

Code generation exercises

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

small: local variable 0
big: local variable 1
mid: local variable 2

iload_addr, istore_addr, iadd, idiv, isub, iconst_2

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

```
[val mid = small + (big - small) / 2; mid]  
ireturn
```

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

```
[ [val mid = small + (big - small) / 2] ] 12  
iload_2  
ireturn
```

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

```
[ [small] ] [ [big] ] [ [small] ] [ [-] ] [ [2] ] [ [/] ] [ [+]  
istore_2  
iload_2  
ireturn
```

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

```
[ [small] ]  [ [big] ]  [ [small] ]  [ [-] ]  [ [2] ]  [ [/] ]  
iadd  
istore_2  
iload_2  
ireturn
```

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

iload_0	idiv
iload_1	iadd
iload_0	istore_2
isub	iload_2
iconst_2	ireturn

Function body

Transform the following code into java bytecode:

```
def middle(small: Int, big: Int): Int = {  
    val mid = small + (big - small) / 2  
    return mid  
}
```

iload_0	idiv
iload_1	iadd
iload_0	ireturn
isub	
iconst_2	

binarySearch

```
def binarySearch(array: Array[Int], value: Int, left: Int,  
right: Int): Int = {  
    if (left > right)  
        return -1  
    val middle = (left + right) / 2  
    if (array(middle) == value)  
        return middle  
    else  
        if (array(middle) > value)  
            return binarySearch(array, value, left, middle - 1)  
        else  
            return binarySearch(array, value, middle + 1, right)  
}
```

binarySearch - bit.ly/1aCuKZz

Stack state:

```
def binarySearch(array: Array[Int], value: Int, left: Int,  
right: Int): Int
```

Methods :

#1 : binarySearch

Local variables mapping:

0:object itself

1:array

2:value

3:left

4:right

5:mid

binarySearch - bit.ly/1aCuKZz

```
[[if (left > right) return -1; ... ]  
     iload_3  
     iload_4  
     if_icmple goto after1:  
     iconst_1  
     ineg  
     goto return  
after1: [[...]]  
return: ireturn
```

binarySearch - bit.ly/1aCuKZz

```
[ [val middle = (left + right) / 2; ...] ]  
          => left right + 2 /  
  
after1: iload_3  
        iload_4  
        iadd  
        iconst_2  
        idiv  
        istore_5  
        [ [...] ]  
  
return: ireturn
```

binarySearch - bit.ly/1aCuKZz

```
[ [if (array(middle) == value) return middle else ...] ]
    aload_1
    iload_5
    iaload
    iload_2
    if_cmpne goto after2:
    iload_5
    goto return
after2: [ [...] ]
return: ireturn
```

binarySearch - bit.ly/1aCuKZz

```
[[if (array(middle) > value)
return binarySearch(array, value, left, middle - 1)
else ...]]

after2:  aload_1
         iload_5
         iaload
         iload_2
         if_cmple goto after3:
         aload_0 // Object itself
         aload_1
         iload_2
         iload_3
         iload_5
         icanst_1
         isub
         invokevirtual #1
         goto return:

after3: [[...]]
return: ireturn
```

binarySearch - bit.ly/1aCuKZz

```
[[return binarySearch(array, value, middle+1, right)]]  
after3:  aload_0  
        aload_1  
        iload_2  
        iload_5  
        iconst_1  
        iadd  
        iload_4  
        invokevirtual #1  
return: ireturn
```

Branching conditions

```
Boolean b;  
  
int f(int x, int y, int z) {  
    while (( !b && (x > 2*(y+z)) )  
    || (x < 2*y + z) ) {  
        x = x + 3  
    }  
    return x;  
}
```

Use destination passing style

Context: b is field 0 of object. (aload, getfield)
x -> 1, y -> 2, z -> 3

Branching conditions

```
Boolean b;

int f(int x, int y, int z) {
    while ( (!b && (x > 1 * (y+z)) || (x < 2*y + z)) {
        x = x + 3
    }
    return x;
}

lLoop: branch(condition, body, lAfter)
body: [x=x+3]
      goto lLoop
lAfter: iload_0
         ireturn
```

Translating the body

[x=x+3]

```
iload_1  
iconst_3  
iadd  
istore_1
```

Translating complex branching

```
( (!b && (x > 2 * (y+z)) || (x < 2*y + z) ) {  
lLoop: branch(c1 || c2, body, lAfter)  
=>
```

```
lLoop: branch(c1, ???, ???)  
c1No: branch(c2, ???, ???)  
body: [x=x+3]  
      goto lLoop  
lAfter: iload_0  
        ireturn
```

