% improvement on a merge-join benchmark (*merge-join*)
- improvements on other benchmarks in the 5%-60% range

# Memory out-of-bounds access profiling

Example program: posterize filter



# Resetting Truffle frame-slot values

---

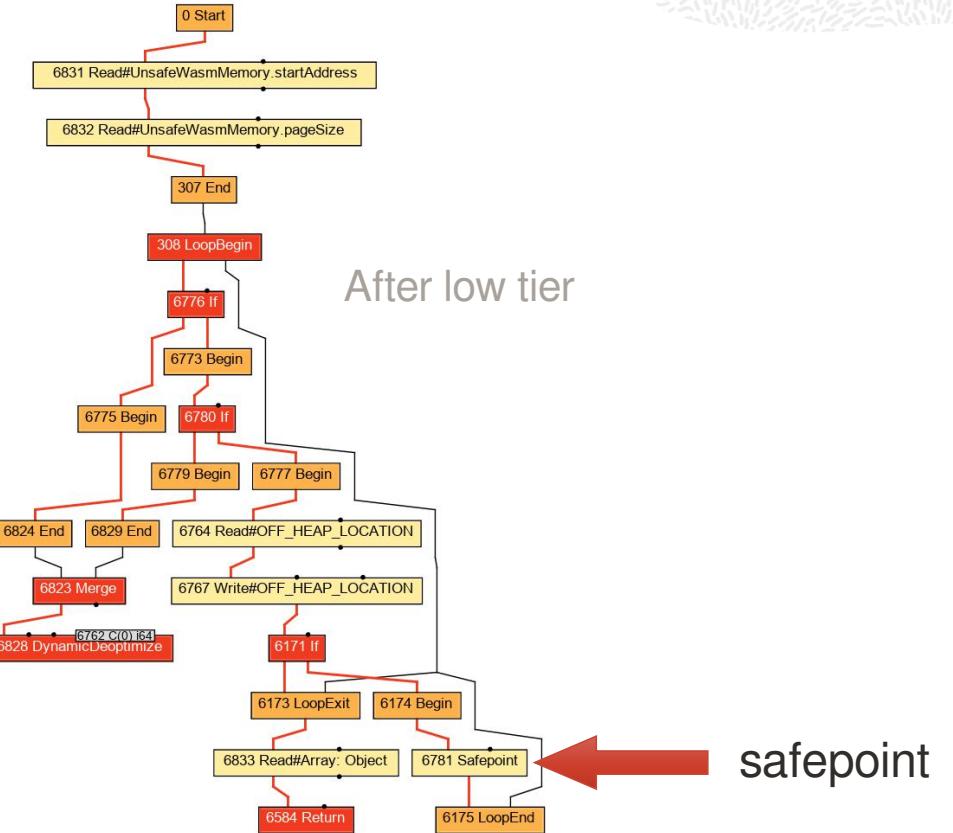
## Example program: posterize filter

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {
    uint32_t color = image[pixel % IMAGE_SIZE];
    uint8_t R = (color & 0xFF000000) >> 24;
    uint8_t G = (color & 0x00FF0000) >> 16;
    uint8_t B = (color & 0x0000FF00) >> 8;
    double luminance = (0.2126 * R + 0.7152 * G + 0.0722 * B);
    result[pixel % IMAGE_SIZE] = luminance > 127 ? UINT32_MAX : 0xFF;
}
```

# Truffle frame-slot values

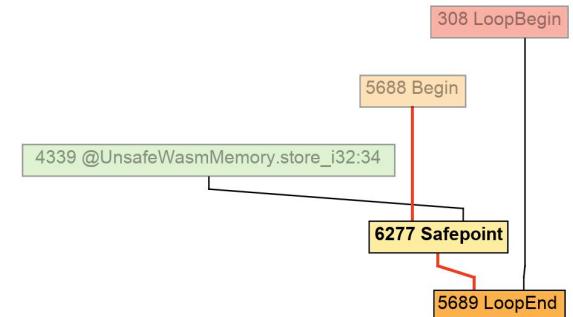
In GraalVM, a program can deoptimize at safepoints

When this happens, the values of the local variables must be copied back to the corresponding frame slots of the interpreter.



# Truffle frame-slot values

Frame-state nodes encode local variable state of the Java program

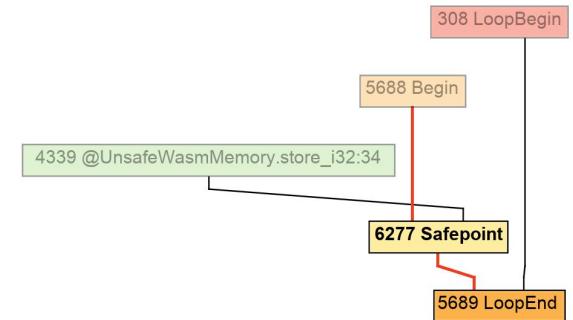


After low tier

# Truffle frame-slot values

Frame-state nodes encode local variable state of the Java program

A Truffle interpreter is a Java program, so the values of its local variables must be linked to the frame state node.



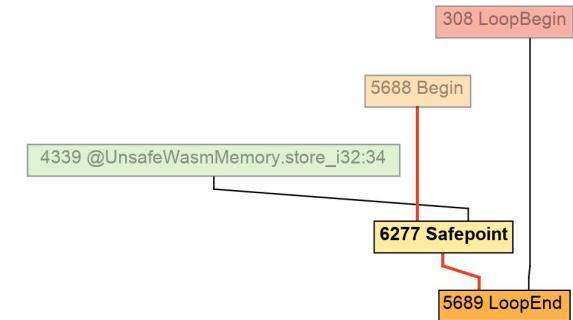
After low tier

# Truffle frame-slot values

Frame-state nodes encode local variable state of the Java program

A Truffle interpreter is a Java program, so the values of its local variables must be linked to the frame state node.

```
int execute(VirtualFrame frame) {  
    int offset = initialOffset;  
    ...
```



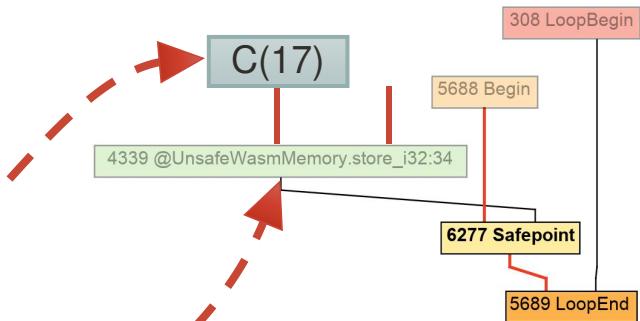
After low tier

# Truffle frame-slot values

Frame-state nodes encode local variable state of the Java program

A Truffle interpreter is a Java program, so the values of its local variables must be linked to the frame state node.

```
int execute(VirtualFrame frame) {  
    int offset = initialOffset;  
    ...
```



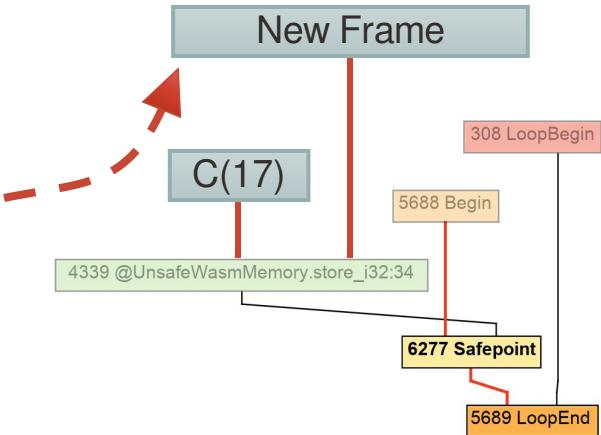
After low tier

# Truffle frame-slot values

Truffle Frames contain the values of the local variables of the guest program

One of the local variables in the interpreter is the Truffle frame, which represents the local variables of the guest program. The frame is represented as an allocation of the Frame object.

```
int execute(VirtualFrame frame) {  
    int offset = initialOffset;  
    ...
```



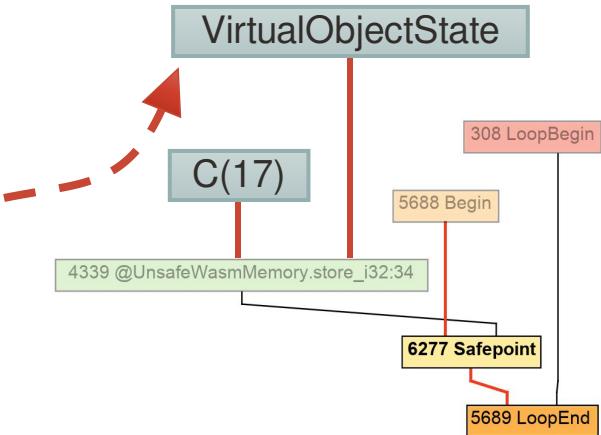
“After PE”

# Truffle frame-slot values

Truffle Frames contain the values of the local variables of the guest program

The escape analysis replaces the allocation node of the Frame with a virtual object state - a node that represents the state of the Frame, but does not require an allocation in the final code.

```
int execute(VirtualFrame frame) {  
    int offset = initialOffset;  
    ...
```



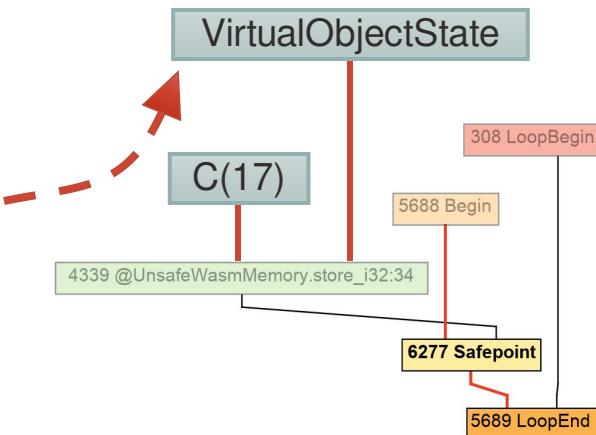
"After escape analysis"

# Resetting Truffle frame-slot values

Truffle Frames contain the values of the local variables of the guest program

If the compiled code of the guest program gets deoptimized, then the Truffle Frame object **must be allocated** (this is the same as with any other object in a Java program).

```
int execute(VirtualFrame frame) {  
    int offset = initialOffset;  
    ...
```



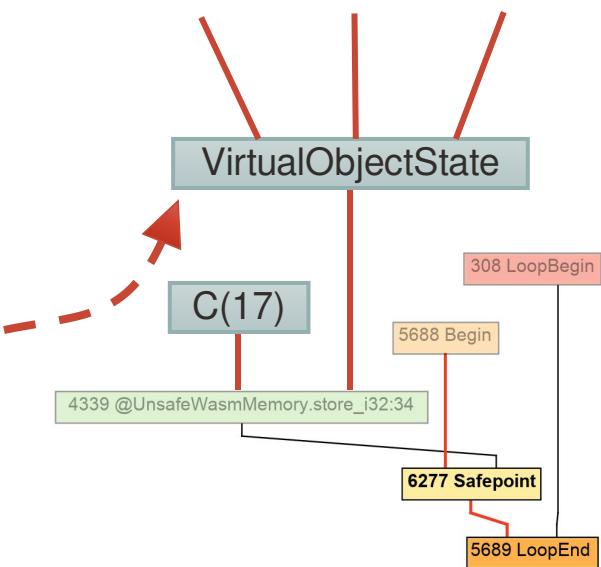
After low tier

# Truffle frame-slot values

Truffle Frames contain the values of the local variables of the guest program

By looking at the inputs of the VirtualObjectState node that corresponds to the Truffle Frame, we can see which values in the guest program were “saved”.

