

Exercises 3

1 Loop semantics

Compute and simplify the relation corresponding to the following programs:

```
y = 0
while (z > 0) {
  y = y + x
  z = z - 1
}

c = 0
while (b >= a) {
  b = b - a
  c = c + 1
}
```

2 Loop invariants

In the following program, provide the necessary loop invariant for verification to succeed:

```
def binarySearch(a: Array[BigInt], key: BigInt): Int = ({
  require(a.length > 0 && forall { (i: Int, j: Int) =>
    (i >= 0 && j >= 0 && i < a.length && j < a.length && i < j) ==> (a(i) <= a(j))
  })

  var low = 0
  var high = a.length - 1
  var res = -1

  (while(low <= high && res == -1) {
    val o = if ((high & 1) == 1 && (low & 1) == 1) 1 else 0
    val i = high / 2 + low / 2 + o
    val v = a(i)

    if(v == key)
      res = i

    if(v > key)
      high = i - 1
    else if(v < key)
      low = i + 1
  }) invariant(TODO)
  res
}) ensuring(res ==> {
  if(res == -1)
    forall((i: Int) => (0 <= i && i < a.length) ==> (a(i) != key))
  else
    a(res) == key
})
```

3 Proof construction

Show that $\text{list flatMap } f \text{ flatMap } g == \text{list flatMap } (x \Rightarrow f(x) \text{ flatMap } g)$ using the following axioms:

1. $\text{Nil flatMap } f == \text{Nil}$
2. $(x :: xs) \text{ flatMap } f == f(x) ++ (xs \text{ flatMap } f)$
3. $\text{Nil} ++ xs == xs$
4. $xs ++ \text{Nil} == xs$
5. $(x :: xs) ++ ys == x :: (xs ++ ys)$

Use the proof strategies you have seen in *Welder*, such as structural induction and equational reasoning. Make sure each step in your reasoning is clearly indicated.

Hint: It may be useful to introduce some auxiliary lemmas.