

Recall the example from the course:

```
class World {
    boolean z;
    int u;
    int f(boolean y) {
        z = y;
        if (u > 0) {
            int z;
            z = f(u) + 3;
            return z+z;
        } else {
            return 0;
        }
    }
}
```

Let Γ_{env} be the environment obtained from the *Class* declaration (with methods and global variables):

$$\Gamma_0 = (z : \text{Bool}), (u : \text{Int}), (f : \text{Bool} \rightarrow \text{Int})$$

Let us type check the code inside the method *f*: we add a binding for *f*'s parameter in the environment:

$$\Gamma_1 = \Gamma_0 \oplus (y : \text{Bool})$$

We obtain the tree Tree_0 :

$$\begin{array}{c} \text{Tree}_0 : \\ \frac{\begin{array}{c} (z, \text{Bool}) \in \Gamma_1 & (y, \text{Bool}) \in \Gamma_1 & \frac{\begin{array}{c} (u, \text{Int}) \in \Gamma_1 & \frac{\vdash 0 : \text{Int}}{\Gamma_1 \vdash u : \text{Int}} & \frac{\vdash 0 : \text{Int}}{\Gamma_1 \vdash 0 : \text{Int}} \\ \hline \Gamma_1 \vdash u > 0 : \text{Bool} & \Gamma_1 \vdash \{ \text{int } z; z = f(u) + 3; \text{return } z + z; \} : \text{Int} & \Gamma_1 \vdash \{ \text{return } 0; \} : \text{Int} \end{array}}{\Gamma_1 \vdash \{ \text{int } z; z = f(u) + 3; \text{return } z + z; \} \text{ else } \{ \text{return } 0; \} : \text{Int}} \\ \hline \Gamma_1 \vdash z = y : \text{void} & \Gamma_1 \vdash \{ \text{int } z; z = f(u) + 3; \text{return } z + z; \} \text{ else } \{ \text{return } 0; \} : \text{Int} \\ \hline \Gamma_1 \vdash z = y; \text{ if } (u > 0) \{ \text{int } z; z = f(u) + 3; \text{return } z + z; \} \text{ else } \{ \text{return } 0; \} : \text{Int} \end{array} \quad \frac{\text{Tree}_1}{\Gamma_1 \vdash 0 : \text{Int}} \end{array}$$

Unfortunately, our program is not type correct. We can see it trying to complete the derivation (Tree_1):

$$\begin{array}{c} \text{Tree}_1 : \\ \frac{\begin{array}{c} (u, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int}) & \text{!!!}(f, \text{Int} \rightarrow \text{Int}) \notin \Gamma_1 \oplus (z, \text{Int}) \text{!!!} \\ \hline \frac{\begin{array}{c} \Gamma_1 \oplus (z, \text{Int}) \vdash u : \text{Int}, f : \text{Int} \rightarrow \text{Int} & \frac{\vdash 3 : \text{Int}}{\Gamma_1 \oplus (z, \text{Int}) \vdash 3 : \text{Int}} \\ \hline \frac{\begin{array}{c} \Gamma_1 \oplus (z, \text{Int}) \vdash f(u) : \text{Int} & \frac{\begin{array}{c} (z, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int}) & \frac{(z, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int})}{\Gamma_1 \oplus (z, \text{Int}) \vdash z : \text{Int}} \\ \hline \frac{\vdash 3 : \text{Int}}{\Gamma_1 \oplus (z, \text{Int}) \vdash z + z : \text{Int}} & \frac{\text{Gamma closed}}{\Gamma_1 \oplus (z, \text{Int}) \vdash \text{return } z + z : \text{Int}} \end{array}}{\Gamma_1 \oplus (z, \text{Int}) \vdash f(u) + 3 : \text{Int}} \\ \hline \Gamma_1 \oplus (z, \text{Int}) \vdash z = f(u) + 3 : \text{void} & \Gamma_1 \oplus (z, \text{Int}) \vdash \{ z = f(u) + 3; \text{return } z + z; \} : \text{Int} \end{array} \end{array} \end{array}$$

If we change the call $f(u)$ in $f(u > 0)$, we can type our program:

$$\begin{array}{c}
 \text{Tree}_1 : \\
 \frac{\frac{\frac{(u, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int})}{\Gamma_1 \oplus (z, \text{Int}) \vdash u : \text{Int}} \quad \frac{\vdash 0 : \text{Int}}{\Gamma_1 \oplus (z, \text{Int}) \vdash 0 : \text{Int}} \quad (f, \text{Bool} \rightarrow \text{Int}) \in \Gamma_1 \oplus (z, \text{Int})}{\Gamma_1 \oplus (z, \text{Int}) \vdash u > 0 : \text{Bool}, f : \text{Bool} \rightarrow \text{Int}} \quad \frac{\vdash 3 : \text{Int}}{\Gamma_1 \oplus (z, \text{Int}) \vdash 3 : \text{Int}}}{\Gamma_1 \oplus (z, \text{Int}) \vdash f(u > 0) + 3 : \text{Int}} \\
 \frac{(z, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int})}{\Gamma_1 \oplus (z, \text{Int}) \vdash z : \text{Int}} \quad \frac{\Gamma_1 \oplus (z, \text{Int}) \vdash f(u > 0) + 3 : \text{Int}}{\Gamma_1 \oplus (z, \text{Int}) \vdash z = f(u > 0) + 3 : \text{void}}}{\Gamma_1 \oplus (z, \text{Int}) \vdash \{z = f(u > 0) + 3; \text{return } z + z;\} : \text{Int}}
 \end{array}$$

$$\frac{\frac{(z, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int})}{\Gamma_1 \oplus (z, \text{Int}) \vdash z : \text{Int}} \quad (z, \text{Int}) \in \Gamma_1 \oplus (z, \text{Int})}{\Gamma_1 \oplus (z, \text{Int}) \vdash z + z : \text{Int}}}{\Gamma_1 \oplus (z, \text{Int}) \vdash \text{return } z + z : \text{Int}}$$

Which rule did we add to the ones presented in the course in order to derive those trees?