```
int execute(VirtualFrame frame) {  
    int offset = initialOffset;  
    ...
```

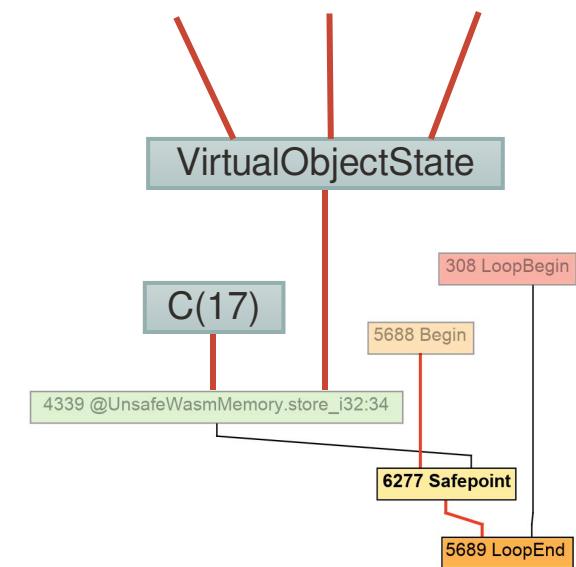


After low tier

# Truffle frame-slot values

The state of the Frame must be tracked in the compiled code

The corresponding values have to stay live at the point where the frame state is used. This means that they must be stored either in a register or on the program stack.



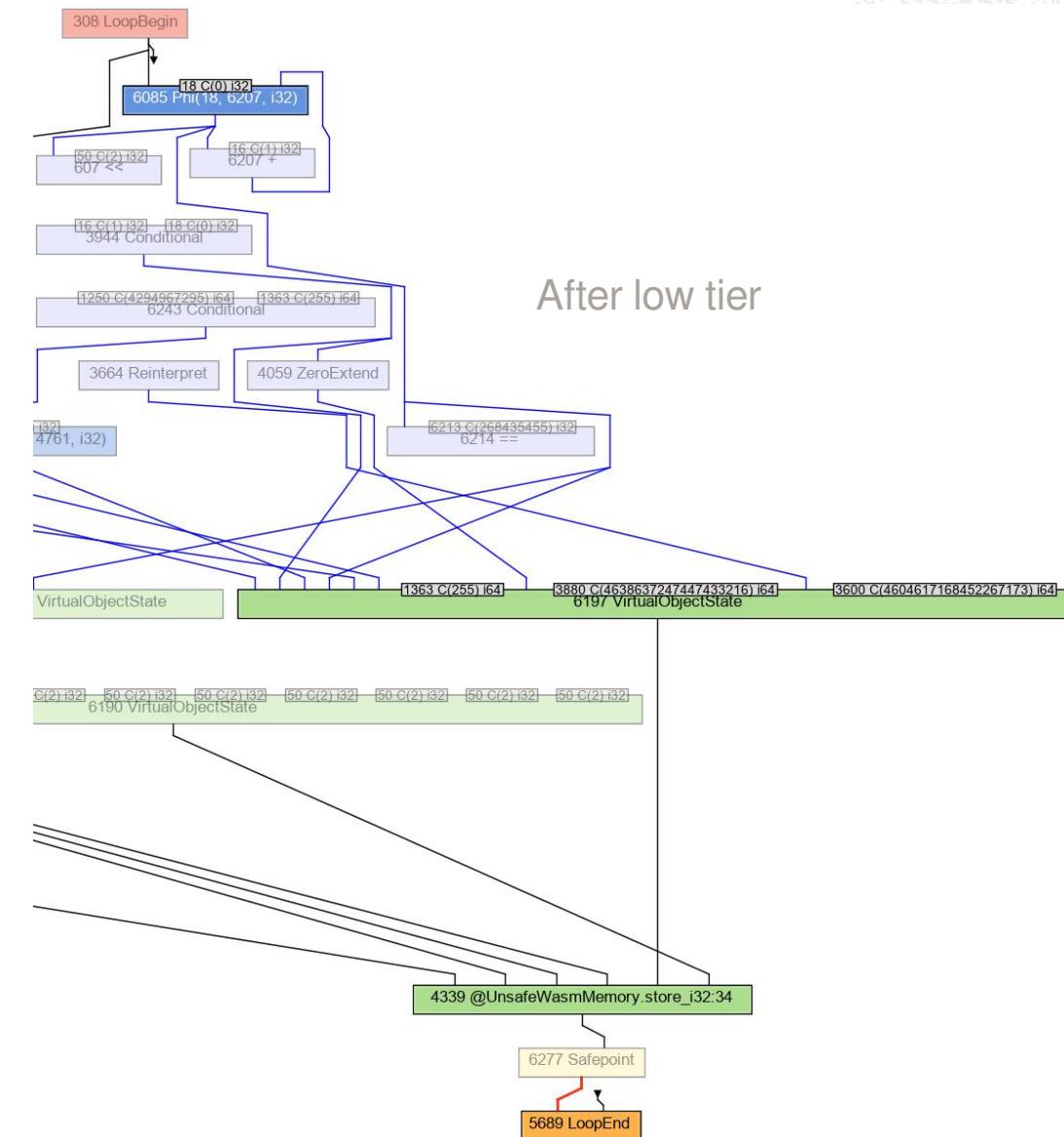
After low tier

# Truffle frame-slot values

By inspecting the state of the frame, we can see what's tracked

A lot of values tracked. Recall, the original C program.

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = (color & 0xFF000000) >> 24;  
    uint8_t G = (color & 0x00FF0000) >> 16;  
    uint8_t B = (color & 0x0000FF00) >> 8;  
    double luminance =  
        (0.2126 * R + 0.7152 * G + 0.0722 * B);  
    result[pixel % IMAGE_SIZE] =  
        luminance > 127 ? UINT32_MAX : 0xFF;  
    black_pixels += luminance > 127 ? 0 : 1;  
}
```

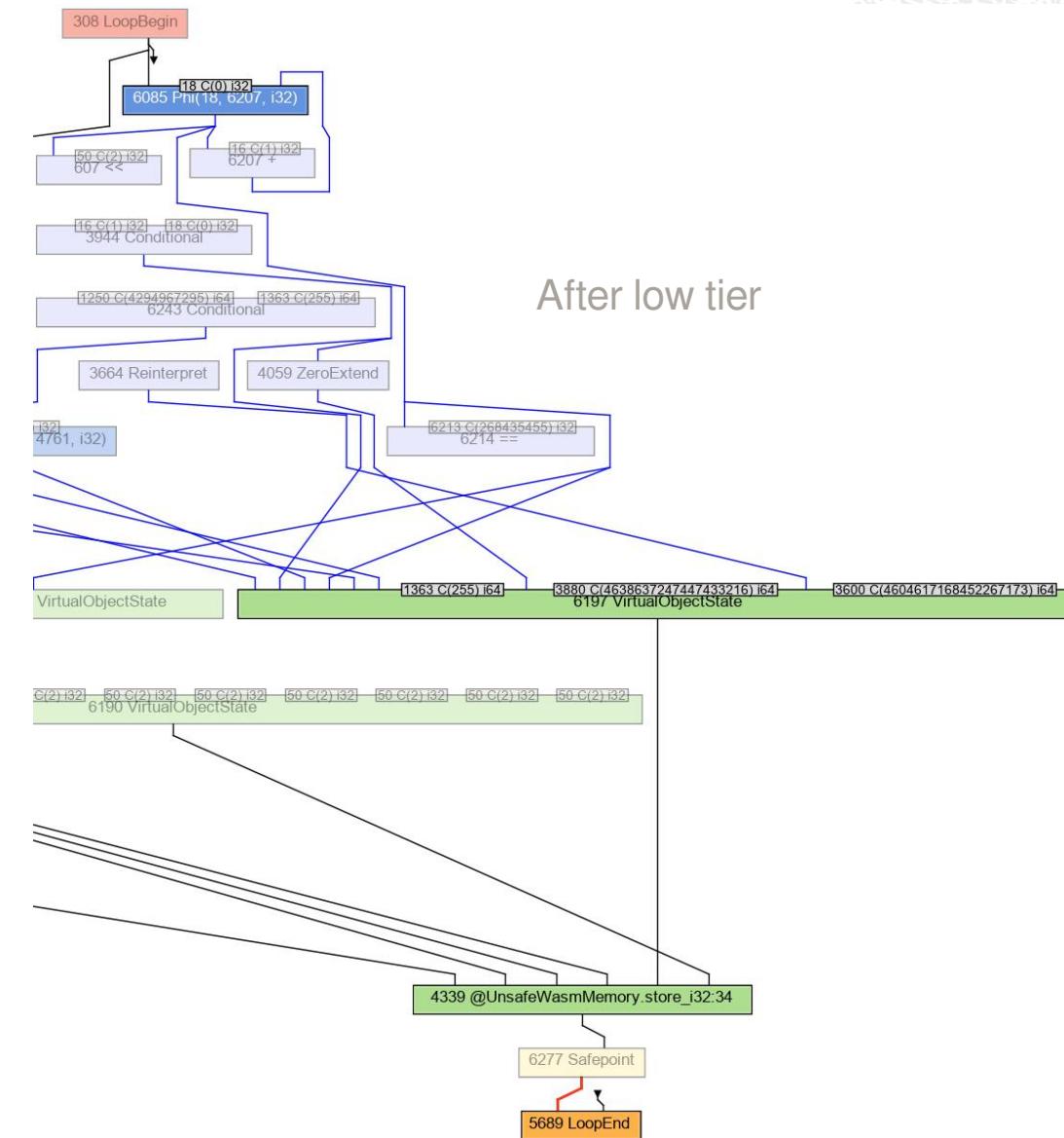


# Truffle frame-slot values

By inspecting the state of the frame, we can see what's tracked

We need to at least track these 2 values at loop end.

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = (color & 0xFF000000) >> 24;  
    uint8_t G = (color & 0x00FF0000) >> 16;  
    uint8_t B = (color & 0x0000FF00) >> 8;  
    double luminance =  
        (0.2126 * R + 0.7152 * G + 0.0722 * B);  
    result[pixel % IMAGE_SIZE] =  
        luminance > 127 ? UINT32_MAX : 0xFF;  
    black_pixels += luminance > 127 ? 0 : 1;  
}
```

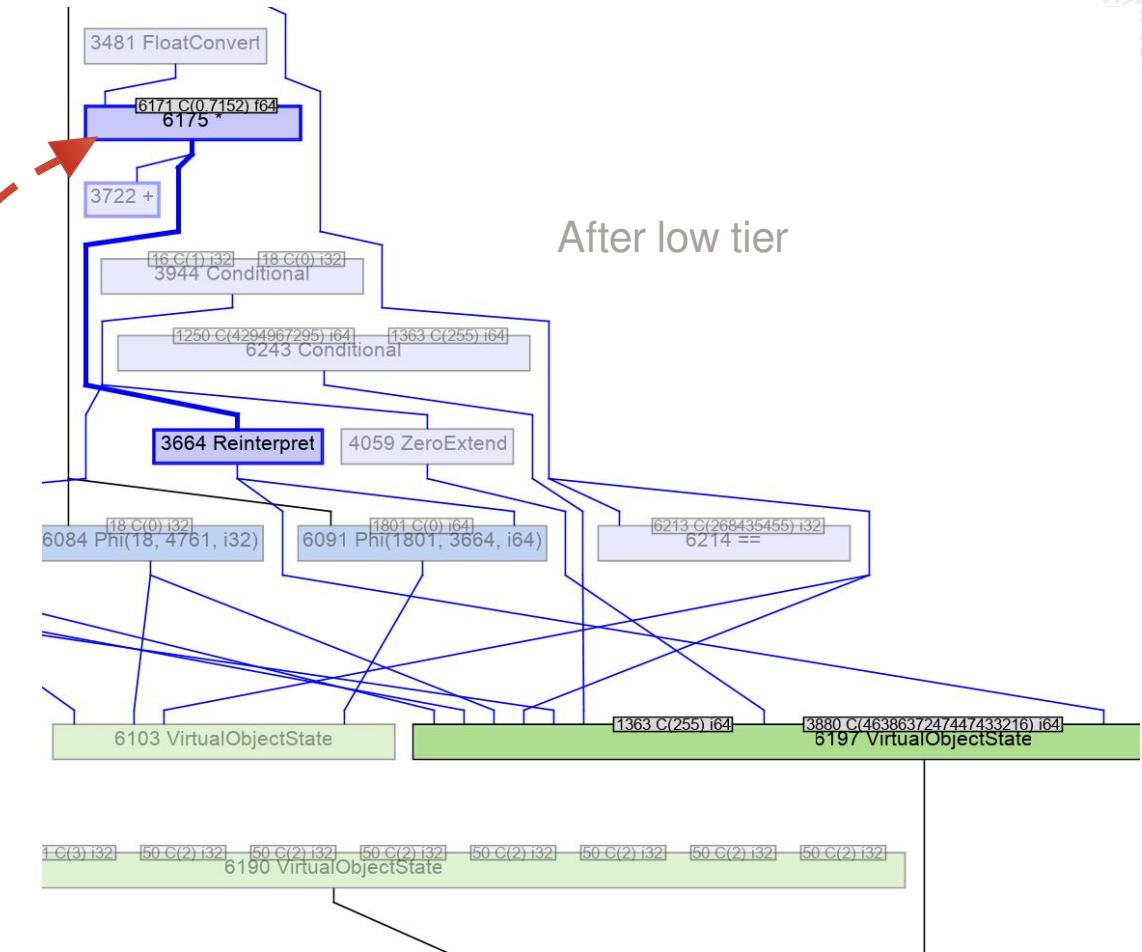


# Truffle frame-slot values

There are more values stored in the frame than what's necessary

What's this?

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = (color & 0xFF000000) >> 24;  
    uint8_t G = (color & 0x00FF0000) >> 16;  
    uint8_t B = (color & 0x0000FF00) >> 8;  
    double luminance =  
        (0.2126 * R + 0.7152 * G + 0.0722 * B);  
    result[pixel % IMAGE_SIZE] =  
        luminance > 127 ? UINT32_MAX : 0xFF;  
    black_pixels += luminance > 127 ? 0 : 1;  
}
```

