

Repair in the Leon tool

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May 1, 2017

- Verification problem:
Given a **correct specification** and an **implementation**, prove if the implementation is correct or not (for every input)
- Synthesis problem:
Given a **correct specification** and **no implementation**, come up with a correct implementation
- Repair problem:
Given a **correct specification** and an **erroneous implementation**, come up with a correct implementation.

Specification: either a logical formula, or input-output examples.

Example: Max Heap merging as a verification problem

Input:

```
def merge(h1: Heap, h2: Heap) : Heap = {
  require(isLegalHeap(h1) && isLegalHeap(h2))
  (h1,h2) match {
    case (Leaf(), _) => h2
    case (_, Leaf()) => h1
    case (Node(v1, l1, r1), Node(v2, l2, r2)) =>
      if(v1 ≤ v2)
        Node(v2, l2, merge(h1, r2))
      else
        Node(v1, l1, merge(r1, h2))
  }
} ensuring { res =>
  isLegalHeap(res) &&
  h1.content ++ h2.content == res.content
}
```

Output: Correct for every input!

Example: Max Heap merging as a synthesis problem

With logical specification:

```
def merge(h1: Heap, h2: Heap) : Heap = {  
  require(isLegalHeap(h1) && isLegalHeap(h2))  
  choose( (res: Heap) ⇒  
    isLegalHeap(res) &&  
    h1.content ++ h2.content == res.content  
  )  
}
```

Example: Max Heap merging as a synthesis problem

With examples:

```
def merge(h1: Heap, h2: Heap) : Heap = {  
  require(isLegalHeap(h1) && isLegalHeap(h2))  
  choose( (res: Heap)  $\Rightarrow$   
    ((h1, h2), res) passes {  
      case (Leaf(), Leaf())  $\Rightarrow$  Leaf()  
      case (Leaf(), Node(0, Leaf(), Leaf()))  $\Rightarrow$   
        Node(0, Leaf(), Leaf())  
      case (  
        Node(1, Leaf(), Leaf()),  
        Node(0, Leaf(), Leaf())  
      )  $\Rightarrow$   
        Node(  
          1,  
          Leaf(),  
          Node(0, Leaf(), Leaf()))))}}
```

Output: Implementation of previous slide.

Example: Max Heap merging as a repair problem

Input:

```
def merge(h1: Heap, h2: Heap) : Heap = {  
  require(isLegalHeap(h1) && isLegalHeap(h2))  
  (h1,h2) match {  
    case (Leaf(), _) => h2  
    case (_, Leaf()) => h1  
    case (Node(v1, l1, r1), Node(v2, l2, r2)) =>  
      if(v1 ≥ v2)  
        Node(v2, l2, merge(h1, r2))  
      else  
        Node(v1, l1, merge(r1, h2))  
  }  
} ensuring { res =>  
  isLegalHeap(res) &&  
  h1.content ++ h2.content == res.content  
}
```

Example: Max Heap merging as a repair problem

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  (h1,h2) match {
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    case (_, Leaf()) => h1
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      if(v1 ≥ v2)
        Node(v2, l2, merge(h1, r2))
      else
        Node(v1, l1, merge(r1, h2))
  }
} ensuring { res =>
  isLegalHeap(res) &&
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```

Output: the above code where \geq has been replaced with \leq .

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- Try to find a small change that would fix the snippet. If that fails, throw it away and write it from scratch.
- Rerun the test suite (or verifier!). If there are still issues, repeat.

Stages of the repair algorithm

- Test generation and (trace) minimization
- Fault Localization
- Synthesis of similar expressions
- Verification of the solution

Test generation and verification

Our algorithm needs at least one *failing test*, which leads to erroneous program execution

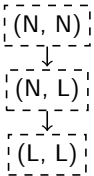
We obtain tests from various sources:

- Input-output examples given by the user
- Enumeration of programs
- Counterexamples from SMT solver

Trace minimization

In the presence of recursive functions, a given test may fail within one of its recursive invocations.

```
def merge(h1: Heap, h2: Heap) : Heap = {  
  require(isLegalHeap(h1) && isLegalHeap(h2))  
  (h1,h2) match {  
    case (Leaf(), _) => h1 // Buggy  
    ...  
  }  
}
```

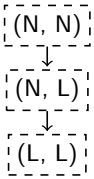


The diagram illustrates the sequence of recursive calls. It starts with a state (N, N) in a dashed box. An arrow points down to a state (N, L) in a dashed box. Another arrow points down to a state (L, L) in a dashed box.

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  (h1,h2) match {  
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    ...  
  }  
}
```



A failing test should also be blamed for the failure of all other tests that invoke it transitively.

In this case, only (Leaf(), Leaf()) is maintained as a failing example.

Follow the trace of failing tests to find in which branch of the program they lead us.

Suppose we have identified as failing tests:

Node(1, Leaf(), Leaf()), Node(0, Leaf(), Leaf())

Node(2, Leaf(), Leaf()), Node(0, Leaf(), Leaf())

```
(h1,h2) match {  
  case (Leaf(), _)  $\Rightarrow$  h2  
  case (_, Leaf())  $\Rightarrow$  h1  
  case (Node(v1, l1, r1), Node(v2, l2, r2))  $\Rightarrow$   
    if(v1  $\geq$  v2)  
      Node(v2, l2, merge(h1, r2))  
    else  
      Node(v1, l1, merge(r1, h2))  
}
```


A realistic set of failing tests is

Node(1, Leaf(), Leaf()), Node(0, Leaf(), Leaf())

Node(0, Leaf(), Leaf()), Node(1, Leaf(), Leaf())

```
(h1,h2) match {  
  case (Leaf(), _)  $\Rightarrow$  h2  
  case (_, Leaf())  $\Rightarrow$  h1  
  case (Node(v1, l1, r1), Node(v2, l2, r2))  $\Rightarrow$   
    if(v1  $\geq$  v2)  
      Node(v2, l2, merge(h1, r2))  
    else  
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```

Focusing on if-condition

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To find out, replace the condition with `havoc` and run the tests,
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test.

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To find out, replace the condition with `havoc` and run the tests,
i.e. nondeterministically consider both branches of the `if` for each
test.

If testing succeeds now,
i.e. there exists a valid execution exists for each failing test,
it means that the answer to the question is true.

Synthesis - Term Grammars

We have localized the error on the if-condition. Now, we have to synthesize an alternative solution.

We describe interesting programs with a term grammar. For repair, the grammar should describe small variations to the original program and simple arbitrary programs.

E.g.

$Boolean ::= Int \geq \mathbf{v2} \mid \mathbf{v1} \geq Int \mid \mathbf{v2} \geq \mathbf{v1} \mid \mathbf{true} \mid \mathbf{false} \mid \dots$

Synthesis - The CEGIS algorithm

Once we have a grammar representing interesting programs, we can synthesize a solution with the CEGIS algorithm.

Basic idea of CEGIS:

- 1 Use concrete tests to filter out candidate programs.
- 2 From those remaining, pick one and send it to the verifier.
- 3 If verification successful, we are done
- 4 Otherwise, the verifier generates a counterexample.
Add it to the set of tests and jump back to (1).
(note: step (1) will now filter out more programs)

CEGIS will generate $v2 \geq v1$ and verify it as the correct solution

Demos!