

# Exercises on Chomsky Normal Form and CYK parsing

1. Convert the following grammar to CNF

$S \rightarrow A(S)B \mid \epsilon$

$A \rightarrow S \mid SB \mid x \mid \epsilon$

$B \rightarrow SB \mid y$

## Exercise 2

Which of the following properties of a grammar are preserved by the “Epsilon Production Elimination” algorithm

- **Ambiguity** No, counter-example ?
  - If a grammar is “ambiguous” does it remain ambiguous after removing epsilon productions ?
- **LL(1)** No, counter-example ?
- **No Left Recursion** No, counter-example ?
- **Unambiguity ?** Yes, proof ?

# Exercise 2

## [ Part 2 ]

Which of the following properties of a grammar are preserved by the “Unit Production Elimination” algorithm

- Ambiguity No, counter-example ?
- Left Recursion May be
  - What about these rules:  $B \rightarrow A \mid a$  ,  $A \rightarrow B$  ?
- LL(1) Yes, proof ?
- Unambiguity Yes, proof ?

## Exercise 3

Given a grammar  $G$  in CNF, how many steps does it require to derive a string of size  $n$ .

## Exercise 4

Assume a grammar in CNF has  $n$  non-terminals. Show that if the grammar can generate a word with a derivation having at least  $2^n$  steps, then the recognized language should be infinite

## Exercise 5

Show the CYK parsing table for the string “aabbab”  
for the grammar

$S \rightarrow AB \mid BA \mid SS \mid AC \mid BD$

$A \rightarrow a$

$B \rightarrow b$

$C \rightarrow SB$

$D \rightarrow SA$

What should be done to construct a parse tree for the  
string

## Exercise 6

Should a grammar be strictly in CNF form for CYK to work ? If not, what are the properties that can be relaxed ?

# A CYK for Any Grammar

grammar  $G$ , non-terminals  $A_1, \dots, A_K$ , tokens  $t_1, \dots, t_L$

input word:  $w = w_{(0)}w_{(1)} \dots w_{(N-1)}$

$w_{p..q} = w_{(p)}w_{(p+1)} \dots w_{(q-1)}$

Triple  $(A, p, q)$  means:  $A \Rightarrow^* w_{p..q}$ ,  $A$  can be:  $A_i$ ,  $t_j$ , or  $\varepsilon$

$P = \{(w_{(i)}, i, i+1) \mid 0 \leq i < N-1\}$

repeat {

  choose rule  $(A ::= B_1 \dots B_m) \in G$

  if  $((A, p_0, p_m) \notin P \ \&\&$

$((m=0 \ \&\& \ p_0=p_m) \ || \ (B_1, p_0, p_1), \dots, (B_m, p_{m-1}, p_m) \in P))$

$P := P \cup \{(A, p_0, p_m)\}$

} until no more insertions possible

Accept if  $(S, 0, n)$  belongs to  $P$

What is the maximal number of steps?

How long does it take to check step for a rule?

} for grammar in given normal form

# Observation

- How many ways are there to split a string of length  $Q$  into  $m$  segments?

$$\binom{Q+m}{m} = \frac{(Q+m)!}{Q!m!}$$

- Exponential in  $m$ , so algorithm is exponential.
- For binary rules,  $m=2$ , so algorithm is efficient.

# Closure Properties of CFG

- Concatenation

- If  $L1$  and  $L2$  are context-free languages is  $L = \{xy \mid x \in L1, y \in L2\}$  also context-free ?

- Union

- Closure

- Complement ?

- Not always a CFG, but sometimes possible

- We can convert any regular expression to a CFG

# Exercise 7

- Compute the complement of the grammar
  - $S \rightarrow A \mid B$
  - $A \rightarrow a A \mid \epsilon$
  - $B \rightarrow b B \mid \epsilon$