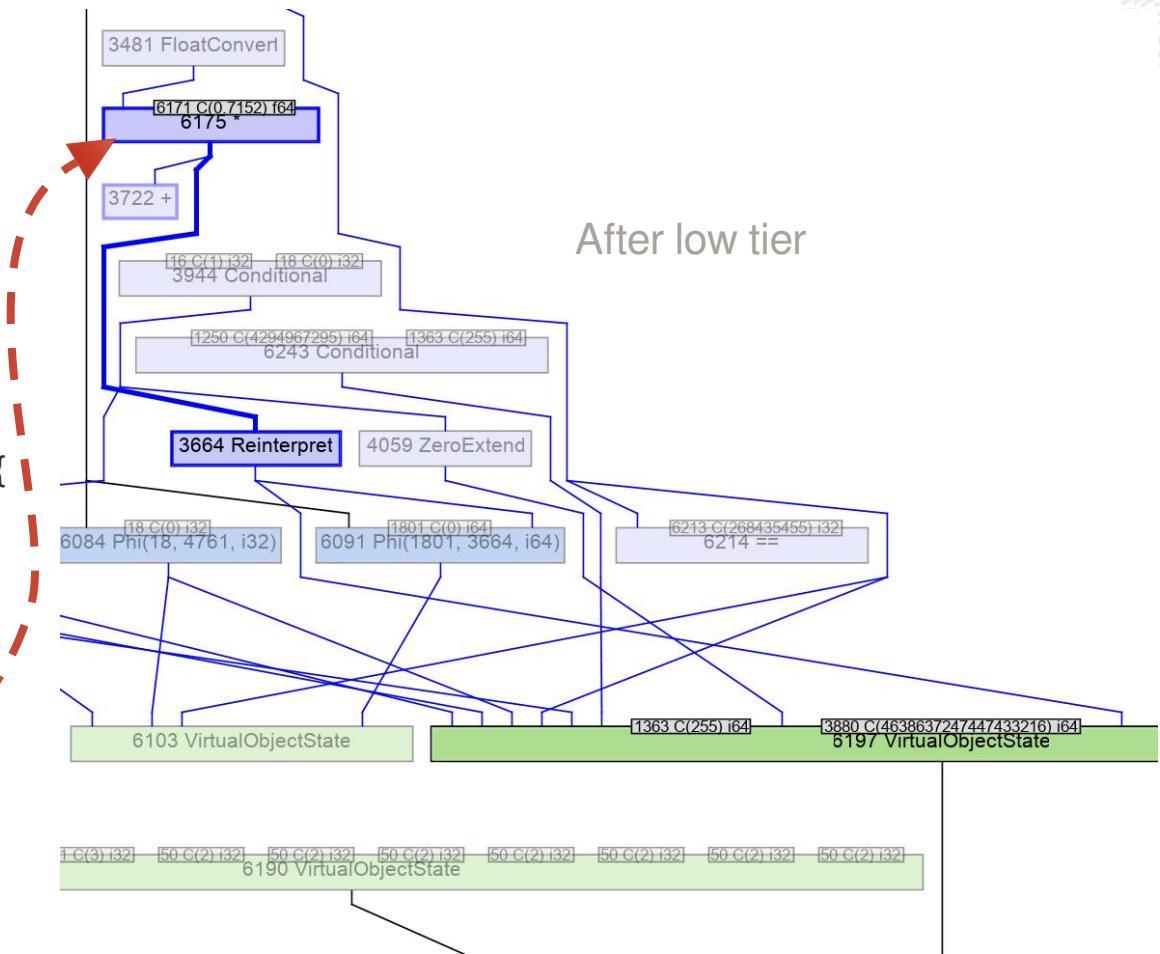


# Truffle frame-slot values

There are more values stored in the frame than what's necessary

No need to store the intermediate expression result.

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = (color & 0xFF000000) >> 24;  
    uint8_t G = (color & 0x00FF0000) >> 16;  
    uint8_t B = (color & 0x0000FF00) >> 8;  
    double luminance =  
        (0.2126 * R + 0.7152 * G + 0.0722 * B);  
    result[pixel % IMAGE_SIZE] =  
        luminance > 127 ? UINT32_MAX : 0xFF;  
    black_pixels += luminance > 127 ? 0 : 1;  
}
```

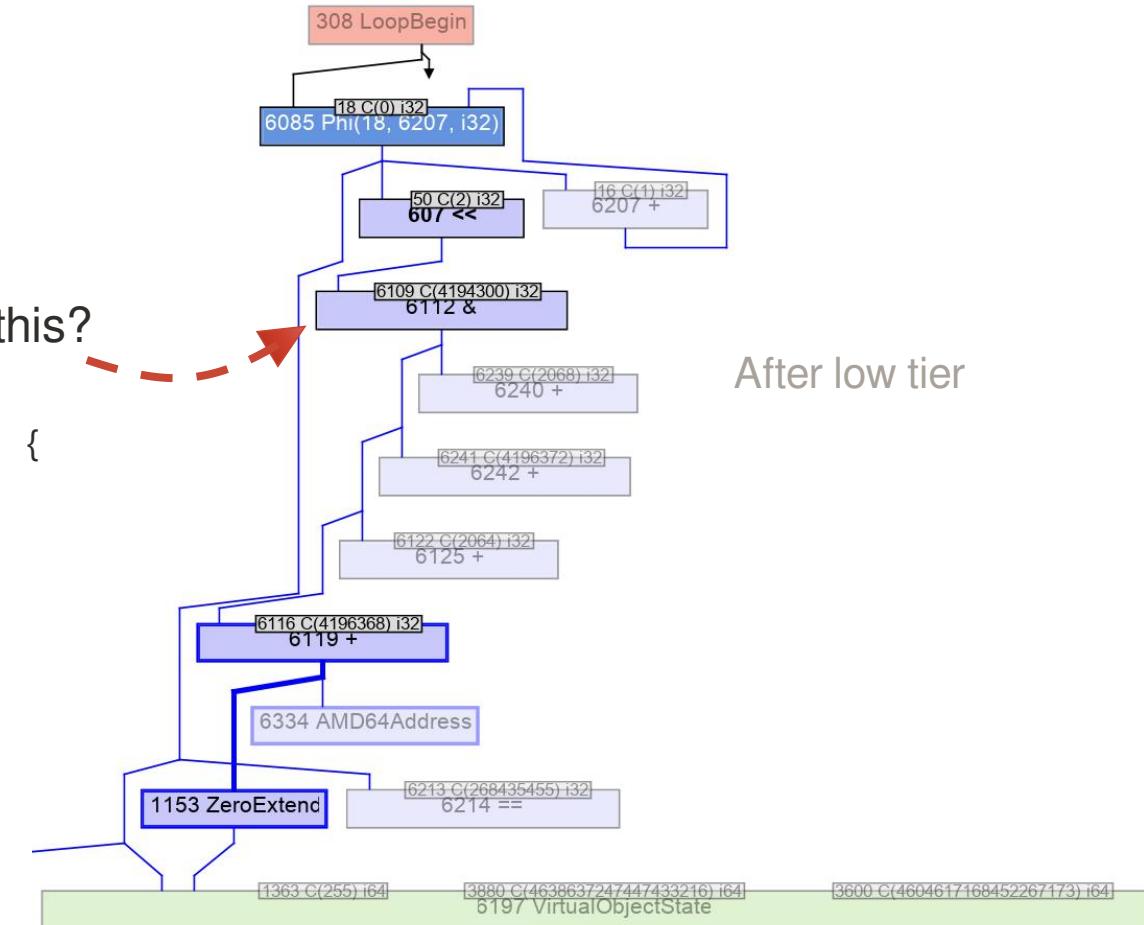


# Resetting Truffle frame-slot values

There are more values stored in the frame than what's necessary

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {  
    uint32_t color = image[pixel % IMAGE_SIZE];  
    uint8_t R = (color & 0xFF000000) >> 24;  
    uint8_t G = (color & 0x00FF0000) >> 16;  
    uint8_t B = (color & 0x0000FF00) >> 8;  
    double luminance =  
        (0.2126 * R + 0.7152 * G + 0.0722 * B);  
    result[pixel % IMAGE_SIZE] =  
        luminance > 127 ? UINT32_MAX : 0xFF;  
    black_pixels += luminance > 127 ? 0 : 1;  
}
```

What's this?

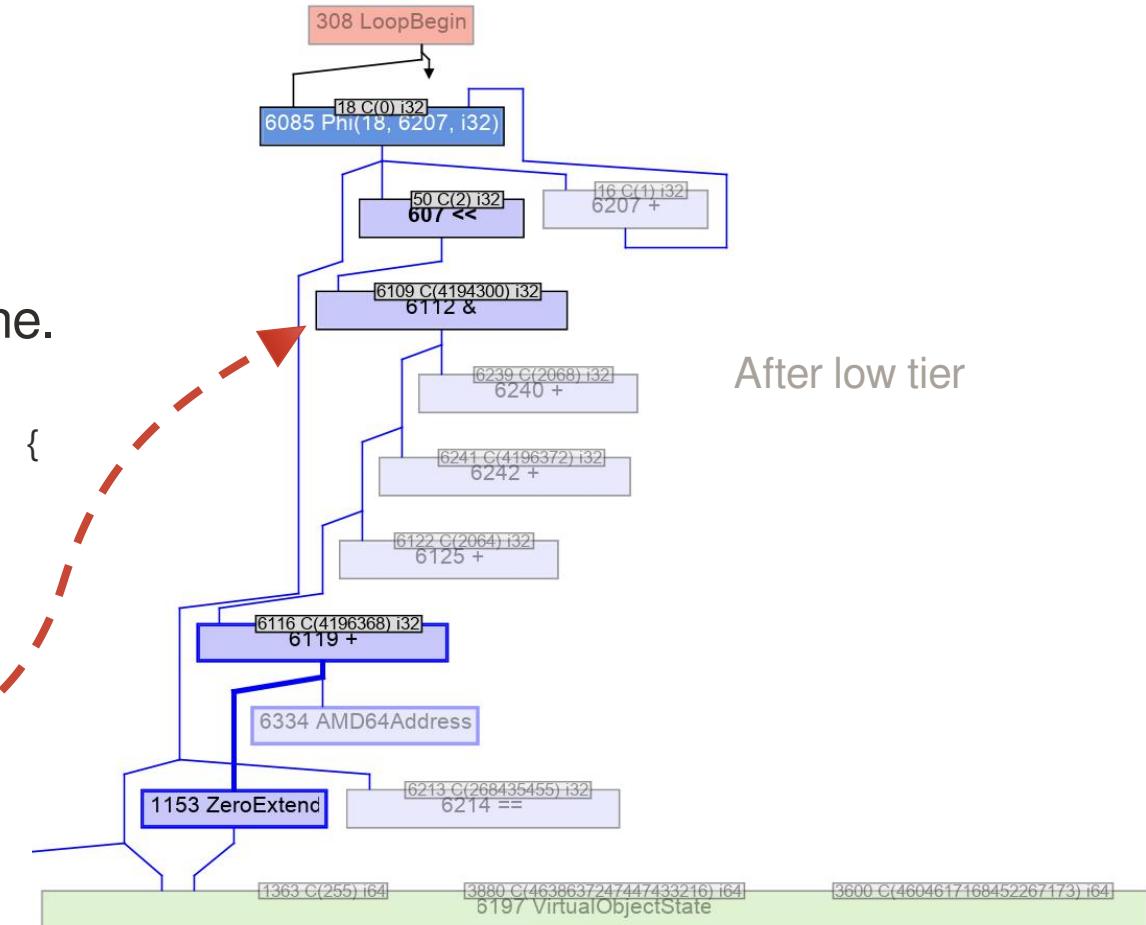


# Resetting Truffle frame-slot values

There are more values stored in the frame than what's necessary

No need to store the array offset after the write is done.

```
for (uint32_t pixel = 0; pixel != IMAGE_SIZE; ++pixel) {
    uint32_t color = image[pixel % IMAGE_SIZE];
    uint8_t R = (color & 0xFF000000) >> 24;
    uint8_t G = (color & 0x00FF0000) >> 16;
    uint8_t B = (color & 0x0000FF00) >> 8;
    double luminance =
        (0.2126 * R + 0.7152 * G + 0.0722 * B);
    result[pixel % IMAGE_SIZE] =
        luminance > 127 ? UINT32_MAX : 0xFF;
    black_pixels += luminance > 127 ? 0 : 1;
}
```



# Resetting Truffle frame-slot values

## Why do these extra values persist?

To understand why, we need to look at the WebAssembly code that the C program was compiled to.

```
loop ; label = @1
local.get 2
i32.const 2
i32.shl
i32.const 4194300
i32.and
local.tee 0
i32.const 4196368
i32.add
i32.const -1
i32.const 255
local.get 0
i32.const 2064
i32.add
i32.load
local.tee 0
i32.const 8
i32.shr_u
i32.const 255
i32.and
f64.convert_i32_s
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722;)
f64.mul
local.get 0
i32.const 24
i32.shr_u
f64.convert_i32_s
f64.const 0x1.b367a0f9096bc9-3 (:=0.2126;)
f64.mul
local.get 0
i32.const 16
i32.shr_u
i32.const 255
i32.and
f64.convert_i32_s
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)
f64.mul
f64.add
f64.add
f64.const 0x1.fcp+6 (:=127;)
f64.gt
local.tee 0
select
i32.store
local.get 1
local.get 0
i32.const 1
i32.xor
i32.add
local.set 1
local.get 2
i32.const 1
i32.add
local.tee 2
i32.const 268435456
i32.ne
br_if 0 (@1;)
end
```

