

Example to Type Check

```

object World {
  var z : Boolean
  var u : Int
  def f(y : Boolean) : Int {
    z = y
    if (u > 0) {
      u = u - 1
      var z : Int
      z = f(!y) + 3
      z+z
    } else { 0 }
  }
}

```

$$\Gamma_0 = \{$$

(z, Boolean),
 (u, Int),
 (f, Boolean \rightarrow Int) }

$$\Gamma_1 = \Gamma_0 \oplus \{(y, \text{Boolean})\}$$

$$\frac{\Gamma_1 \vdash z: \text{Boolean} \quad \Gamma_1 \vdash y: \text{Boolean}}{\Gamma_1 \vdash (z=y): \text{void}}$$

Exercise:

$$\frac{\text{???}}{\Gamma_1 \vdash \text{if } (n > 0) \{ \text{body} \} \text{ else } \{ 0 \} : \text{Int}}$$

Overloading of Operators

$\text{Int} \times \text{Int} \rightarrow \text{Int}$

$$\frac{\Gamma \vdash e_1: \text{Int} \quad \Gamma \vdash e_2: \text{Int}}{\Gamma \vdash (e_1 + e_2): \text{Int}}$$

Not a problem for type checking from leaves to root

$\text{String} \times \text{String} \rightarrow \text{String}$

$$\frac{\Gamma \vdash e_1: \text{String} \quad \Gamma \vdash e_2: \text{String}}{\Gamma \vdash (e_1 + e_2): \text{String}}$$

Arrays

Using array as an expression, on the right-hand side

$$\frac{\Gamma \vdash a: \text{Array}(T) \quad \Gamma \vdash i: \text{Int}}{\Gamma \vdash a[i]: T}$$

Assigning to an array

$$\frac{\Gamma \vdash a: \text{Array}(T) \quad \Gamma \vdash i: \text{Int} \quad \Gamma \vdash e: T}{\Gamma \vdash (a[i]) = e: \text{void}}$$

Example with Arrays

```
def next(a : Array[Int], k : Int) : Int = {  
  a[k] = a[a[k]]  
}
```

Given $\Gamma = \{(a, \text{Array}(\text{Int})), (k, \text{Int})\}$, check $\Gamma \vdash a[k] = a[a[k]] : \text{Int}$

$$\frac{\Gamma \vdash a : \text{Array}(\text{Int}) \quad \frac{\Gamma \vdash a : \text{Array}(\text{Int}) \quad \Gamma \vdash k : \text{Int}}{\Gamma \vdash a[k] : \text{Int}}}{\Gamma \vdash a[a[k]] : \text{Int}} \quad \frac{\Gamma \vdash a : \text{Array}(\text{Int}) \quad \Gamma \vdash k : \text{Int}}{\Gamma \vdash a[k] = a[a[k]] : \text{Int}}$$

Type Rules (1)

$$\frac{(x: T) \in \Gamma}{\Gamma \vdash x: T} \quad \text{variable}$$

$$\frac{}{\text{IntConst}(k): \text{Int}} \quad \text{constant}$$

$$\frac{\Gamma \vdash e_1 : T_1 \quad \dots \quad \Gamma \vdash e_n : T_n \quad \Gamma \vdash f : (T_1 \times \dots \times T_n \rightarrow T)}{\Gamma \vdash f(e_1, \dots, e_n) : T} \quad \text{function application}$$

$$\frac{\Gamma \vdash e_1 : \text{Int} \quad \Gamma \vdash e_2 : \text{Int}}{\Gamma \vdash (e_1 + e_2) : \text{Int}} \quad \text{plus} \quad \frac{\Gamma \vdash e_1 : \text{String} \quad \Gamma \vdash e_2 : \text{String}}{\Gamma \vdash (e_1 + e_2) : \text{String}}$$

$$\frac{\Gamma \vdash b : \text{Boolean} \quad \Gamma \vdash e_1 : T \quad \Gamma \vdash e_2 : T}{\Gamma \vdash (\text{if}(b) e_1 \text{ else } e_2) : T} \quad \text{if}$$

$$\frac{\Gamma \vdash b : \text{Boolean} \quad \Gamma \vdash s : \text{void}}{\Gamma \vdash (\text{while}(b) s) : \text{void}} \quad \text{while}$$

$$\frac{(x, T) \in \Gamma \quad \Gamma \vdash e : T}{\Gamma \vdash (x=e) : \text{void}} \quad \text{assignment}$$

Type Rules (2)

$$\frac{\Gamma \vdash e: T}{\Gamma \vdash \{e\}: T}$$

$$\frac{}{\Gamma \vdash \{\}: \text{void}}$$

$$\frac{\Gamma \oplus \{(x, T_1)\} \vdash \{t_2; \dots; t_n\}: T}{\Gamma \vdash \{\text{var } x : T_1; t_2; \dots; t_n\}: T}$$

block

$$\frac{\Gamma \vdash s_1: \text{void} \quad \Gamma \vdash \{t_2; \dots; t_n\}: T}{\Gamma \vdash \{s_1; t_2; \dots; t_n\}: T}$$

$$\frac{\Gamma \vdash a: \text{Array}(T) \quad \Gamma \vdash i: \text{Int}}{\Gamma \vdash a[i]: T}$$

array use

$$\frac{\Gamma \vdash a: \text{Array}(T) \quad \Gamma \vdash i: \text{Int} \quad \Gamma \vdash e: T}{\Gamma \vdash a[i] = e}$$

array
assignment

Type Rules (3)

Γ^c - top-level environment of class C

```
class C {  
  var x: Int;  
  def m(p: Int): Boolean = {...}  
}
```



$\Gamma^c = \{ (x, \text{Int}), (m, C \times \text{Int} \rightarrow \text{Boolean}) \}$

$$\frac{\Gamma \vdash e : C \quad \Gamma^C \vdash m : T \times T_1 \times \dots \times T_n \rightarrow T_{n+1} \quad \Gamma \vdash e_i : T_i \quad 1 \leq i \leq n}{\Gamma \vdash e.m(e_1, \dots, e_n) : T_{n+1}} \quad \text{method invocation}$$

$$\frac{\Gamma \vdash e : C \quad \Gamma^C \vdash f : T}{\Gamma \vdash e.f : T} \quad \text{field use}$$

$$\frac{\Gamma \vdash e : C \quad \Gamma^C \vdash f : T \quad \Gamma \vdash x : T}{\Gamma \vdash (e.f = x) : \text{void}} \quad \text{field assignment}$$