Translating complex branching

```
( (!b && (x > 2 * (y+z)) || (x < 2*y + z) ) {  
lLoop: branch(c1 || c2, body, lAfter)  
=>
```

```
lLoop: branch(c1, body, c1No)  
c1No: branch(c2, body, lAfter)  
body: [x=x+3]  
      goto lLoop  
lAfter: iload_0  
        ireturn
```

Translating complex branching

```
( !b && (x > 2 * (y+z)) {  
lLoop: branch(c11 && c12, body, c1No)  
=>  
  
lLoop: branch(c11, ???, ???)  
c11Yes: branch(c12, ???, ???)  
c1No: branch(x < 2*y + z, body, lAfter)  
body: [x=x+3]  
      goto lLoop  
lAfter: iload_0  
        ireturn
```

Translating complex branching

```
( !b && (x > 2 * (y+z)) {  
lLoop: branch(c11 && c12, body, c1No)  
=>  
  
lLoop: branch(c11, c1Yes, c1No)  
c1Yes: branch(c12, body, c1No)  
c1No: branch(x < 2*y + z, body, lAfter)  
body: [x=x+3]  
      goto lLoop  
lAfter: iload_0  
        ireturn
```

Translating complex branching

```
lLoop: branch(!b, c1Yes, c1No)
c1Yes: branch(x > 2 * (y+z), body, c1No)
c1No: branch(x < 2*y + z, body, lAfter)
body: [x=x+3]
      goto lLoop
lAfter: iload_0
        ireturn
```

Translating complex branching

```
lLoop: branch(b, c1No, c1Yes)
c1Yes: branch(x > 2 * (y+z), body, c1No)
c1No: branch(x < 2*y + z, body, lAfter)
body: [x=x+3]
      goto lLoop
lAfter: iload_0
        ireturn
```

Translating complex branching

```
lLoop: branch(b, c1No, c1Yes)
```

```
c1Yes: [x]
```

```
[2 * (y+z) ]
```

```
if_cmpgt body
```

```
c1No: [x]
```

```
[2*y + z]
```

```
if_cmplt body
```

```
goto lAfter)
```

```
body: [x=x+3]
```

```
goto lLoop
```

```
lAfter: iload_0
```

```
ireturn
```

Translating complex branching

```
lLoop:  [b]
        [0]
        if_cmpne c11No
c11Yes: [x]
        [2 * (y+z)]
        if_cmplt body
c11No:  [x]
        [2*y + z]
        if_cmplt body
        goto lAfter
body:   [x=x+3]
        goto lLoop
lAfter: iload_0
        ireturn
```

Translating complex branching

```
lLoop:  aload_0
        getfield
        iconst_0
        if_cmpne c11No
c11Yes: iload_1
        [2]  [y]  [z]  [+]
        if_cmplt body
c1No:   iload_1
        [2]  [y]  [*]  [z]  [+]
        if_cmplt body
        goto lAfter)
body:   [x=x+3]
        goto lLoop
lAfter: iload_0
        ireturn
```

Translating complex branching

```
lLoop:  aload_0                      istore_1
        getfield
        iconst_0
        if_cmpne c11No
        goto lLoop
c11Yes: iload_1
         iconst_2
         iload_2
         iload_3
         iadd
         imul
         if_cmpgt body
c1No:   iload_1
         iconst_2
         iload_2
         imul
         iload_3
         iadd
         if_cmplt body
         goto lAfter
body:   iload_1
         iconst_3
         iadd
```

Designing Code Generators

- Can we design a byte-code translation for the construct

```
repeat {  
    S1  
} until (cond)
```

The statement S1 should be repeatedly executed as long as the condition does not hold.

Repeat-until construct

[repeat { S1 } until(c)] after =

top: [S1]

branch(c, after, top)

Designing Code Generators 2

- Design a byte-code translation for a switch expression defined as follows:

```
switch(e) {  
    case c1 => e1  
    case c2 => e2  
    ...  
    case cn => en  
}
```

- If the value of the expression ‘e’ is equal to ‘ci’ (for some i) then ‘ei’ will be the result of the switch expression. If ‘e’ does not equal any ‘ci’ then ‘e’ will be the result of the expression

Code generation for switch

```
[e]                                end: //the value of 'e' will be in the top of stack  
case1: dup  
       [c1]  
       ifcmpne case2  
       pop  
       [e1]  
       goto end  
  
case2: dup  
       ifcmpne case3  
       pop  
       [e2]  
       goto end  
  
...
```

‘dup’ duplicates the value on the top of the stack and
‘pop’ pops the element at the top of the stack