# Resetting Truffle frame-slot values

Part of the guest language Frame is used as an expression stack

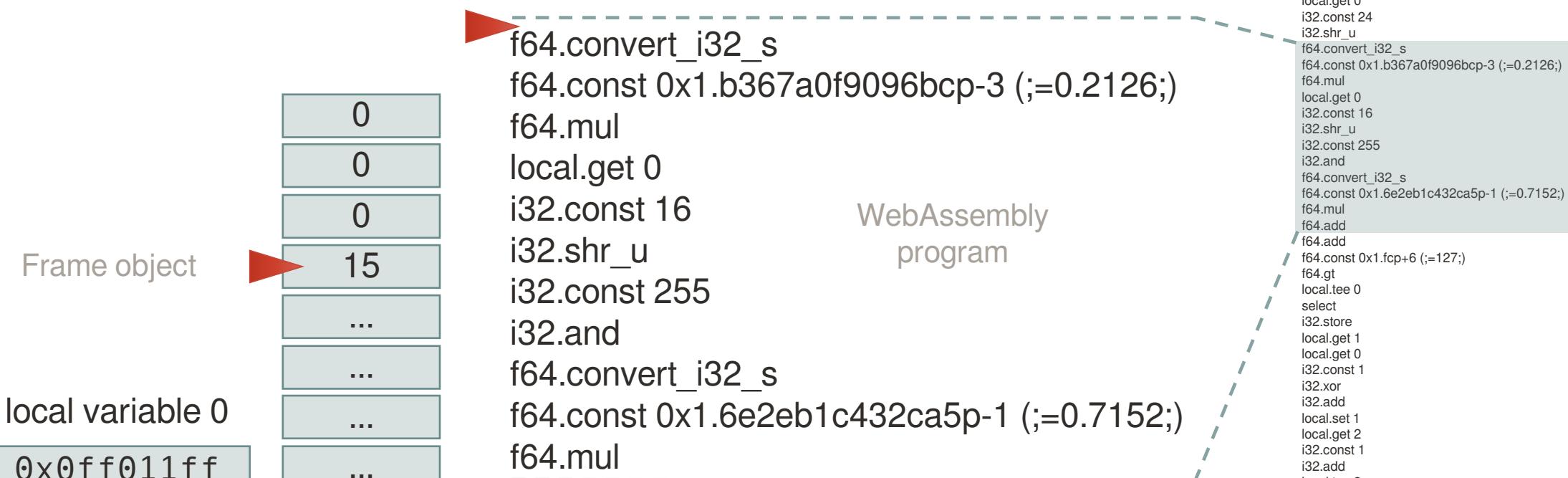
▶ f64.convert\_i32\_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126;  
f64.mul  
local.get 0  
i32.const 16  
i32.shr\_u  
i32.const 255  
i32.and  
f64.convert\_i32\_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;  
f64.mul

WebAssembly  
program

```
loop ; label = @1
local.get 2
i32.const 2
i32.shl
i32.const 4194300
i32.and
local.tee 0
i32.const 4196368
i32.add
i32.const -1
i32.const 255
local.get 0
i32.const 2064
i32.add
i32.load
local.tee 0
i32.const 8
i32.shr_u
i32.const 255
i32.and
f64.convert_i32_s
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722;
f64.mul
local.get 0
i32.const 24
i32.shr_u
f64.convert_i32_s
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126;
f64.mul
local.get 0
i32.const 16
i32.shr_u
i32.const 255
i32.and
f64.convert_i32_s
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;
f64.mul
f64.add
f64.add
f64.const 0x1.fcp+6 (:=127;
f64.gt
local.tee 0
select
i32.store
local.get 1
local.get 0
i32.const 1
i32.xor
i32.add
local.set 1
local.get 2
i32.const 1
i32.add
local.tee 2
i32.const 268435456
i32.ne
br_if 0 (@1;
end
```

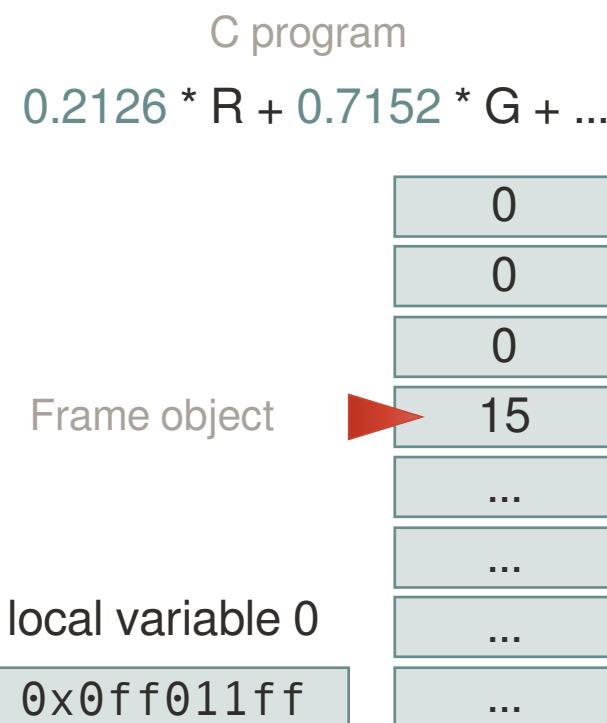
# Resetting Truffle frame-slot values

Part of the guest language Frame is used as an expression stack



```
loop ;; label = @1
local.get 2
i32.const 2
i32.shl
i32.const 4194300
i32.and
local.tee 0
i32.const 4196368
i32.add
i32.const -1
i32.const 255
local.get 0
i32.const 2064
i32.add
i32.load
local.tee 0
i32.const 8
i32.shr_u
i32.const 255
i32.and
f64.convert_i32_s
f64.const 0x1.27bb2fec56d5dp-4 ( $=0.0722;$ )
f64.mul
local.get 0
i32.const 24
i32.shr_u
f64.convert_i32_s
f64.const 0x1.b367a0f9096bcp-3 ( $=0.2126;$ )
f64.mul
local.get 0
i32.const 16
i32.shr_u
i32.const 255
i32.and
f64.convert_i32_s
f64.const 0x1.6e2eb1c432ca5p-1 ( $=0.7152;$ )
f64.mul
f64.add
f64.add
f64.const 0x1.fcp+6 ( $=127;$ )
f64.gt
local.tee 0
select
i32.store
local.get 1
local.get 0
i32.const 1
i32.xor
i32.add
local.set 1
local.get 2
i32.const 1
i32.add
local.tee 2
i32.const 268435456
i32.ne
br_if 0 (@1)
end
```

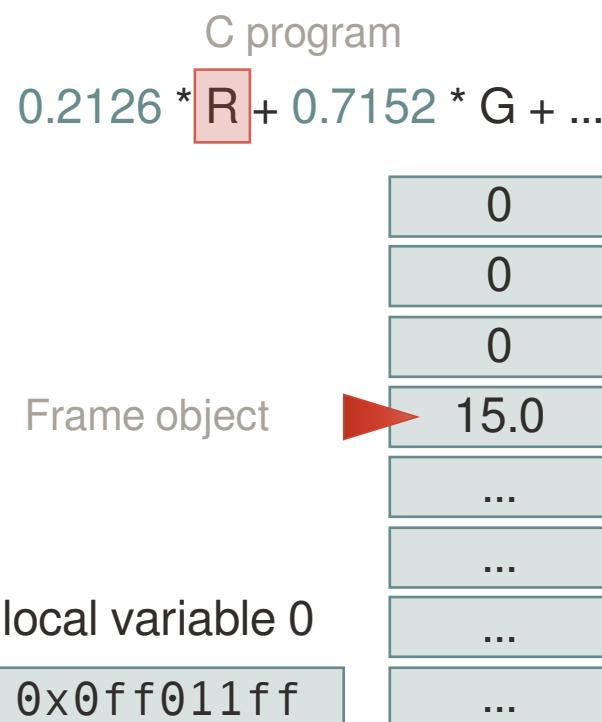
# Resetting Truffle frame-slot values



WebAssembly program

```
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc... (=0.2126);  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (=0.7152);  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (=127);  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1); end
```

# Resetting Truffle frame-slot values

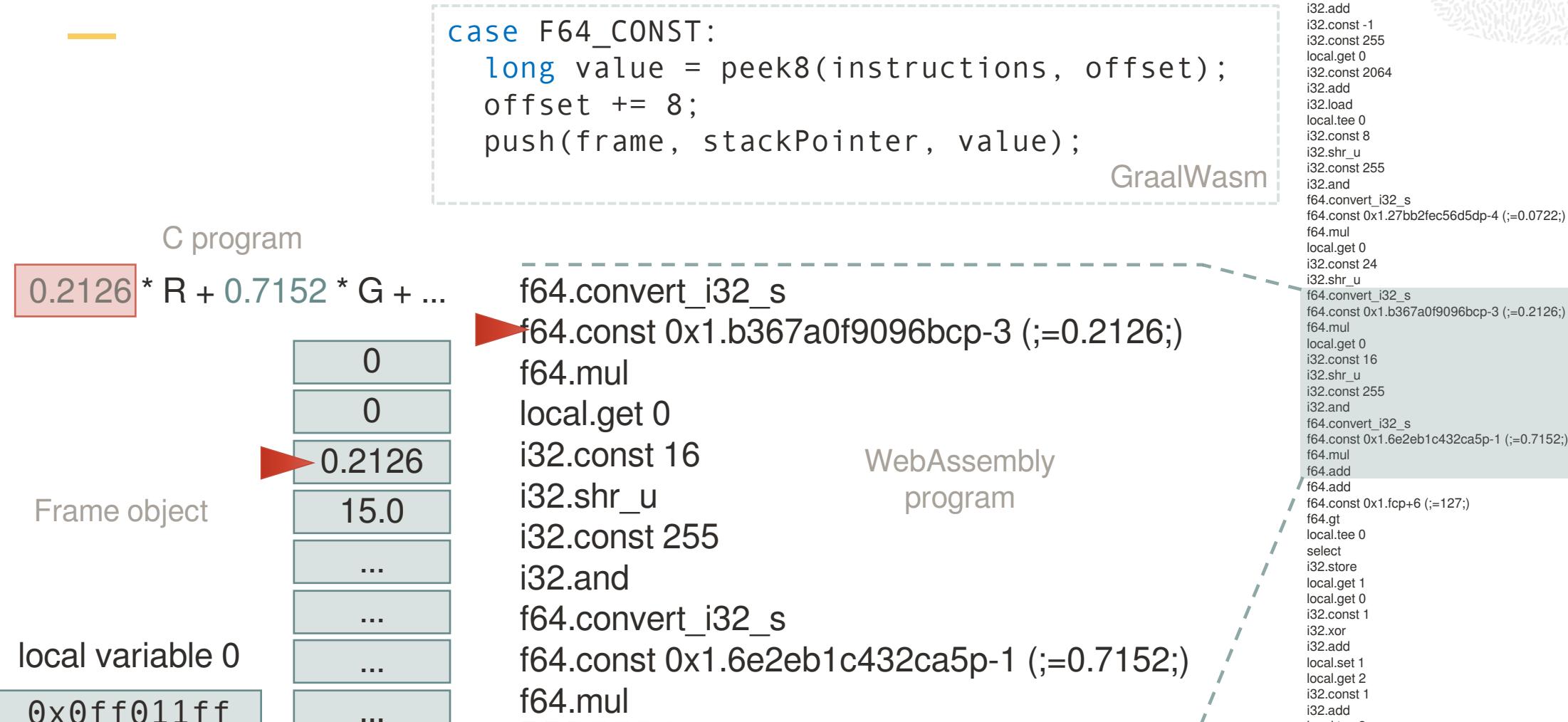


```
case F64_CONVERT_I64_S:           GraalWasm
case F64_CONVERT_I64_U:
    long x = pop(frame, stackPointer);
    double result = (double) x;
    pushDouble(frame, stackPointer, result);
```

f64.convert\_i32\_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr\_u  
i32.const 255  
i32.and  
f64.convert\_i32\_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127;)  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br\_if 0 (@1;)  
end

WebAssembly  
program

# Resetting Truffle frame-slot values



# Resetting Truffle frame-slot values

—

C program

0.2126 \* R + 0.7152 \* G + ...

Frame object

local variable 0

0x0ff011fff

```
case F64_MUL:  
    double x = popDouble(frame, stackPointer);  
    double y = popDouble(frame, stackPointer);  
    double result = y * x;  
    pushDouble(frame, stackPointer, result);
```

f64.convert\_i32\_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr\_u  
i32.const 255  
i32.and  
f64.convert\_i32\_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul

WebAssembly  
program

```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722;)  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127;)  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1;)  
end
```

# Resetting Truffle frame-slot values



# Resetting Truffle frame-slot values

C program

```
G = (color & 0x00FF0000) >> 16;
```

```
0.2126 * R + 0.7152 * G + ...
```

Frame object

local variable 0

```
0x0ff011ff
```

```
case I32_CONST:  
    int value = peek4(instructions, offset);  
    offset += 4;  
    push(frame, stackPointer, value);
```

GraalWasm

```
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);
```

f64.mul

local.get 0

```
i32.const 16
```

i32.shr\_u

```
i32.const 255
```

i32.and

```
f64.convert_i32_s
```

```
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);
```

```
f64.mul
```

WebAssembly  
program

```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722);  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127);  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1);  
end
```

# Resetting Truffle frame-slot values

C program

```
G = (color & 0x00FF0000) >> 16;  
0.2126 * R + 0.7152 * G + ...
```

Frame object

local variable 0

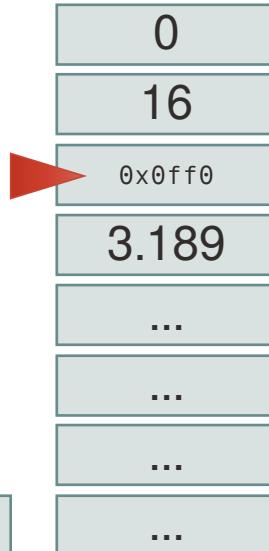
0x0ff011fff

case I32\_SHR\_U:  
int x = popInt(frame, stackPointer);  
int y = popInt(frame, stackPointer);  
int result = y >>> x;  
pushInt(frame, stackPointer, result);

GraalWasm

f64.convert\_i32\_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);  
f64.mul  
local.get 0  
i32.const 16  
i32.shr\_u  
i32.const 255  
i32.and  
f64.convert\_i32\_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127);  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br\_if 0 (@1);  
end

WebAssembly  
program



# Resetting Truffle frame-slot values

C program

```
G = (color & 0x00FF0000) >> 16;
```

```
0.2126 * R + 0.7152 * G + ...
```

Frame object

local variable 0

```
0x0ff011ff
```

```
case I32_CONST:  
    int value = peek4(instructions, offset);  
    offset += 4;  
    push(frame, stackPointer, value);
```

GraalWasm

```
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);
```

f64.mul

local.get 0

i32.const 16

i32.shr\_u

i32.const 255

i32.and

```
f64.convert_i32_s
```

```
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);
```

f64.mul

WebAssembly  
program

```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722);  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127);  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1);  
end
```

# Resetting Truffle frame-slot values

C program

```
G = (color & 0x00FF0000) >> 16;  
0.2126 * R + 0.7152 * G + ...
```

Frame object

local variable 0

0x0ff011ff

```
case I32_AND:  
    int x = popInt(frame, stackPointer);  
    int y = popInt(frame, stackPointer);  
    int result = y & x;  
    pushInt(frame, stackPointer, result);
```

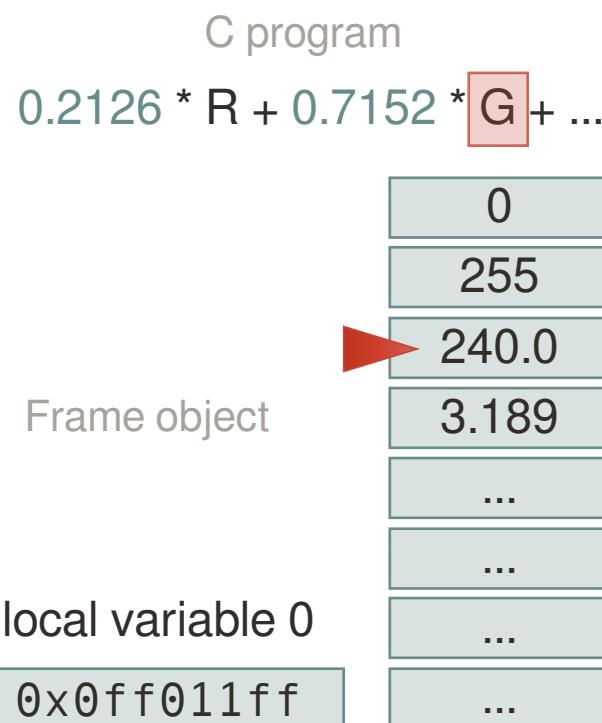
GraalWasm

```
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);  
f64.mul
```

WebAssembly  
program

```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722);  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bc-p-3 (:=0.2126);  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152);  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127);  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1);  
end
```

# Resetting Truffle frame-slot values



GraalWasm

```
case F64_CONVERT_I64_S:  
case F64_CONVERT_I64_U:  
    long x = pop(frame, stackPointer);  
    double result = x;  
    pushDouble(frame, stackPointer, result);
```

WebAssembly  
program

```
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul
```

```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722;)  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127;)  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1;)  
end
```

# Resetting Truffle frame-slot values

	C program							
	$0.2126 * R + 0.7152 * G + \dots$							
Frame object	<table><tr><td>0</td></tr><tr><td>0.7152</td></tr><tr><td>240.0</td></tr><tr><td>3.189</td></tr><tr><td>...</td></tr><tr><td>...</td></tr><tr><td>...</td></tr></table>	0	0.7152	240.0	3.189	...	...	...
0								
0.7152								
240.0								
3.189								
...								
...								
...								
local variable 0	<table><tr><td>0x0ff011fff</td></tr><tr><td>...</td></tr></table>	0x0ff011fff	...					
0x0ff011fff								
...								

```
case F64_CONST:  
    long value = peek8(instructions, offset);  
    offset += 8;  
    push(frame, stackPointer, value);
```

GraalWasm

```
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul
```

WebAssembly  
program

```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 (:=0.0722;)  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bcp-3 (:=0.2126;)  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 (:=0.7152;)  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 (:=127;)  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1;)  
end
```

# Resetting Truffle frame-slot values

—

C program

$0.2126 * R + 0.7152 * G + \dots$

Frame object

local variable 0

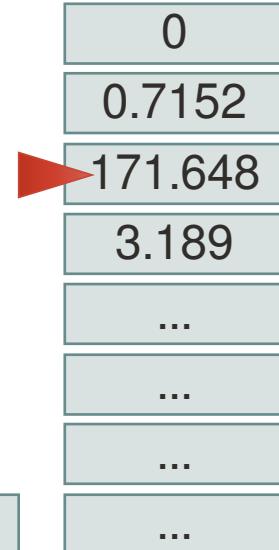
`0x0ff011ff`

```
case F64_MUL:  
    double x = popDouble(frame, stackPointer);  
    double y = popDouble(frame, stackPointer);  
    double result = y * x;  
    pushDouble(frame, stackPointer, result);
```

GraalWasm

f64.convert\_i32\_s  
f64.const 0x1.b367a0f9096bcp-3 ( $=0.2126;$ )  
f64.mul  
local.get 0  
i32.const 16  
i32.shr\_u  
i32.const 255  
i32.and  
f64.convert\_i32\_s  
f64.const 0x1.6e2eb1c432ca5p-1 ( $=0.7152;$ )  
f64.mul  
f64.add  
f64.add

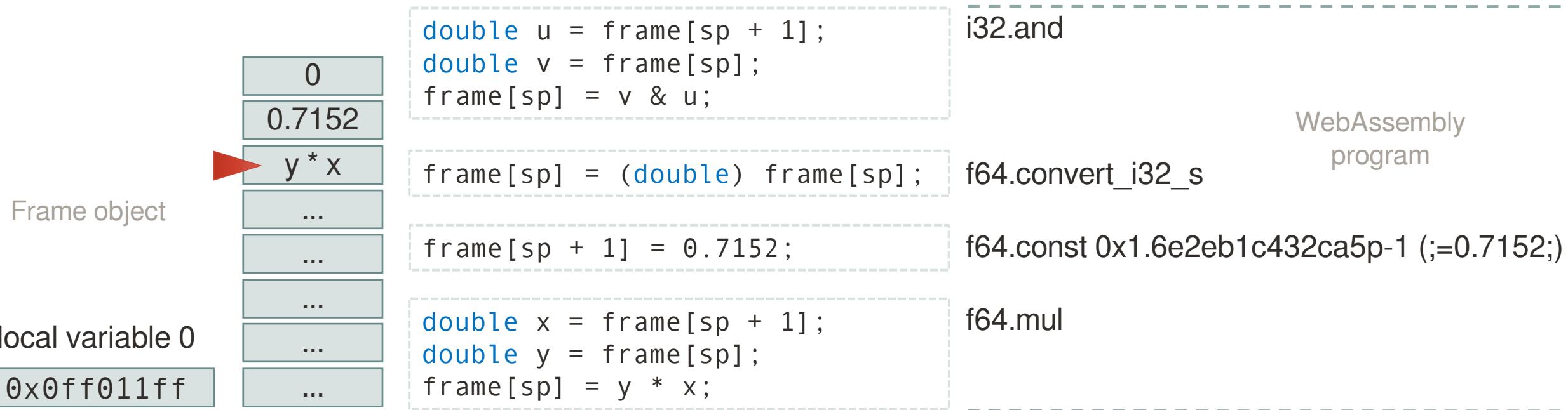
WebAssembly  
program



```
loop ; label = @1  
local.get 2  
i32.const 2  
i32.shl  
i32.const 4194300  
i32.and  
local.tee 0  
i32.const 4196368  
i32.add  
i32.const -1  
i32.const 255  
local.get 0  
i32.const 2064  
i32.add  
i32.load  
local.tee 0  
i32.const 8  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.27bb2fec56d5dp-4 ( $=0.0722;$ )  
f64.mul  
local.get 0  
i32.const 24  
i32.shr_u  
f64.convert_i32_s  
f64.const 0x1.b367a0f9096bcp-3 ( $=0.2126;$ )  
f64.mul  
local.get 0  
i32.const 16  
i32.shr_u  
i32.const 255  
i32.and  
f64.convert_i32_s  
f64.const 0x1.6e2eb1c432ca5p-1 ( $=0.7152;$ )  
f64.mul  
f64.add  
f64.add  
f64.const 0x1.fcp+6 ( $=127;$ )  
f64.gt  
local.tee 0  
select  
i32.store  
local.get 1  
local.get 0  
i32.const 1  
i32.xor  
i32.add  
local.set 1  
local.get 2  
i32.const 1  
i32.add  
local.tee 2  
i32.const 268435456  
i32.ne  
br_if 0 (@1;)  
end
```

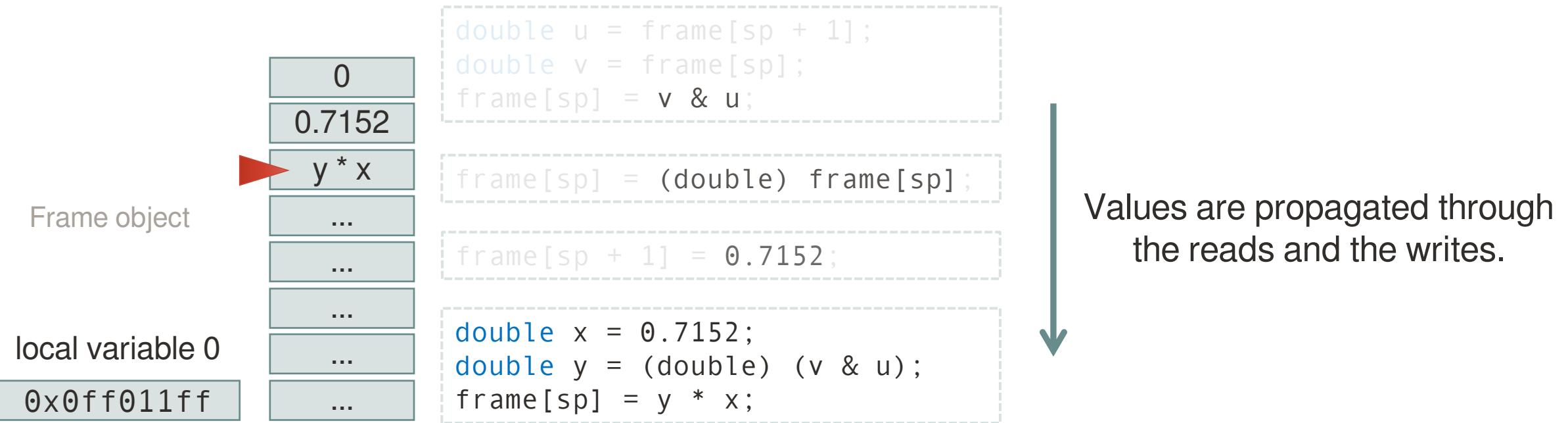
# Resetting Truffle frame-slot values

The partial evaluator stitches these interpreter snippets together.



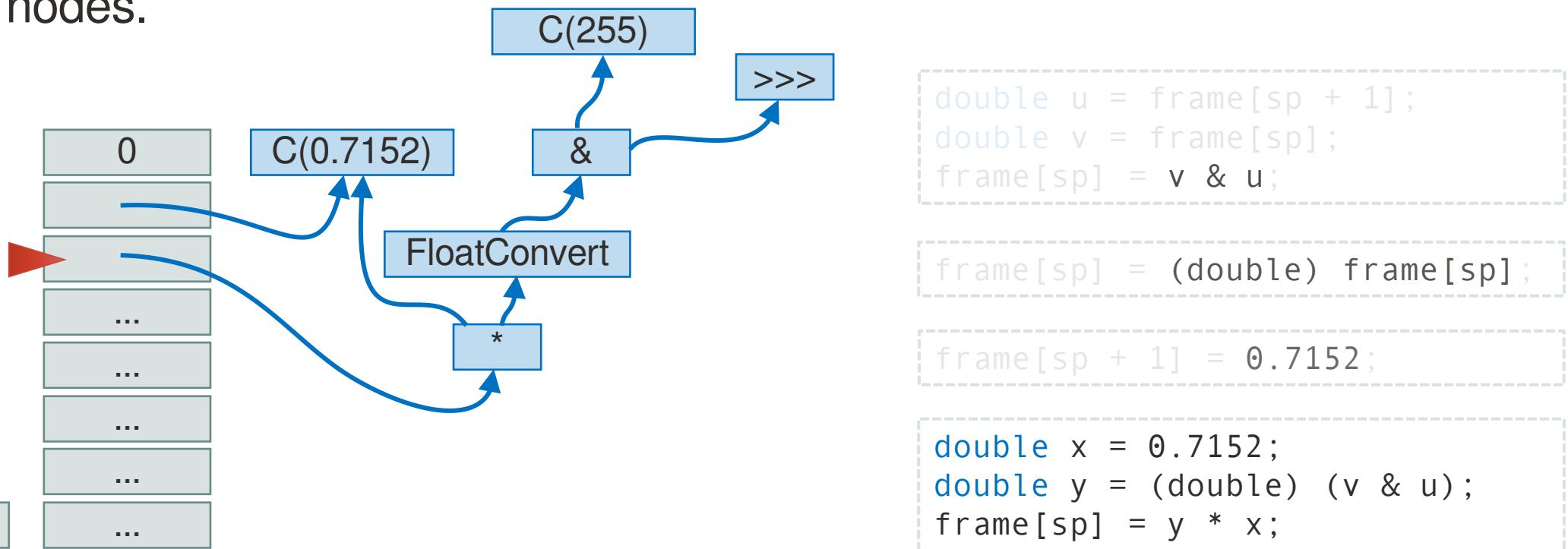
# Resetting Truffle frame-slot values

Reads and writes on the frame slots are eliminated during escape analysis.



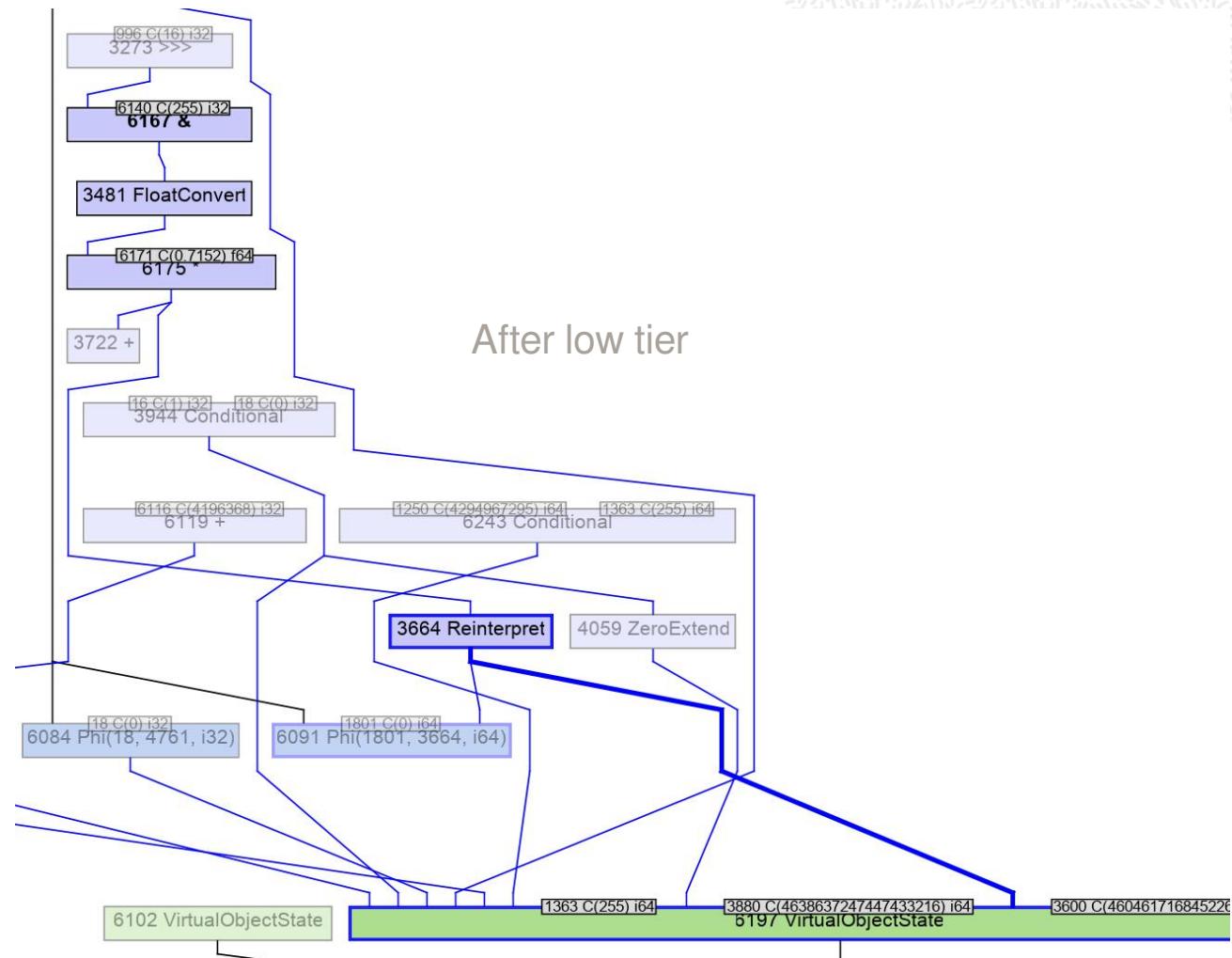
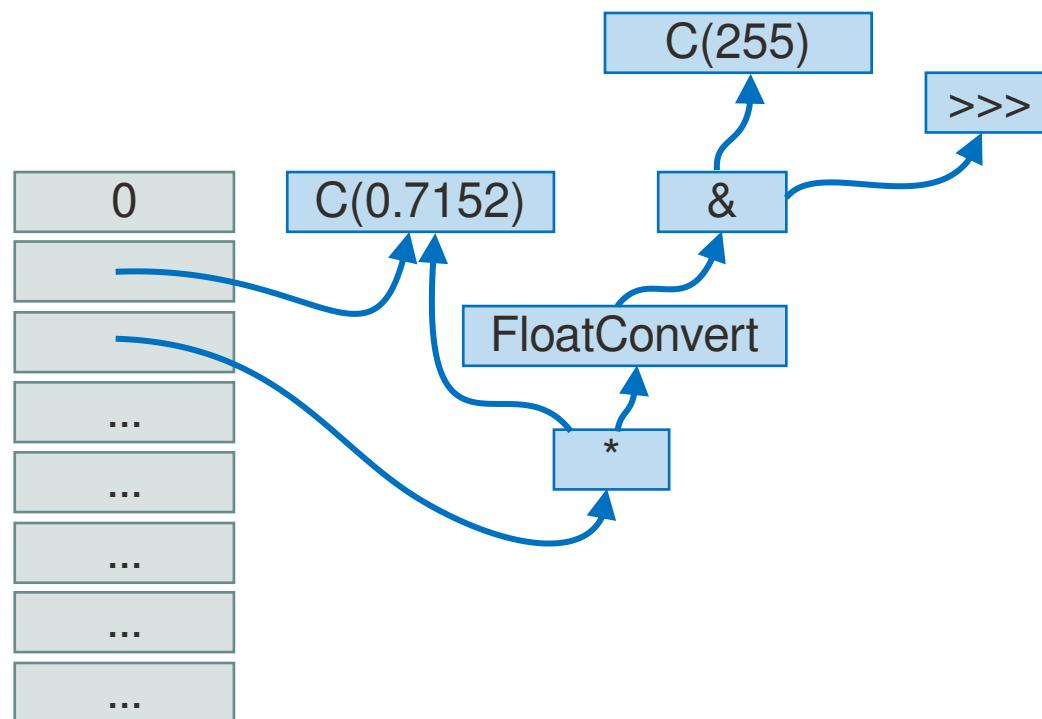
# Resetting Truffle frame-slot values

From the compiler's standpoint, the contents of the Frame object are various expression nodes.



# Resetting Truffle frame-slot values

This is exactly what we see in the IR.

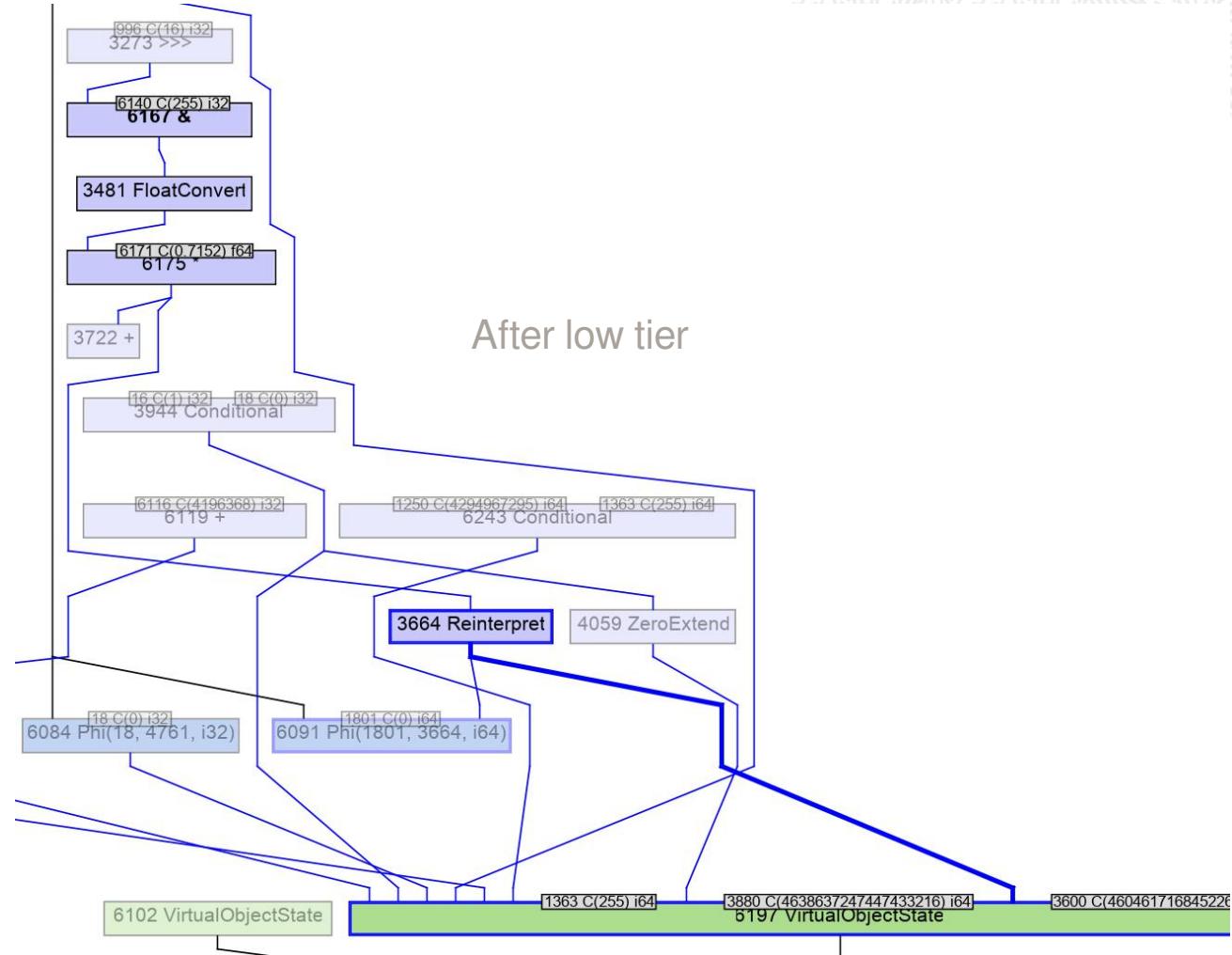


# Resetting Truffle frame-slot values

## How do we prevent this?

Set the Frame slots to their default values after the value is no longer needed, because the virtual object state does not need to track default values.

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long pop(VirtualFrame frame, int slot) {  
    long result = frame.getLong(slot);  
    return result;  
}
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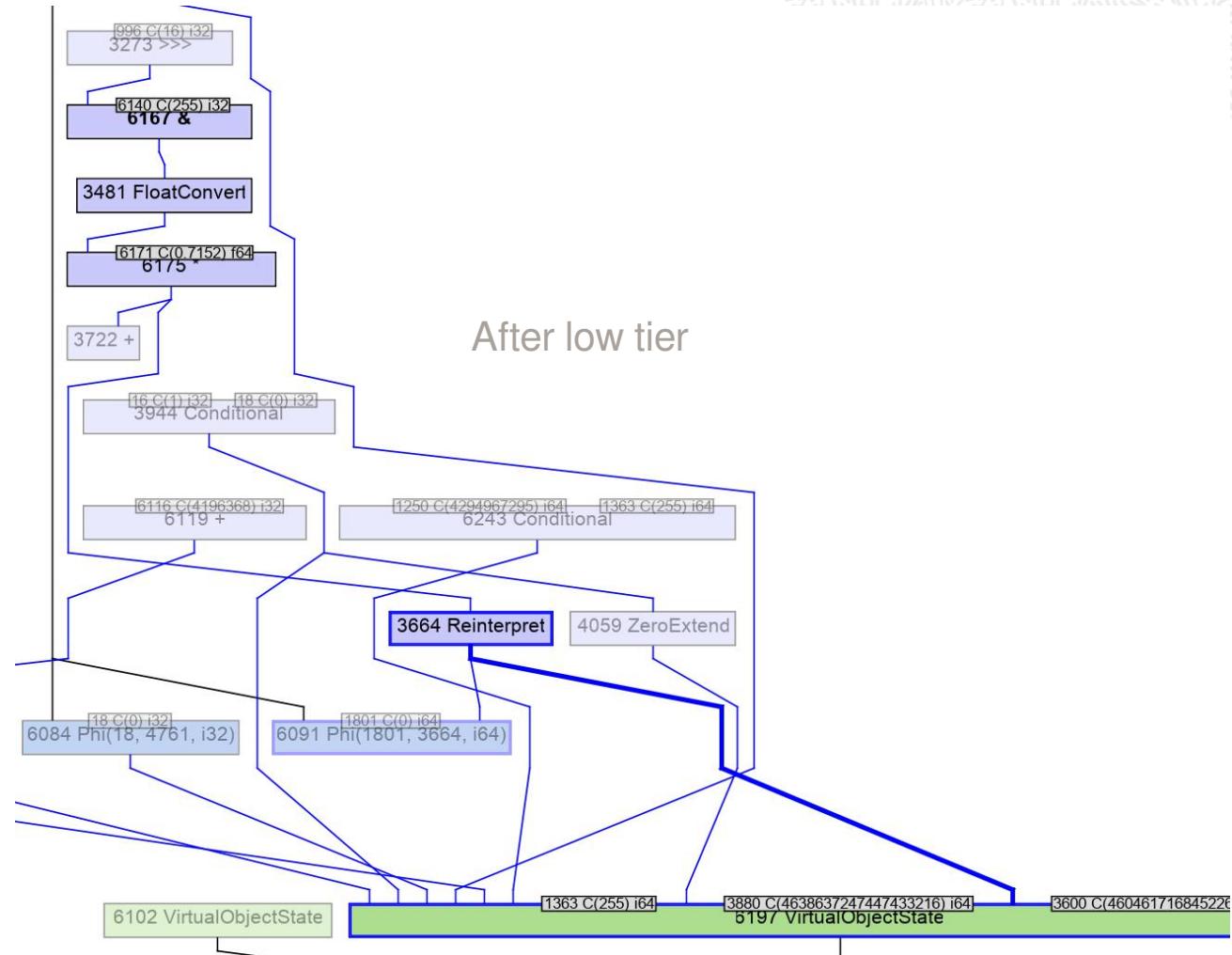


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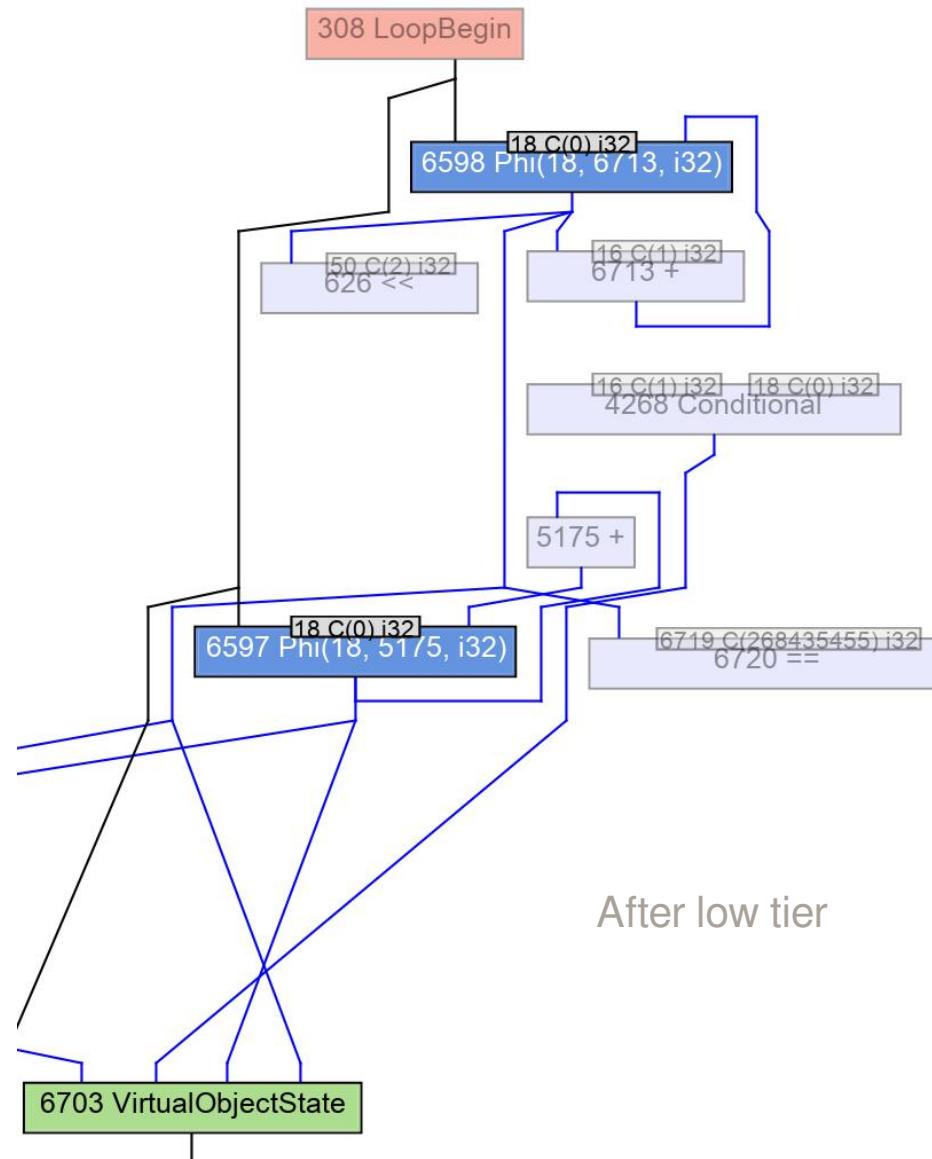


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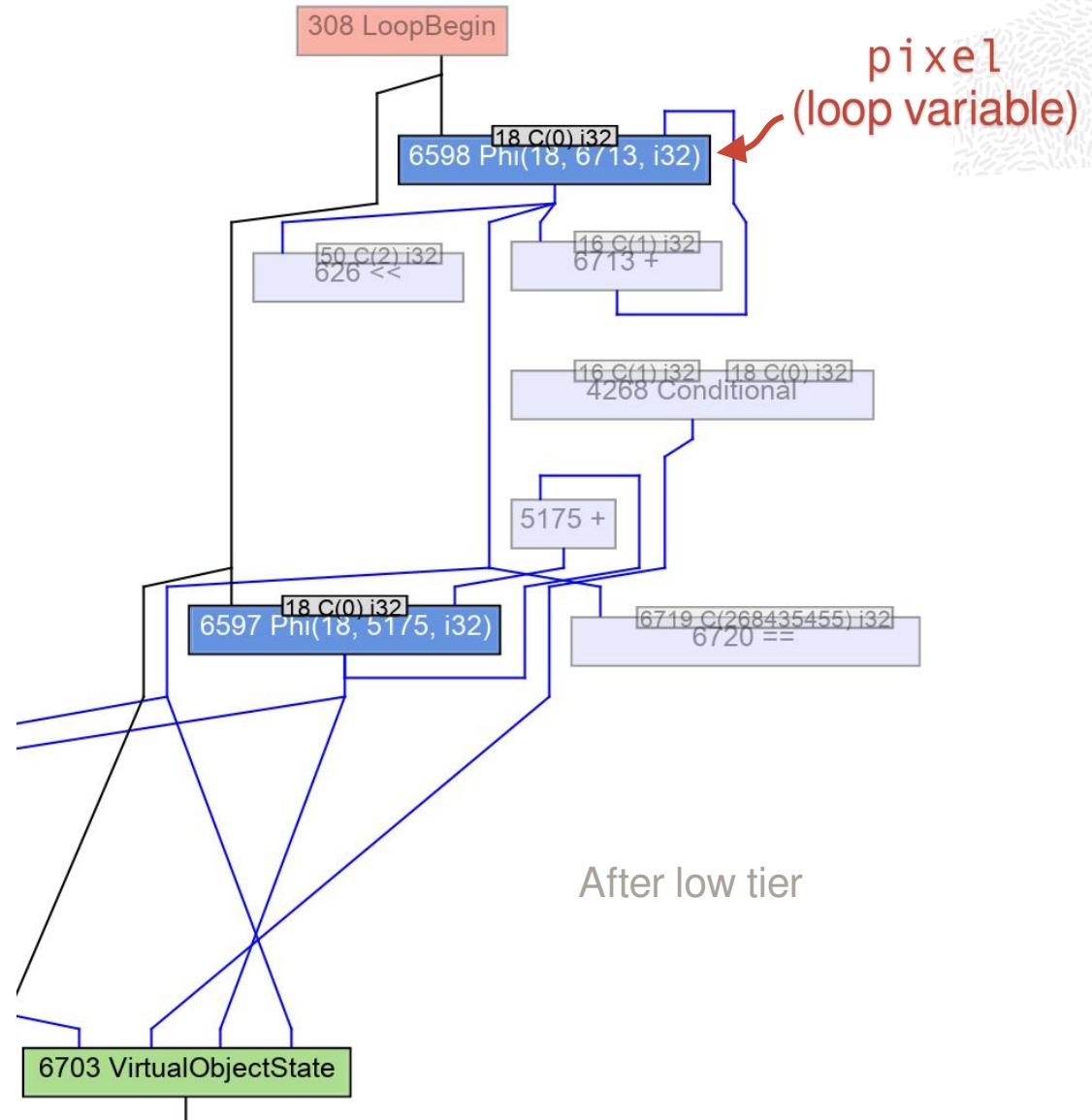


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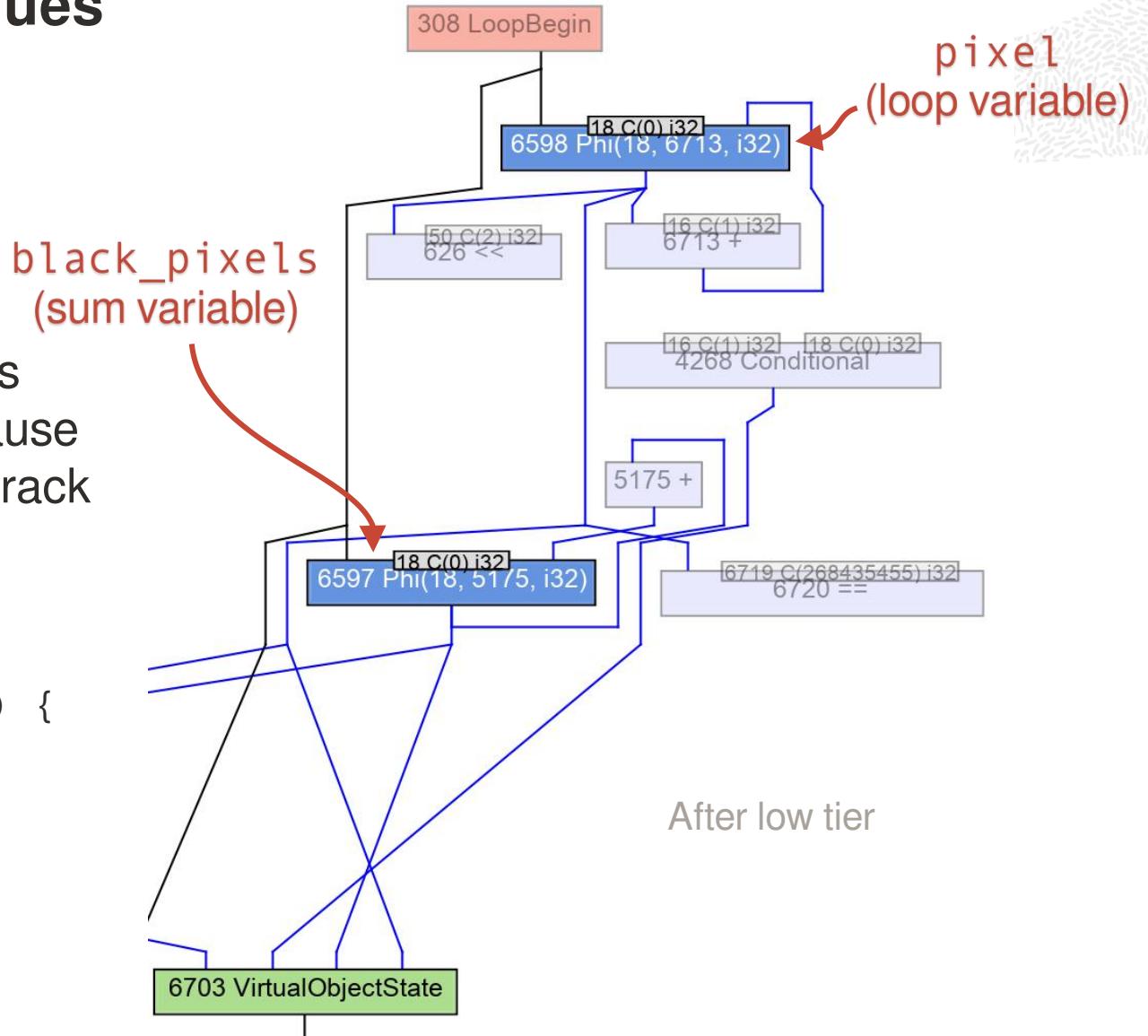


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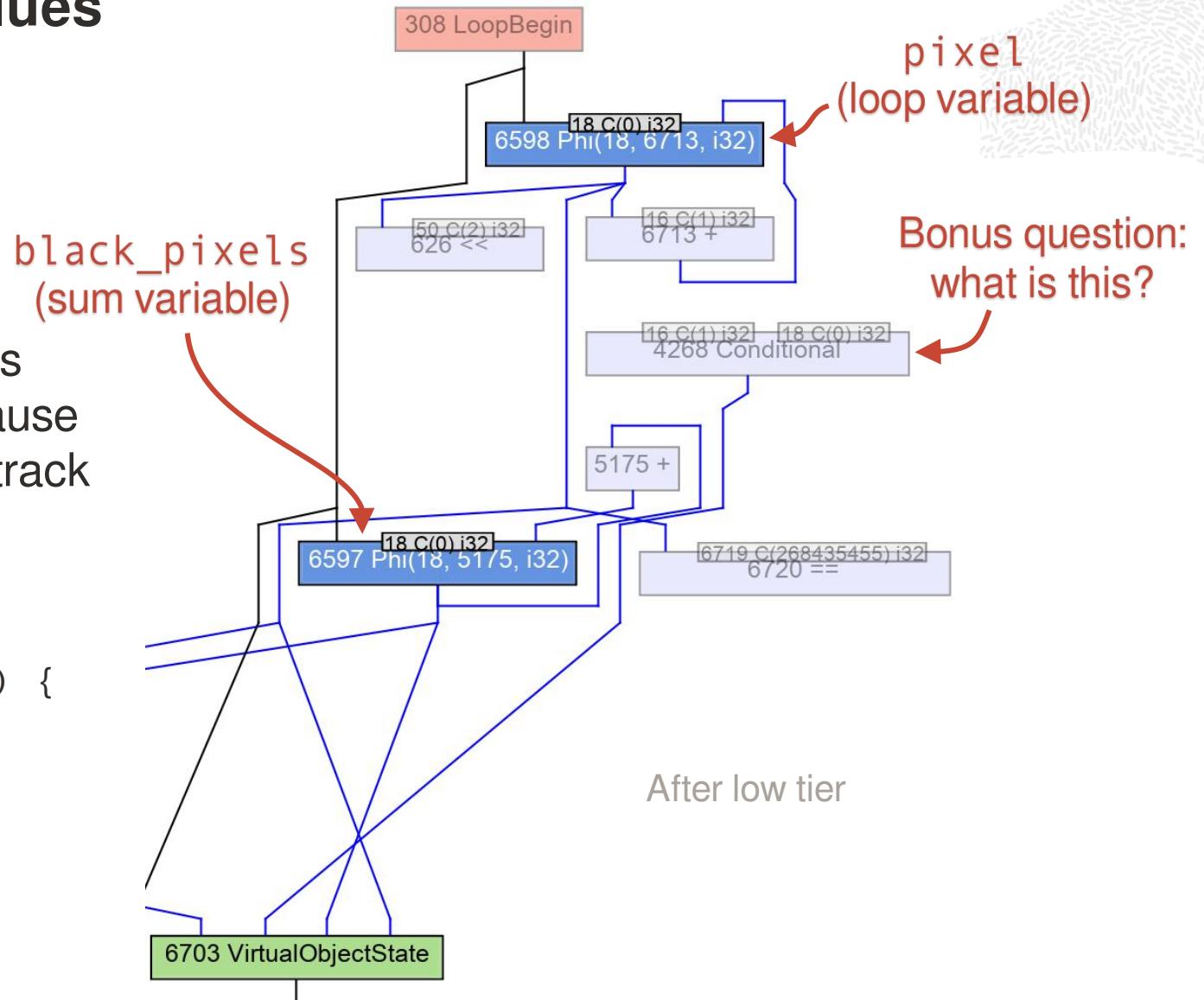


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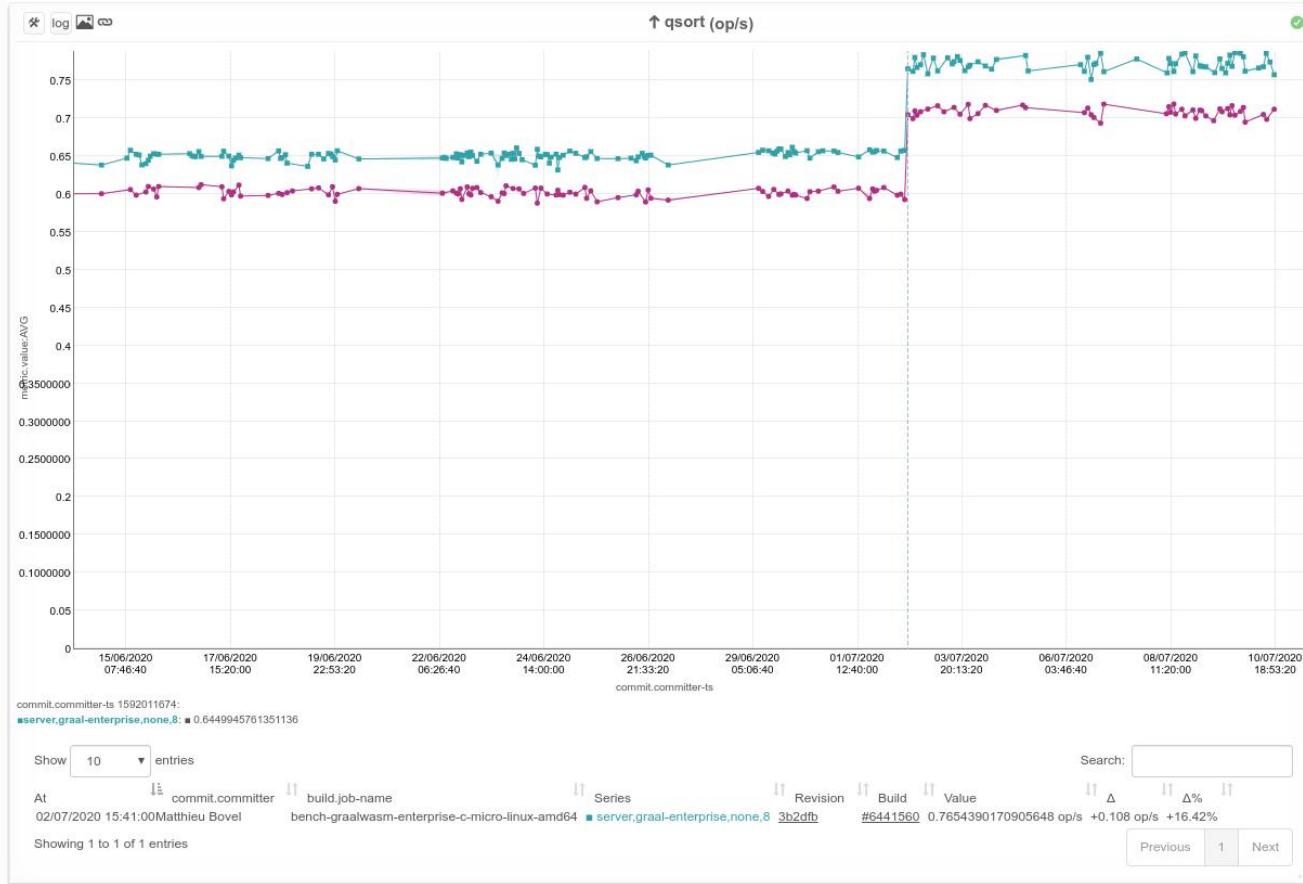


# Resetting Truffle frame-slot values



- 21% improvement on a discrete-event simulation benchmark (*event-sim*)

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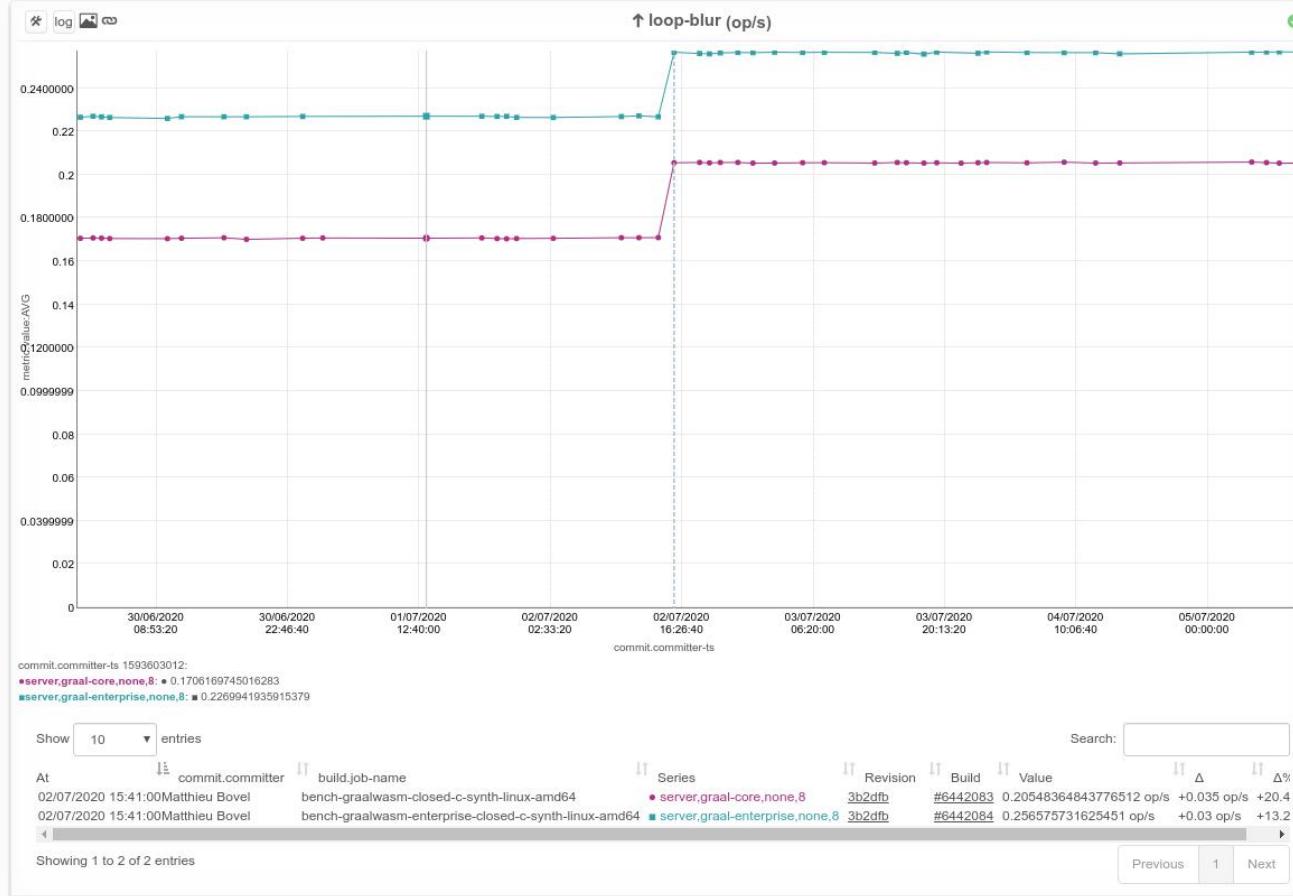
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- 21% improvement on a discrete-event simulation benchmark (*event-sim*)
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- 21% improvement on a discrete-event simulation benchmark (*event-sim*)
- 16% improvement on the quicksort benchmark (*qsort*)
- 13% improvement on a hash-join benchmark (*hash-join*)
- 20% improvement on an image-processing microbenchmark (*loop-blur*)



We've only scratched the surface.

# Internships in all areas of GraalVM

# Internship Program

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## Our Research Center Locations



Zurich, Switzerland



Linz, Austria



California, USA



Prague, Czech Republic



Brno, Czech Republic



Lviv, Ukraine



Casablanca, Morocco



Belgrade, Serbia

<https://www.graalvm.org/community/internship/>

# Thank you!

Questions?